# Duality and pricing

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Section 6.1.1-6.1.2, 6.2.1 in Convex Optimization of Power Systems.

### 1 The power industry

- Initially: competitive, no regulation (Edison vs. Westinghouse, later others)
- 1935: Samuel Insull, PUHCA
- 1935-1978: Vertically integrated, gov't regulated monopoly
- 1978: PURPA utilities must buy from IPPs
- 1978 present: deregulation

### Current format:

- Generators (OPG) sell power to:
  - large loads (steel mills) and utilities (HydroOne) via bilateral contracts
  - system operators (IESO) in spot markets
- System operators sell power to large loads and utilities in spot markets
- Utilities sell power to smaller loads.

#### Today: Spot markets

- Add liquidity, increase competition, lower barrier to entry for new generators
- Typically 3 markets:
  - Day-ahead: bulk quantities

- Hour-ahead: moderate adjustments
- Real-time (5 min. ahead): balancing
- Basic operation:
  - Generators sell power to ISO, ISO sells to loads
  - Transactions occur at nodal prices AKA LMPs.
  - Prices are dual variables of OPF.

# 2 Duality

Primal problem

$$P = \min_{x} \quad f(x)$$
  
subject to 
$$g_{i}(x) \leq 0$$
$$h_{i}(x) = 0$$

Lagrangian:

$$L(x, \alpha, \beta) = f(x) + \sum_{i} \alpha_{i} g_{i}(x) + \sum_{i} \beta_{i} h_{i}(x), \quad \alpha_{i} \ge 0$$

Dual function

$$\mathcal{L}(\alpha,\beta) = \min_{x} L(x,\alpha,\beta).$$

Always concave because minimum over family of affine functions (draw). If x is primal feasible  $(g_i(x) \leq 0, h_i(x) = 0)$ , then since  $\alpha \geq 0$ ,

$$L(x, \alpha, \beta) \le f(x).$$

True for all x, so

 $\mathcal{L}(\alpha,\beta) \le P,$ 

Weak duality. Want largest lower bound ... dual problem:

$$D = \max_{\alpha,\beta} \quad \mathcal{L}(\alpha,\beta)$$
  
subject to  $\alpha \ge 0$ 

Concave maximization regardless of primal. Strong duality:

$$D = P$$

Holds when primal is convex (+ constraint qualification).

#### 2.1 Example: LP duality

Standard form LP:

$$\begin{array}{ll} \underset{x}{\text{minimize}} & l^T x\\ \text{subject to} & Ax = b\\ & x \ge 0 \end{array}$$

Dual function:

$$\mathcal{L}(\alpha,\beta) = \min_{x} l^{T}x + \beta^{T}(b - Ax) - \alpha^{T}x$$
$$= \beta^{T}b + \min_{x} (l - A^{T}\beta - \alpha)^{T}x$$

Equals  $-\infty$  unless if  $A^T\beta + \alpha = l$  ... a constraint in the maximization. Standard form LP dual:

$$\begin{array}{ll} \underset{\alpha,\beta}{\text{maximize}} & \beta^T b\\ \text{subject to} & A^T \beta + \alpha = l\\ & \alpha \geq 0 \end{array}$$

Note that the dual variable  $\alpha$  can be eliminated by consolidating the two constraints into  $A^T \beta \leq l$ . Strong duality guaranteed by convexity.

#### 2.2 Complementary slackness

Strong duality, D = P, implies at optimal solution  $(x^*, \alpha^*, \beta^*)$ ,

$$f(x^*) = \mathcal{L}(\alpha^*, \beta^*)$$
  
=  $\min_x L(x, \alpha^*, \beta^*)$   
 $\leq L(x^*, \alpha^*, \beta^*)$   
=  $f(x^*) + \sum_i \alpha_i^* g_i(x^*) + \sum_i \beta_i^* h_i(x^*)$   
 $\leq f(x^*)$ 

because the latter terms must be negative due to feasibility. Therefore

$$\alpha_i^* g_i \left( x^* \right) = 0.$$

Central in transmission economics.

# 3 Pricing the linearized power flow

Recall linearized power flow:

$$\begin{array}{ll} \underset{p,\theta}{\text{minimize}} & \sum_{i} f_{i}(p_{i}) \\ \text{subject to} \\ \lambda_{i}: & p_{i} = \sum_{j} b_{ij}(\theta_{i} - \theta_{j}) \\ \chi_{ij} \geq 0: & b_{ij}(\theta_{i} - \theta_{j}) \leq \overline{s}_{ij} \end{array}$$

Lagrangian:

$$L(p,\theta,\lambda,\chi) = \sum_{i} f_{i}(p_{i}) + \lambda_{i} \left( -p_{i} + \sum_{j} b_{ij}(\theta_{i} - \theta_{j}) \right) + \sum_{j} \chi_{ij} \left( b_{ij}(\theta_{i} - \theta_{j}) - \overline{s}_{ij} \right)$$

Stationarity conditions:

$$\frac{\partial f_i(p_i)}{\partial p_i} - \lambda_i = 0 \quad \forall i$$
$$\sum_{j:j\sim i} b_{ij}(\lambda_i - \lambda_j + \chi_{ij} - \chi_{ji}) = 0 \quad \forall i$$

 $\lambda_i$ : the price at node *i*. Agent *i* solves:

$$\underset{p_i}{\text{minimize}} \quad f_i(p_i) - \lambda_i p_i$$

- Strong duality means the solution to this matches centralized if  $\lambda_i$  is optimal multiplier.
- Nodal or locational marginal pricing
- An instance of the Fund. Theorems of Welfare Econ. also used to price comm. channels, Kelly mechanism.

$$\sum_{i} \lambda_{i} p_{i} = \sum_{i} \lambda_{i} \sum_{j:j \sim i} b_{ij} (\theta_{i} - \theta_{j})$$

$$= \sum_{i} \theta_{i} \sum_{j:j \sim i} b_{ij} (\lambda_{i} - \lambda_{j})$$

$$= \sum_{i} \theta_{i} \sum_{j:j \sim i} b_{ij} (\chi_{ji} - \chi_{ij})$$

$$= -\sum_{i} \sum_{j:j \sim i} \chi_{ij} b_{ij} (\theta_{i} - \theta_{j})$$

$$= -\sum_{i} \sum_{j:j \sim i} \chi_{ij} \overline{s}_{ij} \quad \text{by comp. slackness}$$

Therefore

$$\sum_{i} \lambda_i p_i + \sum_{i \sim j} \chi_{ij} \overline{s}_{ij} = 0.$$

Intuition:

- 1st term: SO's budget balance
- 2nd term: always positive
- $\bullet$   $\Rightarrow$  SO always makes money

#### 3.1 Example: 2 bus network

$$\begin{array}{ll} \underset{p,\theta}{\text{minimize}} & f_1(p_1) + f_2(p_2) \\ \text{subject to} \\ \lambda_1 : & p_1 = b(\theta_1 - \theta_2) \\ \lambda_2 : & p_2 = b(\theta_2 - \theta_1) \\ \chi_{12} \ge 0 : & b_{12}(\theta_1 - \theta_2) \le \overline{s}_{12} \\ \chi_{21} \ge 0 : & b_{12}(\theta_2 - \theta_1) \le \overline{s}_{12} \end{array}$$

Stationarity:

$$\frac{\partial f_i(p_i)}{\partial p_i} - \lambda_i = 0$$
  
$$b_{12}(\lambda_1 - \lambda_2 + \chi_{12} - \chi_{21}) = 0$$

Suppose line is uncongested (define). Then comp. slack. implies

- $\chi_{12} = \chi_{21} = 0$
- which implies  $\lambda_1 = \lambda_2$
- SO budget is  $\lambda_1 p_1 + \lambda_2 p_2 = \lambda_1 (p_1 p_1) = 0$ , balanced.

Now suppose congested (define). Then comp. slack. implies

•  $\chi_{12} > 0$  (if power flow from 1 to 2),

$$\lambda_1 - \lambda_2 = -\chi_{12} < 0$$

- which implies  $\lambda_2 > \lambda_1$
- SO budget is  $\lambda_1 p_1 + \lambda_2 p_2 = p_1(\lambda_1 \lambda_2) < 0$ , extra money.

### 3.2 More discussion

- In radial networks,  $\lambda_1 \lambda_2 = \chi_{21} \chi_{12}$ , nodal price differences equal sum of shadow prices along line
- Not true in general networks!!!
- $\chi$  helps define financial transmission rights (in two weeks)
- [1] gives good discussion

### Convex relaxations

- Also have strong duality, can define real + reactive power prices
- Reactive power pricing is dubious though
- Congestion is present, but above relationships don't hold

### Limitations

- Strong duality is required. Exact power flow not convex.
- Unit commitment many costs there, not captured. Lead to generators being underpaid.
- Uplift or make-whole payments are made to generators in this case.

# References

 Felix Wu, Pravin Varaiya, Pablo Spiller, and Shmuel Oren. Folk theorems on transmission access: Proofs and counterexamples. *Journal of Regulatory Economics*, 10:5–23, 1996.