

MG14 02/23/05

P. 1

Consider interpolating a function within the interval $[-1, 1]$. Let N be an integer and

x_0, x_1, \dots, x_N be a set of distinct points in $[-1, 1]$.

Say $1 \geq x_0 > x_1 > x_2 > \dots > x_{N-1} > x_N \geq -1$

Equispaced points : $x_j = 1 - \frac{2j}{N}$ $0 \leq j \leq N$

Chebyshev zero points : $x_j = \cos \frac{(j-1/2)\pi}{N}$ $1 \leq j \leq N$

Chebyshev collocation pts : $x_j = \cos \frac{j\pi}{N}$ $0 \leq j \leq N$
(extreme)

Legendre zero points : $x_j = j$ th zero of P_N ($1 \leq j \leq N$)

Legendre extreme points : $x_j = j$ th extreme of P_N ($0 \leq j \leq N$)

density of nodes :

$N \rightarrow \infty$ $\mu(x) = \frac{N}{2}$ for equally spaced points

For Chebyshev zero points

$$\frac{\partial x_j}{\partial j} = - \frac{\sin(j - \frac{1}{2})\pi/N \cdot (\frac{\pi}{N})}{\sin(j - \frac{1}{2})\pi/N}$$

$$\text{density } \mu(x_j) = - \frac{1}{\frac{\partial x_j}{\partial j}} = \frac{N/\pi}{\sin(j - \frac{1}{2})\pi/N} = \frac{N/\pi}{\sqrt{1-x_j^2}}$$

$$\therefore \boxed{\mu(x) = \frac{N}{\pi \sqrt{1-x^2}}}$$

Why is it a good idea to use Chebyshev zeros or extreme pts for grids to interpolate / approximate a function?

Why not use equally spaced grids?

(polynomial)
If we use trigonometric interpolation in equispaced points,

say for an ^{analytic} ~~periodic~~ function f in $[-1, 1]$

Gibbs phenomenon: $\|f - P_N\| = \mathcal{O}(1)$ as $N \rightarrow \infty$
(Runge) even if f is analytic
 $= \mathcal{O}(2^N)$ for polynomial interpolation

\Rightarrow doubling the domain and reflecting (periodic boundary conditions) can reduce Gibbs phenomenon.

Using Chebyshev points or Legendre points

$\|f - P_N\| \sim \mathcal{O}(\text{constant}^{-N})$ if f is analytic

Definition of Lebesgue constant

$$\Lambda_N = \|I_N\|_\infty$$

I_N is the interpolation operator $I_N: f \mapsto P_N$

A small Lebesgue constant means that the interpolation is not much worse than best (minimax) approx.

$$\|f - P_N\| \leq (\Lambda_N + 1) \|f - P_N^*\|$$

P_N^* is the best (minimax) approx.

Theorem:

For Equispaced points: $\Lambda_N \sim 2^N / eN \log N$

Chebyshev pts: $\Lambda_N \sim \log N$

Legendre pts: $\Lambda_N \sim \sqrt{N}$

Say $f_N(x) = \cos(\alpha N x)$ is an approximation as $N \rightarrow \infty$.

$f_N(x)$ changes but the number of points per wavelength remains constant.

Will $\|f_N - P_N\| \rightarrow 0$ as $N \rightarrow \infty$

For equispaced points: convergence is achieved if

there are at least 6 pts / wavelength

For Chebyshev points: convergence if π points / wavelength
(on average)

on average because the grid is non-uniform.

$\Rightarrow \frac{\pi}{2}$ times less dense in the middle than the equispaced pts
with the same number of N pts.

\Rightarrow in the center, only $\frac{2 \text{ pts}}{\text{wave length}}$ are needed

\Rightarrow familiar Nyquist limit

Example: Show that $x^n = \sum_{j=0}^n C_{j,n} P_j(x)$

for appropriate coefficients $\{C_{j,n}\}$
 the Legendre polynomial $P_j(x)$ is an orthogonal polynomial family

because $(P_i, P_j) = \delta_{ij}$

$$\Rightarrow \begin{bmatrix} P_0(x) \\ P_1(x) \\ P_2(x) \\ \vdots \\ P_n(x) \end{bmatrix} = \begin{bmatrix} a_{00} & a_{01} & \dots & a_{0n} \\ a_{10} & a_{11} & \dots & a_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n0} & a_{n1} & \dots & a_{nn} \end{bmatrix} \begin{bmatrix} 1 \\ x \\ \vdots \\ x^n \end{bmatrix}$$

$$= A \cdot \begin{bmatrix} 1 \\ x \\ \vdots \\ x^n \end{bmatrix}$$

$$\det(A) \neq 0$$

$$\therefore \begin{bmatrix} 1 \\ x \\ \vdots \\ x^n \end{bmatrix} = A^{-1} \begin{bmatrix} P_0 \\ P_1 \\ \vdots \\ P_n \end{bmatrix} \quad \therefore C_{ij} = (A^{-1})_{ij}$$

$$\Rightarrow \begin{aligned} P_0(x) &= 1 \\ P_1(x) &= x \\ P_2(x) &= -\frac{1}{2} + \frac{3}{2}x^2 \end{aligned} \quad \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -\frac{1}{2} & 0 & \frac{3}{2} \end{bmatrix} \begin{bmatrix} 1 \\ x \\ x^2 \end{bmatrix}$$

$$\begin{bmatrix} 1 \\ x \\ x^2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ \frac{1}{3} & 0 & \frac{2}{3} \end{bmatrix} \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$

$$\therefore \begin{cases} 1 = P_0 \\ X = P_1 \\ X^2 = \frac{1}{3}P_0 + \frac{2}{3}P_2 \end{cases}$$

How about $n=3$?

$$\begin{bmatrix} P_0 \\ P_1 \\ P_2 \\ P_3 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ -\frac{1}{2} & 0 & \frac{3}{2} & 0 \\ 0 & -\frac{3}{2} & 0 & -\frac{5}{2} \end{bmatrix} \begin{bmatrix} 1 \\ X \\ X^2 \\ X^3 \end{bmatrix}$$

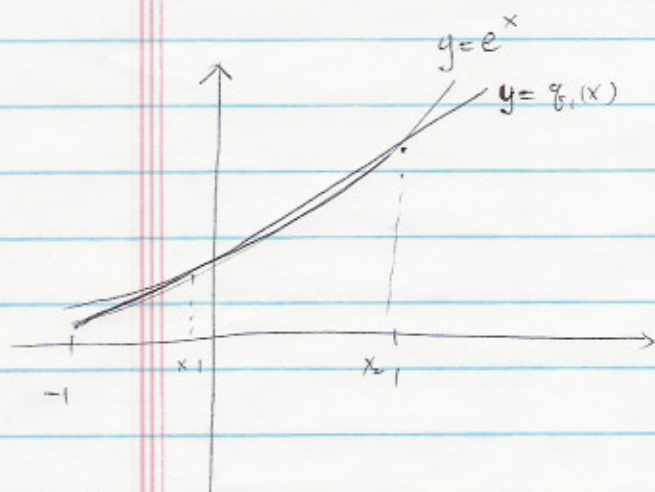
$$\therefore \begin{bmatrix} 1 \\ X \\ X^2 \\ X^3 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ \frac{1}{3} & 0 & \frac{2}{3} & 0 \\ 0 & \frac{3}{5} & 0 & \frac{2}{5} \end{bmatrix} \begin{bmatrix} P_0 \\ P_1 \\ P_2 \\ P_3 \end{bmatrix}$$

$$\therefore \begin{cases} 1 = P_0 \\ X = P_1 \\ X^2 = \frac{1}{3}P_0 + \frac{2}{3}P_2 \\ X^3 = \frac{3}{5}P_1 + \frac{2}{5}P_3 \end{cases}$$

Minimax Approximation

\Rightarrow minimax polynomial approx. to e^x on $[-1, 1]$

Let $g_1^*(x) = a_0 + a_1x$ be the linear minimax approx.



define the error $\varepsilon(x)$ as

$$\varepsilon(x) = e^x - (a_0 + a_1 x)$$

$$\rho_1 \equiv \max_{-1 \leq x \leq 1} (\varepsilon(x))$$

shift $g_1^*(x)$ s.t.

$$\varepsilon(-1) = \varepsilon(1) = -\varepsilon(x_3) = \rho_1$$

with $x_1 < x_3 < x_2$

$$\therefore e^{-1} - (a_0 - a_1) = \rho_1$$

$$e^{x_3} - (a_0 + a_1 x_3) = -\rho_1$$

$$e^1 - (a_0 + a_1) = \rho_1$$

$$\varepsilon'(x_3) = 0, \quad e^{x_3} - a_1 = 0$$

four unknowns

four equations

$$\boxed{a_1 = \frac{e^1 - e^{-1}}{2}}, \quad \boxed{x_3 = \ln\left(\frac{e - e^{-1}}{2}\right)}$$

$$e^{x_3} - e - a_1 x_3 + a_1 = -2\rho_1$$

$$2a_1 - e - a_1 x_3 = -2\rho_1$$

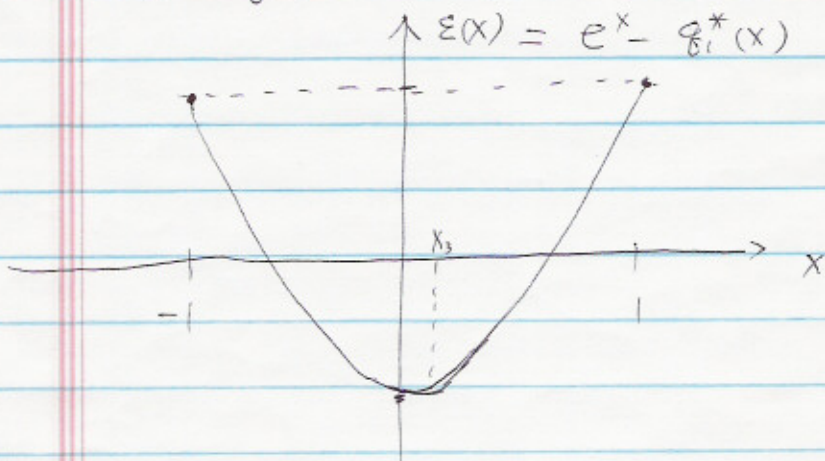
$$\boxed{\rho_1 = \frac{1}{2}e^{-1} + \frac{x_3}{4}(e - e^{-1})}$$

≈ 0.2788

$$a_0 = \rho_1 + (1 - x_3)a_1$$

$$\therefore g_1^*(x) = 1.2643 + 1.1752x$$

$$\varepsilon(x) = e^x - g_1^*(x)$$



- error is evenly distributed over the interval
- error is oscillating between $+P_n$ & $-P_n$

Thm: Let $f \in C[a, b]$ and $n \geq 0$.

Suppose we have a polynomial $Q(x)$ of degree $\leq n$ which satisfies

$$f(x_j) - Q(x_j) = (-1)^j e_j \quad j = 0, 1, \dots, n+1$$

with all e_j non zero and of the same sign and with
 $a \leq x_0 \leq x_1 \leq \dots \leq x_{n+1} \leq b$

$$\text{Then } \min_{0 \leq j \leq n+1} |e_j| \leq P_n(f) \equiv \underbrace{\|f - Q_n^*\|_\infty}_{\text{by definition}} \leq \|f - Q\|_\infty$$

$$\text{Assume } P_n(f) < \min_{0 \leq j \leq n+1} |e_j|$$

By definition of $P_n(f)$, there exists a polynomial $P(x)$ of degree $\leq n$
 s.t.

$$P_n(f) \leq \|f - P\|_\infty < \min |e_j|$$

all $e_j > 0$

$\Rightarrow R(x) \equiv Q(x) - P(x)$ a polynomial of degree $\leq n$.

$$\begin{aligned} R(x_0) &= Q(x_0) - P(x_0) = [f(x_0) - P(x_0)] - [f(x_0) - Q(x_0)] \\ &= [f(x_0) - P(x_0)] - e_0 < 0 \end{aligned}$$

Next

$$R(x_1) = Q(x_1) - P(x_1) = [f(x_1) - P(x_1)] - (-e_1) > 0$$

Therefore the sign of $R(x_j)$ is $(-1)^{j+1}$ for $j = 0, 1, \dots, n+1$

$\Rightarrow R(x)$ has $n+1$ zeros \Rightarrow since $\deg(R) \leq n$, $R \equiv 0$.

$$P \equiv Q$$

This contradicts

$$\rho_n(f) \leq \|f - P\|_{\infty} < \min |e_j|$$

 \Rightarrow because ~~$\rho_n(f)$~~

$$f(x_j) - Q(x_j) = (-1)^j e_j$$

 \Rightarrow therefore

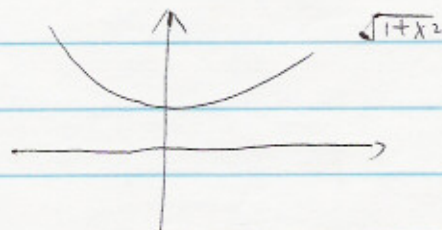
$$\rho_n(f) \geq \min_{0 \leq j \leq n+1} |e_j|$$

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(a) show that the linear minimax approx. to $\sqrt{1+x^2}$ on $[0, 1]$

$$\text{is } g_1^*(x) = .955 + .414x$$

$$\text{write } g_1^*(x) = a_0 + a_1x$$

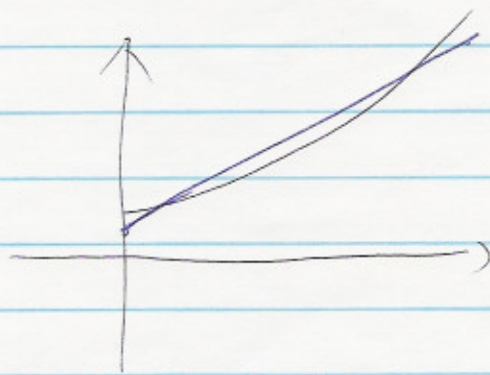


$$\sqrt{1+1} - (a_0 + a_1) = \rho_1$$

$$\sqrt{1+0} - (a_0 + 0) = \rho_1$$

$$\sqrt{1+x_3^2} - (a_0 + a_1x_3) = -\rho_1$$

$$\frac{1 \cdot 2x_3}{2\sqrt{1+x_3^2}} - a_1 = 0$$



$$a_1 = \sqrt{2} - 1 = .414, \quad x_3 = \left(\frac{1}{2+2\sqrt{2}} \right)^{1/2} = 0.45509$$

$$a_0 = \frac{1}{2} (1 + \sqrt{1+x_3^2} - a_1x_3)$$

$$a_0 = .95514$$

$$\rho_1 = 1 - a_0 = .04486$$

Probl (b)

$$\sqrt{y^2 + z^2} = z \cdot \sqrt{1 + (y/z)^2} \quad 0 \leq y \leq z$$

$$\sqrt{1 + (y/z)^2} \sim g_1^*(y/z) = .955 + .414 \left(\frac{y}{z}\right)$$

$$z \cdot \sqrt{1 + (y/z)^2} \sim z \cdot g_1^*(y/z) = .955z + .414y$$

the error $z \cdot g_1 = z \cdot .04486$ is a function of z .

It is easy to remember how Chebyshev points are defined: they are the projections onto the interval $[-1, 1]$ of equally-spaced points (roots of unity) along the unit circle $|z| = 1$ in the complex plane:

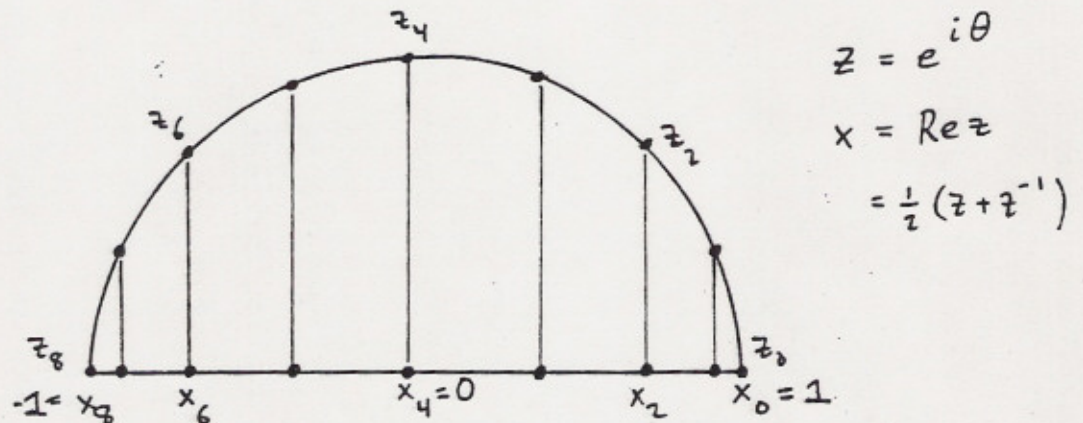


Figure 8.1.1. Chebyshev extreme points ($N = 8$).

To the eye, Legendre points look much the same, although there is no elementary geometrical definition. Figure 8.1.2 illustrates the similarity:

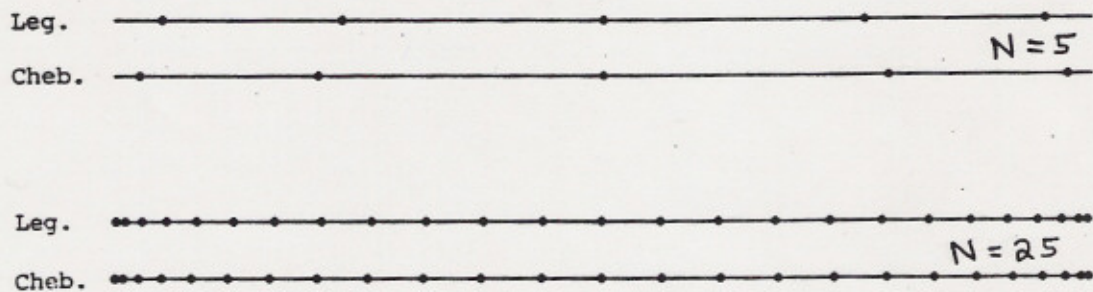


Figure 8.1.2. Legendre vs. Chebyshev zeros.

As $N \rightarrow \infty$, equispaced points are distributed with density

$$\mu(x) = \frac{N}{2} \quad \text{Equally spaced,} \quad (8.1.2)$$

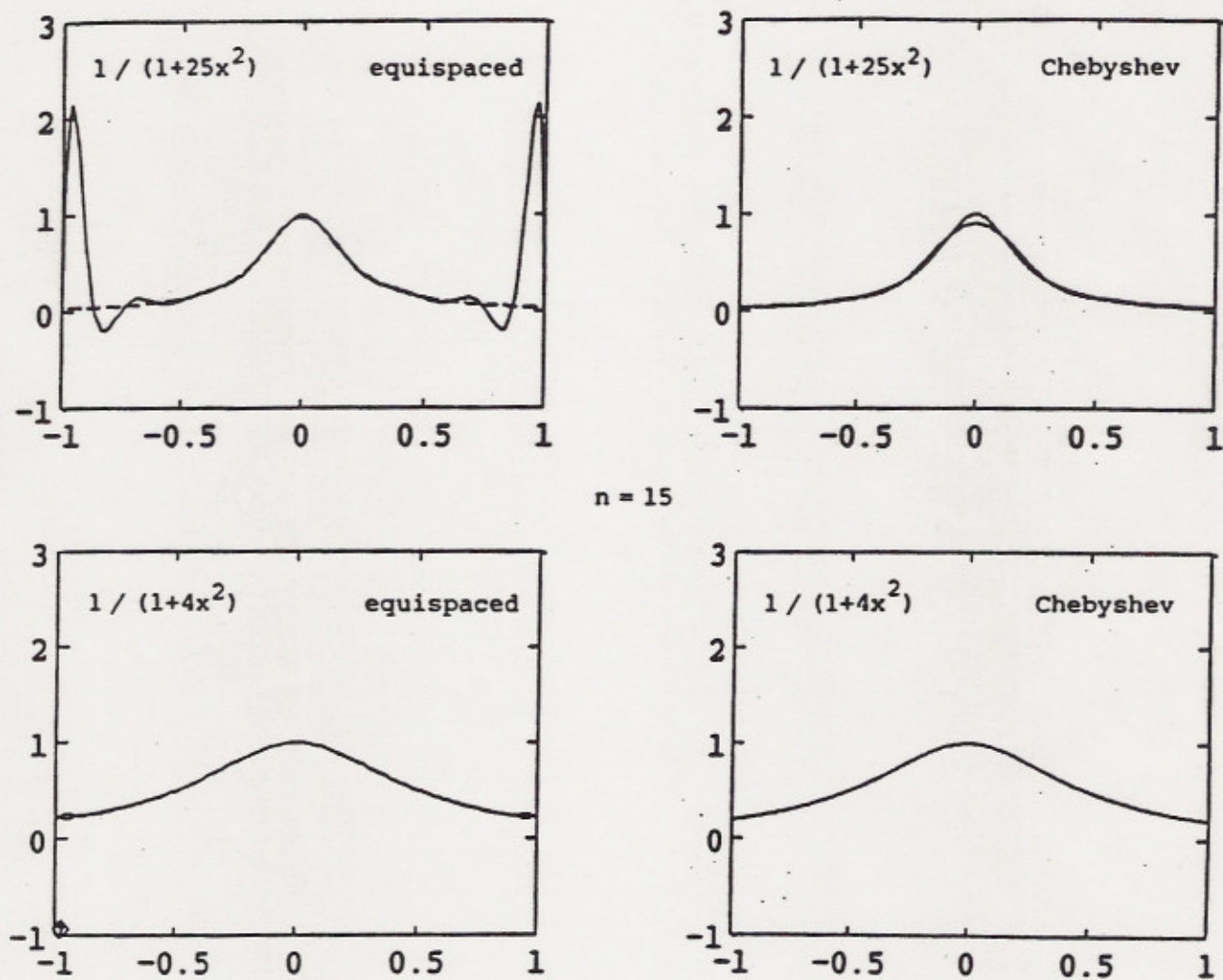
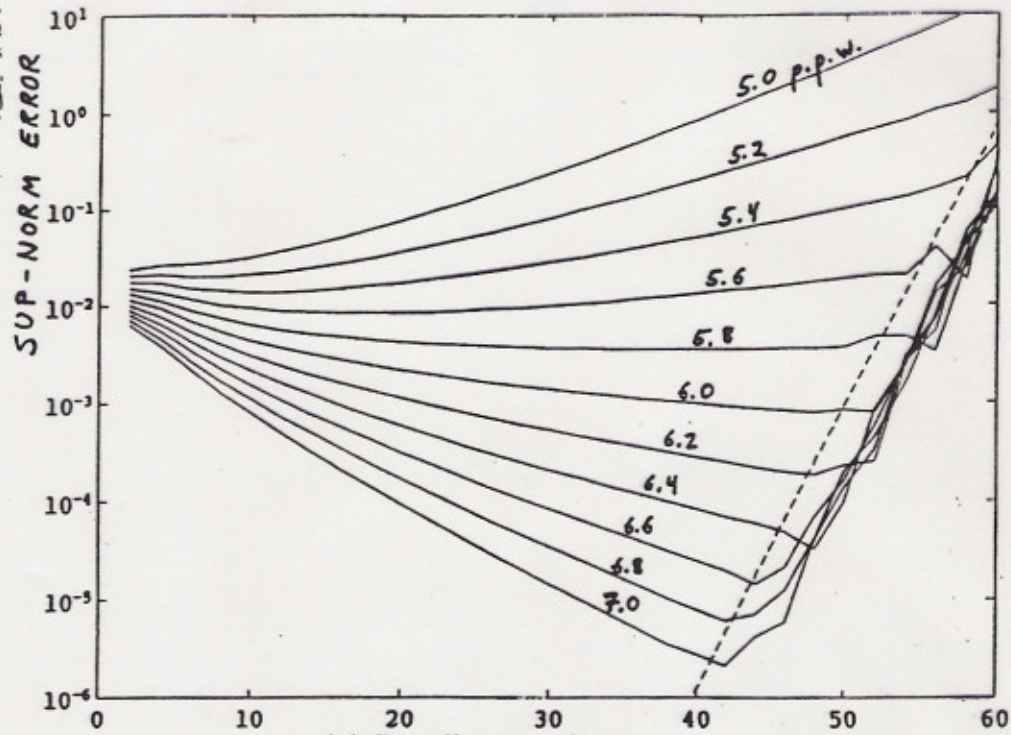
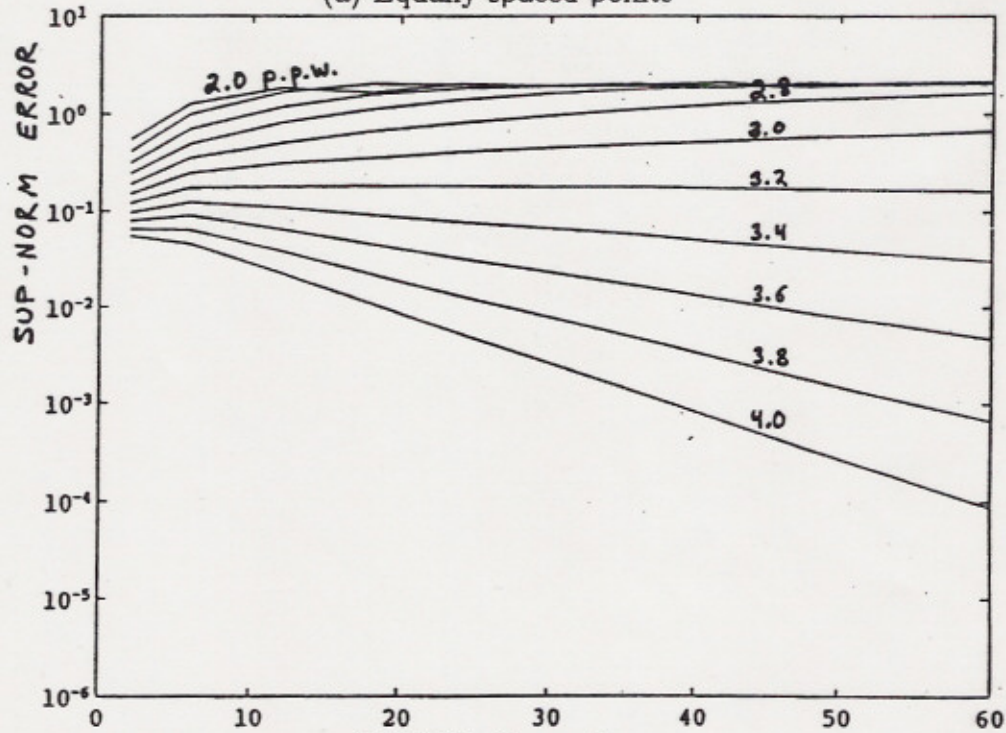


Figure 8.1.3. The Runge phenomenon.



(a) Equally spaced points



(b) Chebyshev points

Figure 8.1.5. Error as a function of N in interpolation of $\cos \alpha Nx$, with α , hence the number of points per wavelength, held fixed.