Generalized Discrete Fourier Transform*

Non-Linear Phase DFT for Improved Multicarrier Communications

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TUTORIAL

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Outline

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Outline

VI. BER Performance Comparisons of DFT and GDFT CDMA for Rayleigh Fading Channels
VII. Variations of CDMA Communications: From DS-CDMA to MC-DC-CDMA
VIII. Emerging 3GPP LTE Mobile Phone Standard
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I. Introduction to Orthogonal Block Transforms:
            A Time-Frequency Perspective
Function/Signal
Shape of a function/signal
Function/Signal = Energy Shape
Energy of a function/signal
Duality (Time-Frequency)
Parseval (Time-Frequency)
Function set (Time-Frequency)
Orthogonality (Time-Frequency)
DFT, DCT, WHT
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Functions (Energy Shape) in Time and Frequency Domains: Duality





















Signal Energy in Time-Frequency: Parseval Theorem

Let
$$\phi_k(n) \longleftrightarrow_{DTFT} \Phi_k(e^{j\omega})$$

 $\Phi_k(e^{j\omega}) = \sum_{k=1}^{\infty} \phi_k(n) e^{-j\omega n}$

 $n = -\infty$

$$\phi_k(n) = \frac{1}{2\pi} \int_{-\pi}^{\pi} \Phi_k(e^{j\omega}) e^{j\omega n} d\omega$$

$$E_{\phi_{k}} = \sum_{n=-\infty}^{n=\infty} \phi_{k}(n)\phi_{k}^{*}(n) = \sum_{n=-\infty}^{n=\infty} |\phi_{k}(n)|^{2}$$

$$=\frac{1}{2\pi}\int_{-\pi}^{\pi}\Phi_{k}(e^{j\omega})\Phi_{k}^{*}(e^{j\omega})d\omega=\frac{1}{2\pi}\int_{-\pi}^{\pi}\left|\Phi_{k}(e^{j\omega})\right|^{2}d\omega$$



Orthogonality in Time-Frequency: Parseval Theorem

Let
$$\phi_k(n) \leftarrow DTFT \rightarrow \Phi_k(e^{j\omega})$$

 $\phi_l(n) \leftarrow DTFT \rightarrow \Phi_l(e^{j\omega})$
Consider a function set $\{\phi_k(n)\}; k = 0, 1, ..., N - 1$
 $\sum_{n=-\infty}^{n=\infty} \phi_k(n) \phi_l^*(n) = \frac{1}{2\pi} \int_{-\pi}^{\pi} \Phi_k(e^{j\omega}) \Phi_l^*(e^{j\omega}) d\omega = \delta(k-l)$
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Correlations in Time-Frequency: Multiuser Communications (Transmultiplexer for OFDM/CDMA/TDMA)

Let
$$\phi_k(n) \longleftrightarrow_{DTFT} \Phi_k(e^{j\omega})$$

 $\phi_l(n) \longleftrightarrow_{DTFT} \Phi_l(e^{j\omega})$

Consider a function set $\{\phi_k(n)\}; k = 0, 1, ..., N-1$ and pairwise cross-correlations defined

$$R_{\phi_k\phi_l}(m) = \sum_{n=-\infty}^{n=\infty} \phi_k(n)\phi_l^*(n-m) = \frac{1}{2\pi} \int_{-\pi}^{\pi} \Phi_k(e^{j\omega})\Phi_l^*(e^{j\omega})e^{j\omega m}d\omega$$

Correlations result in Inter Carrier Interference (ICI), Inter Symbol Imterference (ISI) and MultiUserInterference (MUI) in a Multiuser / Multicarrier Communications System (Transmultiplexer for OFDM/CDMA/TDMA or T-FMA?)



Time Frequency Localization of A Discrete Time Function



Discrete-Time Uncertainty



$$\sigma_{n}\sigma_{\omega} \geq \frac{|1-\mu|}{2}$$

$$\mu \triangleq \frac{|H(e^{j\pi})|^{2}}{E}$$
Class I: $H(e^{j\pi}) = 0 \rightarrow \sigma_{n}\sigma_{\omega} \geq \frac{1}{2}$
Class II: $H(e^{j\pi}) \neq 0 \rightarrow \sigma_{n}\sigma_{\omega} \geq \frac{|1-\mu|}{2}$

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Time Frequency Localization of 31-Length Spread Spectrum KLT Codes

•Binary valued codes have constant spread in time where as multiple valued KLT codes are more spread in time

•Frequency spread for Walsh-like, Gold and KLT code sets is similar where as for Walsh codes spread is lower



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Orthonormal Spectral Analyzer as a Filter Bank



Orthogonality Principle Demonstration



Transform Encoder / Decoder





2D Transform Coding





KLT Basis [AR(1), rho=0.95] in Time & Frequency (N=8)





















 $\frac{y(n)}{f_s/M}$

n

n

71

3

x(n) .

2M

x'(n)

2M

y(n)

2



Interpolation Operation









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4-Band / Two-Levels Filter Bank (Subband) Tree



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Irregular Subband Tree





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A Dyadic (Octave Band) Subband Tree





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Maximally Decimated M-Band Filter Bank (Analysis/Synthesis) Single Input Single Output (SISO)



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M-Band Transmultiplexer (Synthesis/Analysis FB Configuration) Multiple Input Multiple Output (MIMO)



II. Discrete Fourier Transform with Linear Phase

$$e_{k}(n) = e^{j(2\pi/N)kn} \qquad k, n = 0, 1, \dots, N-1$$

$$\frac{1}{N} \sum_{n=0}^{N-1} e_{k}(n) e_{l}^{*}(n) = \frac{1}{N} \sum_{n=0}^{N-1} e^{j(2\pi/N)(k-l)n} = \begin{cases} 1, & k-l = r = mN \\ 0, & k-l = r \neq mN \\ m, n = \text{int } eger \end{cases}$$





DFT Phase Functions (Linear)

Modulo 2π



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III. Orthogonal Transmultiplexer for Multicarrier Communications: OFDMA, TDMA, CDMA







Non-linear Phase (Walsh-like) Binary Orthogonal Codes: Design and Performance

- Walsh codes are linear phase, zero mean with unique number of zero crossings in the set. DC code is part of the set
- Features are useful for source coding and not necessary for spread spectrum applications. Hence, such design restrictions are waived in Walsh-like code design.
- For *n*-length binary code, sample space consists of integer numbers up to 2^n -1. First basis function is selected by representing any integer number in the sample space in radix 2 format with [1,-1] elements



Walsh-like Codes Design

- Select the next basis function by checking the orthogonality with the first basis function and maximum normalized cross-correlation value between the pair is less than 1 for all possible delays
- Repeat this process (n 1) times to get *n* orthogonal codes
- A number of orthogonal sets are formed with first basis function as common basis function for all the orthogonal sets

Walsh-like Codes Design

- By choosing different integer as the first basis function, unique orthogonal sets can be formed. Number of *n x n* orthogonal sets with multiples of 4 as their lengths are obtained in our simulation (8,12,16,20,...)
- Complexity of this algorithm is $n.(2^n-1)$ for *n*-length code


Asynchronous BER Performance Comparison of 8-Length Walsh-like Code Sets in AWG Noise(2 Users)

BER performance of 8length Walsh-like codes is marginally better than 7length Gold codes and far exceeds that of Walsh codes
Number of orthogonal code

sets are available with similar performance



Asynchronous BER Performance Comparison of 32-Length Walsh-like Code Sets in AWG Noise(2 Users)

•BER performance of Walsh-like codes exceeds Walsh codes and closely matches with Gold codes performance at all lengths of the code



Async. BER Performance of 16,20,24,28,32 Length Walsh-like Codes in AWG Noise (2 Users)

• Number of Walsh-like code sets are available for all lengths that are multiples of 4



Rayleigh Flat - Slow Fading Channel Description

- Multipath reflections of the symbol occur in the same symbol interval. This implies coherent bandwidth of the channel is greater than the symbol bandwidth (Flat fading)
- Channel conditions are assumed to remain same during symbol interval (Slow fading)



Rayleigh Flat - Slow Fading Channel Description

• Amplitude of the received signal modeled as y(t) = h(t)*s(t) + n(t),

s(t) transmitted signal, n(t) AWG noise, h(t) channel impulse response and y(t) received signal



Rayleigh Flat - Slow Fading Channel Description

- For flat fading channel, h(t) single tap filter with zero delay
- *h(t)* WSS complex Gaussian waveform with zero mean and unity variance whose amplitude varies as Rayleigh PDF variable
- Fading channel modeled separately for each user in uplink



Sync / Async BER Performance Comparison for Length-32 Walsh-like Codes (2 Users)-Rayleigh Channel

•Performance of orthogonal Walsh and Walsh-like codes is similar in all Rayleigh flat fading conditions

•Performance of nonorthogonal Gold codes is poor in synchronous conditions



Multiple Level Code : Design and Performance

- For p level coding, sample space is pⁿ for an n-length code. Represent numbers in sample space using radix p elements. Map radix elements into corresponding PAM level chip amplitudes
- 4 level representation requires radix 4 elements (0,1,2,3) and PAM chip levels {-3, -1, 1, 3}. Weights of the individual elements for 8-length code are {4⁷, 4⁶,..., 4¹, 4⁰}



Multiple Level Code : Design and Performance

- As an example, number 125 in radix 4 is represented as {0,0,0,0,1,3,3,1}. After PAM mapping, code becomes {-3, -3, -3, -1, 3, 3, -1}
- Number of unique orthogonal code sets are obtained by brute force search method in the sample space with code features similar to Walsh-like code sets
- In the search algorithm, additional constraint of same norm for all basis functions within the set is also imposed



Asynchronous BER performance Comparison of 4-Level , 4-Length Codes in AWG Noise(2 Users)

BER performance of 2- level, 4-length orthogonal codes (Walsh) is poor where as 4-level codes give good performance
Sample 4-level, 4-length code

$$\begin{bmatrix} -3 & -1 & 1 & 3 \\ -1 & 3 & -3 & 1 \\ -1 & 3 & 3 & -1 \\ 3 & 1 & 1 & 3 \end{bmatrix}$$



Asynchronous BER Performance Comparison of Multiple Level, 8-Length Codes in AWG Noise (2 Users)

•BER performance improves as the number of coding levels increase



Multiple Level, 6-Length Codes

3-Level	5-Level	7-Level	9-Level	11-Level	13-Level
Codes	Codes	Codes	Codes	Codes	Codes
Basis	Basis	Basis	Basis	Basis	Basis
Elements	Elements	Elements	Elements	Elements	Elements
{-1,0,1}	{-2, -1, 0, 1, 2}	{-3, -2, -1, 0, 1, 2, 3}	{-4, -3, -2, -1,	{-5, -4, -3, -2, -1,	{6,5,4,
Norm ² -4	Norm ² -10	Norm ² -20	0,1,2,3,4}	0,1,2,3,4,5}	-3, -2, -1, 0,
			Norm ² -26	Norm ² -50	Norm²-50
22	182	2858	9832	17968	95570
396	8378	51974	163032	143082	2433484
404	8476	60218	218158	847144	2469152
524	10660	73438	272578	847218	2779564
528	10714	77530	310372	885134	3144890
670	12932	103798	372786	885216	3356242



Asynchronous BER Performance of Multiple Level, 6-Length Codes in AWG Noise (2 Users)

•BER performance improves as the number of levels increase upto a certain level



Normalized Cross Correlation Metrics for Multiple Level 6-Length Codes

Parameter	3	5	7	9	11	13
	Level	Level	Level	Level	Level	Level
Max Even Correlation	.75	.7	0.8	.73	.64	.72
Max Odd Correlation	.75	.8	.65	.73	.64	.70
Max Aperiodic	.75	.7	0.7	.73	.64	.70
Correlation						
Sum of Square of Even	11	13.1	10	13	14	13
Correlations						
Sum of Square of Odd	12	13.5	13	13	13	14
Correlations						
Sum of Square of	12	13.3	11	13	14	13
Aperiodic Correlations						



Comparison of Multiple Level and Binary Level Codes

•Shorter length codes with higher chip levels perform as good as longer length codes with binary spreading levels



Spread Spectrum KLT Codes

- Integer binary and multiple level spread spectrum codes are obtained by brute force search method
- Karhunen-Loeve Transform (KLT) based analytical method is used to generate multiple value spread spectrum codes for a given covariance or power spectral density (PSD) function
- PSD function can be modeled using AR, ARMA methods, giving many solutions with variable code lengths

Example: Power Spectral Density and 8-Length Auto Correlation sequence

•Power spectral density is first modeled as AR sequence

•Eigen vectors generated from covariance matrix are used as spread spectrum codes





Asynchronous 2 User BER Performance of 8-Length Spread spectrum KLT Codes in AWG Noise

•BER performance of multiple valued 8-length KLT codes is better than binary valued codes.



Typical 31-Length Gold Code and its Auto Correlation Sequence

-1	1	1	-1	-1	1	1	1
-1	-1	-1	-1	1	-1	-1	-1
-1	-1	1	-1	1	-1	1	-1
1	-1	-1	1	1	-1	-1	_

-0.0645 1.0000 -0.0645 -0.0323 0.0645 0.2258 -0.1935 0.2258 -0.0323 -0.1290 0.0000 0.0645 -0.0968 0.1290 0.0968 -0.03230.0968 0.0645 -0.0968 0.0000 -0.03230.1290 0.0968 0 -0.2258 0 0.1613 0 0.0000 0.0323 -0.0968



31-Length Spread Spectrum KLT Auto Correlation Sequence and Power Spectrum

•Normalized auto- and cross-correlation sequences of Gold codes or Walshlike sequences are used to generate eigen vectors



Asynchronous 2 User BER Performance of 31-Length, Spread Spectrum KLT Codes in AWG Noise

BER performance of spread spectrum KLT codes matches with Gold codes performance
Number of unique orthogonal sets can be generated by taking different auto- and cross-correlation sequences of Gold / Walsh-like codes



IV. Correlation Performance Metrics

Aperiodic Correlation Function (ACF)

$$d_{k,l}(m) = \begin{cases} \frac{1}{N} \sum_{n=0}^{N-1-m} e_k(n) e_l^*(n+m), 0 < m \le N-1 \\ \frac{1}{N} \sum_{n=0}^{N-1+m} e_k(n-m) e_l^*(n), 1-N < m \le 0 \\ 0 , |m| \ge N \end{cases}$$

Max of Auto- and Cross-Correlation Sequences

$$d_{\max} = \max\left\{d_{am}, d_{cm}\right\}$$

$$d_{am} = \max \left\{ \left| d_k(m) \right| \right\}$$
$$0 \le k < M$$
$$1 \le m < M$$

$$d_{cm} = \max \left\{ \begin{vmatrix} d_{k,l}(m) \end{vmatrix} \right\}$$

$$0 \le k, l < M \quad k \neq l$$

$$0 \le m < M$$

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MS of Auto- and Cross-Correlation Sequences

 $P_{AC} = \frac{1}{M} \sum_{k=1}^{M} \sum_{m=1-N}^{N-1} \left| d_{k,k}(m) \right|^{2}$ R $m \neq 0$ $=\frac{1}{M(M-1)}\sum_{k=1}^{M}\sum_{l=1}^{M}\sum_{m=1-N}^{M}\left|d_{k,l}(m)\right|^{2}$ $l \neq k$





MS of Auto- and Cross-Correlation Sequences

 $T_{AC} = \frac{1}{M} \sum_{k=1}^{M} \sum_{m=1-N}^{N-1} \left| d_{k,k}(m) \right|^{2}$ R $m \neq 0$ $=\frac{1}{M(M-1)}\sum_{k=1}^{M}\sum_{l=1}^{M}\sum_{m=1-N}^{M}\left|d_{k,l}(m)\right|^{2}$ $l \neq k$



V. GDFT with Nonlinear Phase for Auto- and Cross-Correlation Improvements

- Motivation
 - GDFT
- Design Metrics and EfficientImplementation
- Performance Improvements in BER,PAPR
- Dynamic Code Basis Hopping



Motivation

- Constant modulus transforms are valuable for several applications including communications
- DFT is very popular since everyone uses it
- Can I have my DFT-like transforms?
- Emerging radio applications particularly SW based ones for sensing and P2P communications might benefit from flexible code/carrier libraries
- Introducing dynamic code/basis assignments offers additional improvements for system security (scrambler)



An *Nth* root of unity is a complex number satisfying the equation

$$z^{N} = 1$$
 N = 1, 2, 3,

If Z hold this equation but

$$m z_p \neq 1$$
 m=1,2,3,..., N-1

then Z_p is defined as a primitive Nth root of unity where m and N are coprime integers

$$z_{1} = e^{j(2\pi/N)}$$

is the primitive Nth root of unity with the smallest positive argument.

The are N distinct Nth roots of unity and expressed as

$$z_k = (z_p)^k \quad k = 1, 2, 3, \dots, N \quad \forall p$$



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As an example, for N=4 there are two primitive Nth roots of unity expressed as

$$z_1 = e^{j\frac{2\pi}{4}}$$

$$z_2 = e^{j\frac{3\pi}{2}}$$

All primitive Nth roots of unity satisfy the unique summation property of a geometric series expressed as follows

$$\sum_{n=0}^{N-1} (z_p)^n = \frac{(z_p)^N - 1}{z_p - 1} = \begin{cases} 1, N = 1\\ 0, N > 1 \end{cases} \quad \forall p$$



We now define a periodic, with the period of N, constant modulus, complex sequence as the *rth power* of the first *primitive Nth roots of unity* raised to the *nth power* as

$$e_r(n) \triangleq (z_1^r)^n = e^{j(2\pi r/N)n}$$

 $n = 0, 1, 2, ..., N - 1 \text{ and } r = 0, 1, 2, ..., N - 1$



This complex sequence over a finite discrete-time interval in a geometric series is expressed

$$z_1 = e^{j\omega_0} \qquad \qquad \omega_0 = 2\pi/N$$

$$\frac{1}{N} \sum_{n=0}^{N-1} e_r(n) = \frac{1}{N} \sum_{n=0}^{N-1} (z_1^r)^n = \frac{1}{N} \sum_{n=0}^{N-1} e^{j(2\pi r/N)n}$$
$$= \begin{cases} 1, & r = mN \\ 0, & r \neq mN \\ & m = \text{integer} \end{cases}$$

One defines the discrete Fourier transform (DFT) set with the factorization into two orthogonal exponential functions where

$$< e_{k}(n), e_{l}^{*}(n) >= \frac{1}{N} \sum_{n=0}^{N-1} e_{k}(n) e_{l}^{*}(n) = \frac{1}{N} \sum_{n=0}^{N-1} e^{j(2\pi k/N)n} e^{-j(2\pi l/N)n}$$

$$= \frac{1}{N} \sum_{n=0}^{N-1} e^{j(2\pi/N)(k-l)n} = \begin{cases} 1, & k-l=r=mN \\ 0, & k-l=r\neq mN \\ m=\text{integer} \end{cases}$$
Let's generalize by introducing a product function in the phase defined as

 $\varphi(n) = \varphi_k(n) - \varphi_l(n)$

and expressing a constant amplitude orthogonal set as follows,



$$\frac{1}{N} \sum_{n=0}^{N-1} e^{j(2\pi r/N)n} = \frac{1}{N} \sum_{n=0}^{N-1} e^{j[2\pi\varphi(n)/N]n}$$

$$= \begin{cases} 1, \quad \varphi(n) = \varphi_k(n) - \varphi_l(n) = r = mN \\ 0, \quad \varphi(n) = \varphi_k(n) - \varphi_l(n) = r \neq mN \\ m, n = \text{ integer} \end{cases}$$

$$= \frac{1}{N} \sum_{n=0}^{N-1} e^{j[2\pi(\varphi_k(n) - \varphi_l(n))/N]n}$$

$$= \frac{1}{N} \sum_{n=0}^{N-1} e^{j[2\pi\varphi_k(n)/N]n} e^{-j[2\pi\varphi_l(n)/N]n}$$

$$= \frac{1}{N} \sum_{n=0}^{N-1} e_k(n) e_l^*(n) = \langle e_k(n), e_l^*(n) \rangle$$

Hence, the basis functions of Generalized DFT (GDFT) are defined as

$$e_{k}(n) \triangleq e^{j(2\pi/N)\varphi_{k}(n)n}$$

$$k, n = 0, 1, \dots, N-1$$

As an example, one might define

$$\varphi_{k}(n) = \frac{N_{k}(n)}{D_{k}(n)} = \frac{\sum_{j=1}^{N} a_{kj} n^{b_{kj}}}{\sum_{j=1}^{M} c_{kj} n^{d_{kj}}} \qquad N \le M; \quad k = 0, 1, \dots, N-1$$

 $D_k(n) = 1$

$$\varphi_k(n) = N_k(n) = \sum_{j=1}^N a_{kj} n^{b_{kj}} = a_{k1} n^{b_{k1}} + a_{k2} n^{b_{k2}} + a_{k3} n^{b_{k3}} + \dots + a_{kN} n^{b_{kN}}$$



1) DFT is a special solution of GDFT with

$$\varphi_k(n) = a_{k1} = k \text{ and } a_{k2} = a_{k3} = \dots = a_{kN} = 0$$

$$b_{k1} = b_{k2} = \dots = b_{kN} = 0$$

Having constant valued $\{\varphi_k(n)\}$ functions makes DFT a linear-phase transform

$$\{e_k(n)\} \triangleq e^{j(2\pi/N)kn}$$
 $k, n = 0, 1, ..., N-1$ **NJIT**

 Popular Walsh and Nonlinear Phase Walsh-like orthogonal binary transforms are special solutions of GDFT. As an example,

$$A_{WALSH} = A_{GDFT} = A_{DFT} G_{WALSH}$$

$$G_{WALSH} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0.71e^{j\pi 0.25} & 0.71e^{-j\pi 0.25} \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0.71e^{-j\pi 0.25} & 0.71e^{j\pi 0.25} \end{bmatrix}$$

 $\{b_{k1} = 0, b_{k2} = 1, b_{k3} = 2, b_{k4} = 3 \text{ for } k = 0, 1, 2, 3\}$

$$\begin{cases} \varphi_0^{\text{WALSH}}(n) = 0 \\ \varphi_1^{\text{WALSH}}(n) = 6.22n - 5.33n^2 + 1.11n^3 \\ \varphi_2^{\text{WALSH}}(n) = -1.278n + 1.667n^2 - 0.389n^3 \\ \varphi_3^{\text{WALSH}}(n) = 4.5n - 3n^2 + 0.5n^3 \end{cases} \quad n = 0, 1, 2, 3$$



3) There are infinitely many possible GDFT sets available in the phase space with constant power where one can design the optimal basis for the desired figure of merit.

The availability of rich library of orthogonal constant amplitude transforms with good performance allows us to design adaptive systems where basis assignments as well as code allocations are made dynamically and intelligently to exploit the current channel conditions in order to deliver better communications performance and improved physical layer security.

4) Oppermann, Frank-Zadoff and Chu Sequences are the special cases of his code family.

$$A_{GDFT} = A_{OPP} = A_{DFT}G_{OPP}$$

$$a_1 = \frac{k^m + kN}{2}$$

$$b_1 = 0; \quad a_2 = \frac{1}{2}; \quad b_2 = n - 1$$

$$A_{OPP}(k,i) = (-1)^{ki} \exp\left(\frac{j\pi(k^m i^p + i^n)}{N}\right)$$

$$k, i = 1, 2, \dots, N$$

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Variations of Auto-correlation Metric for Parametric GDFT Solutions (N=8)





Matrix Representation

$$A_{GDFT} = G_1 A_{DFT} G_2$$

$$A_{GDFT} A_{GDFT}^{*T} = I$$

$$G_1 G_1^{*T} = I \text{ and } G_2 G_2^{*T} = I$$



$$G_{1}(k,n) = \begin{cases} e^{j\theta_{kk}}, & k = n \\ 0, & k \neq n \\ k,n = 0,1,...N \end{cases}$$
$$G_{2}(k,n) = \begin{cases} e^{j\gamma_{nn}}, & n = k \\ 0, & n \neq k \\ k,n = 0,1,...N \end{cases}$$

GDFT Kernel (Diagonal G1 & G2)

$$A_{GDFT} = \{e_k(n)\} \triangleq \{e^{j[(2\pi/N)kn + \theta_{kk} + \gamma_{nn}]}\}$$

k, n = 0, 1, ..., N - 1



DESIGN METRICS

Aperiodic Correlation Function (ACF)

$$d_{k,l}(m) = \begin{cases} \frac{1}{N} \sum_{n=0}^{N-1-m} e_k(n) e_l^*(n+m), \ 0 < m \le N-1 \\ \frac{1}{N} \sum_{n=0}^{N-1+m} e_k(n-m) e_l^*(n), \ 1-N < m \le 0 \\ 0 & , & |m| \ge N \end{cases}$$

Max of Auto- and Cross-Correlation Sequences

$$d_{\max} = \max\left\{d_{am}, d_{cm}\right\}$$

$$d_{am} = \max \left\{ \left| d_k(m) \right| \right\}$$
$$0 \le k < M$$
$$1 \le m < M$$

$$d_{cm} = \max \left\{ \left| d_{k,l}(m) \right| \right\}$$

$$0 \le k, l < M \quad k \neq l$$

$$0 \le m < M$$

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MS of Auto- and Cross-Correlation Sequences

 $= \frac{1}{M} \sum_{k=1}^{M} \sum_{m=1-N}^{N-1} \left| d_{k,k}(m) \right|^{2}$ R $m \neq 0$ $=\frac{1}{M(M-1)}\sum_{k=1}^{M}\sum_{l=1}^{M}\sum_{m=1-N}^{N}\left|d_{k,l}(m)\right|^{2}$ $l \neq k$





MS of Auto- and Cross-Correlation Sequences

 $T_{AC} = \frac{1}{M} \sum_{k=1}^{M} \sum_{m=1-N}^{N-1} \left| d_{k,k}(m) \right|^{2}$ R $m \neq 0$ $=\frac{1}{M(M-1)}\sum_{k=1}^{M}\sum_{l=1}^{M}\sum_{m=1-N}^{N}\left|d_{k,l}(m)\right|^{2}$ $l \neq k$



Optimal Design of Phase Shaping Function

$$\hat{\varphi}_{k}(n) = \varphi_{k}(n)n = kn + \psi(n) \text{ for } k = 0, 1, ..., N - 1 \text{ and } n = 1, ..., N - 1$$
$$\psi(n) = \hat{\varphi}_{k}(n) - kn = [\varphi_{k}(n) - k]n \text{ for } k = 0, 1, ..., N - 1 \text{ and } n = 1, ..., N - 1$$
$$\psi(0) \in \mathbb{R} \quad \hat{\varphi}_{k}(0) = \psi(0)$$

Cross Correlation:

$$R_{\widehat{\varphi}_{k}\widehat{\varphi}_{l}}(m) = \sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})\widehat{\varphi}_{k}(n)} e^{-j(\frac{2\pi}{N})\widehat{\varphi}_{l}(n+m)}$$
$$= \sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})[-lm+(k-l)n+\psi(n)-\psi(n+m)]}$$

Ideal Case for CC (Requires Sinc Functions):

$$R_{\widehat{\varphi}_k \widehat{\varphi}_l}(m) = 0; \quad \forall m$$

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Optimal Design of Phase Shaping Function

$$\hat{\varphi}_k(n) = \varphi_k(n)n = kn + \psi(n)$$

Auto Correlation:

$$R_{\widehat{\varphi}_{k}\widehat{\varphi}_{k}}(m) = \sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})\widehat{\varphi}_{k}(n)} e^{-j(\frac{2\pi}{N})\widehat{\varphi}_{k}(n+m)}$$
$$= \sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})[-km+\psi(n)-\psi(n+m)]}$$

Ideal Case:

$$R_{\widehat{\varphi}_k\widehat{\varphi}_k}(m) = \delta(m)$$

NJIT

Optimal Design of Phase Shaping Function

The first two functions of optimal GDFT sets with N=8 along with their DFT counterparts

	OPTIMIZATION METRIC (N=8)			
Numerical Search Tool and Optimal Phase Shaping Function	R _{AC}	R _{CC}		
GDFT (Mathematica, FindMin)	0.0877	0.4219		
$\psi(n)$	{ -1.37 -2.53 -2.21 3.39 0.0 -4.21 -3.19 -0.83 }	{ 1.637 -0.79 -0.54 2.01 1.59 -0.83 1.73 2.44}		
GDFT (MATLAB,fminsearch)	0.086	0.4205		
$\psi(n)$	{ -1.38 -2.56 -2.24 3.42 0.07 - 4.27 -3.27 -0.80 }	{ 1.673 -0.87 -0.51 2.02 1.51 -0.86 1.70 2.46 }		
DFT	4.375	0.8536		



Closed Form Phase Shaping Function $\psi(n)$





Closed Form Phase Shaping Function for GDFT Design

$$\varphi_k(n) = kn + \psi(n)$$

$$\{a_1 = 1, b_1 = 1.75, c_1 = 3.75, a_2 = 1.75, b_2 = 6, c_2 = 0.5\}$$

$$\psi(n) = \exp\left(-\left(\frac{n-1.75}{3.75}\right)^2\right) + 1.75 \exp\left(-\left(\frac{n-6}{0.50}\right)^2\right)$$





Phase Functions: DFT (Linear) vs GDFT (Nonlinear)



DFT Phase Functions (Linear)

Modulo 2π



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N

GDFT Phase Functions (Nonlinear)







Correlation metrics: DFT vs GDFT



Performance of Various Codes

d_{am}, d_{cm})
	d_{am}, d_{cm}

Code	d _{am}	d _{cm}	d _{max}	R _{AC}	R _{CC}	F
Walsh [8x8]	0.875	0.875	0.875	2.375	0.661	0.421
Walsh-like [8x8]	0.625	0.625	0.625	0.875	0.875	1.143
DFT [8x8]	0.875	0.327	0.875	4.375	0.375	0.220
7/8 Gold	0.714	0.714	0.714	0.857	0.878	1.167
Opperrman (opt d _{max}) (m=1, p=1, n=2.98, N=7)	0.425	0.419	0.465	1.278	0.787	0.783
GDFT (opt d _{max})	0.376	0.387	0.387	1.095	0.843	0.912



Magnitude of Auto-correlation Functions for GDFT (solid line) and DFT (dashed line) (N=16)


Magnitude of Cross-correlation Functions for GDFT (solid line) and DFT (dashed line) (N=16)





DS-CDMA BER Performance (2 Users, AWGN)



AM/AM Characteristics of RF PA



GDFT-SLM BASED PAPR REDUCTION



PAPR Reduction Comparisons for N=256, $P(PAPR > \gamma)=10^{-1}$

	PAPR _{DFT}	PAPR ^{EFF} GDFT	PAPR ^{OPT} GDFT	PAPR _{SLM}	$\Delta PAPR[dB]$		
т	[dB]	[dB]	[dB]	[dB]	EFF-GDFT	OPT-GDFT	SLM
2	10.4	8.6	8.1	8.3	1.8	2.3	2.1
4	10.4	8.3	7.6	7.7	2.1	2.8	2.7
6	10.4	8.1	7.2	7.3	2.3	3.2	3.1
8	10.4	7.9	7.0	7.0	2.5	3.4	3.4
10	10.4	7.7	6.7	6.8	2.6	3.7	3.6
12	10.4	7.6	6.5	6.6	2.8	3.9	3.8
14	10.4	7.5	6.4	6.4	2.9	4	4



Average PA Efficiency with PAPR Reduction N=256



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BER with PA Nonlinearities for N=256 (a) QPSK, (b) 16-QAM



BER with PA Nonlinearities for N=256 (QPSK) (a) OPT-GDFT (b) EFF-GDFT



BER with PA Nonlinearities for N=256 (16-QAM) (a) OPT-GDFT (b) EFF-GDFT



Potential GDFT Applications

- PAPR Reduction
- PAPR-ISI-ICI-Spectrum and Power Efficiency Trade-offs
- OFDM Variations SC-OFDM/MC-OFDM/DS-CDMA-OFDM and LTE Types
- Scrambling & Cryptography
- Basis Hopping

VI. BER Performance of DFT and GDFT CDMA for Rayleigh channel and 2 users



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In DS/CDMA, each user assigned a spreading code. The transmitted signal uses the entire frequency band. Spreading is performed in time-domain.

$$y(t) = \sum_{k=1}^{K} A_{k} b_{k}(t) s_{k}(t) + n(t)$$

$$y(t) = \sum_{k=1}^{K} A_{k} b_{k}(t - \tau_{k}) s_{k}(t - \tau_{k}) + n(t) \qquad t \in [0, T]$$

$$Z_{i} = b_{i}(0) + \sum_{k\neq 1}^{K} b_{k}(-1) \rho_{ki}(\tau_{k}) + b_{k}(0) \hat{\rho}_{ki}(\tau_{k})$$

$$\rho_{ki}(\tau) = \int_{0}^{\tau} s_{i}(t) s_{k}(t - \tau) dt$$

$$\hat{\rho}_{ki}(\tau) = \int_{\tau}^{T} s_{k}(t - \tau) s_{i}(t) dt$$

$$d_{k,i}(l) = \begin{cases} \frac{1}{N} \sum_{n=0}^{N-1-l} a_k(n) a_i(n+l), \ 0 < l \le N-1 \\ \frac{1}{N} \sum_{n=0}^{N-1+l} a_k(n-l) a_i(n), \ 1-N < l \le 0 \\ 0 , |l| \ge N \end{cases}$$

$$\rho_{k,i}(\tau) = d_{k,i}(l-N) + [d_{k,i}(l+1-N) - d_{k,i}(l-N)](\tau-l)$$
$$\hat{\rho}_{k,i}(\tau) = d_{k,i}(l) + [d_{k,i}(l+1) - d_{k,i}(l)](\tau-l) \qquad 0 \le l < \tau < l+l \le N$$

$$\theta_{k,i}(l) = \sum_{n=0}^{N-1} a_k(n) a_i(n+l)$$

$$\theta_{k,i}(l) = d_{k,i}(l) + d_{k,i}(l-N)$$

$$\hat{\theta}_{k,i}(l) = d_{k,i}(l) - d_{k,i}(l-N) \quad \text{for } 0 \le l \le N$$

NJIT

Multiple user interference in DS/CDMA system depends on the even and odd correlations between the user spreading codes.

If
$$b_k(-1) \neq b_k(0)$$

 $Z_i = b_i(0) + \sum_{k \neq i}^{K} b_k(0)\hat{\theta}_{k,i}(l_k) + [\hat{\theta}_{k,i}(l_k+1) - \hat{\theta}_{k,i}(l_k)](\tau_k - l_k)$
If $b_k(-1) = b_k(0)$
 $Z_i = b_i(0) + \sum_{k \neq i}^{K} b_k(0)\theta_{k,i}(l_k) + [\theta_{k,i}(l_k+1) - \theta_{k,i}(l_k)](\tau_k - l_k)$



With the advances in Digital Signal Processing and success of Multicarrier Modulation in early broadcast applications motivated researchers to investigate the suitability of multicarrier modulation in mobile wireless communications.

Multicarrier CDMA and MC-DS/CDMA are introduced
emerging spreading spectrum techniques with multicarrier
modulation to serve multiple users even on frequency selective
channels.

In MC-CDMA, the spreading is performed in frequency domain whereas in MC-DS/CDMA, the spreading is in timedomain. In both methods, all users shared the same available bandwidth simultaneously

MC-DS-CDMA offers better performance on Rayleigh channel due to the use of orthogonal sub-carriers each having equally spaced bandwidth in the frequency. MC-DS-CDMA may be considered as N-channel DS-CDMA system.







MC-CDMA Transmitter





4G Technologies aim to increase cell capacity, cell radius, scalability of bandwidth and data rates deploying a completely new technology or emerging existing 3G networks.

Two parallel standardization efforts are IEEE 802.16 (Wimax) and 3GPP LTE.

Both standards use OFDMA for downlink. Wimax uses OFDM for uplink whereas 3GPP LTE suggest using SC-FDMA for uplink. SC-FDMA is preferred due to its advantage of low PAPR over OFDM systems. Both employ MIMO with multiple antennas for UL and DL.



for Long-Term Evolution. John Wiley & Sons, Nov. 2008.

- Roaming framework for Mobile Wimax is completely new whereas 3GPP LTE is based on existing GSM/UMTS communications systems.
- Legacies for Wimax is IEEE 802.16a through IEEE 802.16d, the legacies for 3GPP LTE GSM, GPRS, UMTS, EGPRS, HSPA.



- Nokia Siemens Network, Motorola, Ericson, Freescale Semiconductor and NTT DoCoMo demonstrated successful implementations of LTE networks.
- AT&T, T-Mobile, Verizon Wireless, Vodafone, France Telecom are among those companies which announced their intension to upgrade their current networks to LTE.



SPECS OF WIMAX SYSTEM SIMULATED

FFT size	1024
Number of data subcarriers	720
Number of pilot subcarriers	120
Number of null-Guard	184
subcarriers	
Channel bandwidth	10MHz







- GDFT is a continuation of early work on Nonlinear Phase Walsh-like transforms (binary phase grid)
- We introduced Generalized DFT (GDFT) framework with Nonlinear Phase and Efficient Design Methods (any phase grid) for design of constant modulus sets
- Marked departure from block-circulant correlation and eigen-structures of DFT

- Methodically interconnects constellation, DFT and OFDM frame of interest
- Graceful departure from OFDM to CDMA or any TF-MA
- Basis hopping for better fit to channel (loading) and/or code level security (built-in scrambler) is inherent



- Design flexibilities for possible improvements (e.g. BER, PAPR) over DFT with efficient add on to FFT
- Next Generation Multicarrier Communications Systems (SW based) might benefit from GDFT family
- Currently looking into radar applications
 including MIMO Radar employing GDFT
 concepts



- Extensions to filters banks being studied for complex MUX
- Hilbert pair interpretation of a GDFT subset being formalized



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