Name: .................................................................

ID Number: .................. Exam Type: ...B... Exam Number: ...........

Grade: 1: ... 2: ... 3: ... 4: ... 5: ... Total: .......

Solve all the problems in the space provided

There are 3 (THREE) pages this page included

Read and Sign the statement below at the end of the exam

Unsigned exams will be marked with grade 0 (zero).

STATEMENT

On my honor, I pledge that I have not violated the provision of the NJIT Student Honor Code.

Sign below at the end of the exam

Signature ..........................................................

List of useful formulae

\[ \sum_{i=1}^{n} i = \frac{n(n+1)}{2}, \quad \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4} \]

\[ n! \approx \left( \frac{n}{e} \right)^n, \quad a^{\log_b n} = n^{\log_b a}, \]

For \( x \neq 1 \), we have that

\[ \sum_{i=0}^{n} x^i = \frac{x^{n+1} - 1}{x - 1}, \quad \sum_{i=0}^{n-1} ix^i = \frac{(n-1)x^{n+1} - nx^n + x}{(1-x)^2}, \quad \sum_{i=0}^{\infty} i \cdot (1/2)^i = 2 \]

NOTE: Give an algorithm means, present the algorithm, show that it works as claimed, and analyze its running time by providing an asymptotically tight bound.

Master Method. \( T(n) = aT(n/b) + f(n) \), such that \( a > 0, b > 1, f(n) > 0 \).

M1 If \( f(n) = O(n^{\log_b a - \epsilon}) \) for some constant \( \epsilon > 0 \), then \( T(n) = \Theta(n^{\log_b a}) \).

M2 If \( f(n) = \Theta(n^{\log_b a} \log^k n) \), then \( T(n) = \Theta(n^{\log_b a} \log^{k+1} n) \), where \( k \geq 0 \) is a non-negative constant.

M3 If \( f(n) = \Omega(n^{\log_b a + \epsilon}) \) for some constant \( \epsilon > 0 \), and if \( af(n/b) \leq cf(n) \) for some constant \( c < 1 \) and for large \( n \), then \( T(n) = \Theta(f(n)) \).

* : Read Syllabus for a possible quiz in April and how this could affect Exam 1.
Problem 1. (35 points)
Give a TRUE or FALSE for each one of the statements below. Any answer other than a full TRUE or FALSE will be considered wrong. All algorithms are the book/notes implementations.

(1) MergeSort is stable.

(2) Odd-even transposition sort sorts in-place.

(3) On the reverse-sorted sequence ⟨n, n − 1, . . . , 1⟩, SelectionSort is asymptotically faster than InsertionSort.

(4) On the reverse-sorted sequence ⟨n, n − 1, . . . , 1⟩, InsertionSort has running time that is Ω(n lg n).

(5) The asymptotic solution of \( T(n) = T(n/2) + n \) is \( T(n) = \Theta(n \lg n) \).

(6) The asymptotic solution of \( T(4n) = T(2n) + n \) is \( T(n) = \Theta(n \lg n) \).

(7) The asymptotic solution of \( T(n) = T(n/2) + n^2 \) is \( T(n) = \Theta(n^2) \)

Problem 2. (50 points)
Rank the following functions by order of growth; that is, find an arrangement \( g_1, g_2, g_3, \ldots, g_6, g_7 \) of the functions satisfying \( g_1 = \Omega(g_2) \), \( g_2 = \Omega(g_3) \), \( g_3 = \Omega(g_4) \), \( g_4 = \Omega(g_5) \), \( g_5 = \Omega(g_6) \), \( g_6 = \Omega(g_7) \). Partition your list in equivalence classes such that \( f(n) \) and \( h(n) \) are in the same class if and only if \( f(n) = \Theta(h(n)) \). For example for functions \( \lg n, n, n^2 \), and \( 2^{\lg n} \) you could write: \( n^2, \{n, 2^{\lg n}\}, \lg n \).

\[ \text{lg } n, \; n^2, \; (n - 2)!, \; 2^{2\lg n}, \; \sqrt{n}, \; \lg (n!), \; n \lg n. \]

Problem 3. (40 points)
Show that \( n^3 - n^2 - 2 = \Theta(n^3) \) by providing the \( n_0, c_1, c_2 \) of the definition of \( \Theta \).
Problem 4. (41 POINTS)
No asymptotics! You are given the following sorted sequences of a total (combined) 4n keys: A with 3n sorted keys, B with n/4 sorted keys, C with n/4 sorted keys and D with n/2 keys. How many comparisons does it take to merge (optimally) those sequences into a sorted sequence of 4n keys? Explain and BE precise; no asymptotics. (You may assume n is a multiple of 4.)

Problem 5. (34 POINTS)
Fill in the table below by providing the running time for every one of the indicated operations and indicated data structure (DaSt) of size $n$ containing distinct keys. Fill entries by writing $n$ to denote for example running time of $O(n)$ or $\Theta(n)$, 1 to denote $\Theta(1)$, and so on. We fill in the entry for Succ(z) into an SDL. Since it is $\Theta(1)$ we write 1.

(a) Ins, Sea is for Insertion and Search of a key $x$. (b) FindMin or FindMax determine and return the value of the key with the smallest or largest value among the keys of the DaSt; the key remains in it. (c) ExMin or Exmax perform FindMin or FindMax respectively and in addition remove the corresponding key from DaSt. (d) Del deletes the record referenced/pointed/indexed by $z$ i.e. the location that will be deleted is known. (e) Succ is the successor of $z$, i.e. the smallest among all the elements greater than or equal to the value pointed by reference/pointer/index $z$. For removal operations the data structures must be maintained in likewise form after the operation (i.e. there are no gaps or locations with no keys).

SA, RA Sorted Array (SA) and Reverse-Sorted Array with elements ordered appropriately in increasing or decreasing order.

SLL, RLL Sorted SinglyLinked List (SLL) and Reverse-Sorted SinglyLinked List (RLL) and the first value (smallest or largest) key pointed by a head pointer; no tail pointer.

SDL, RDL Sorted DoublyLinked List (SDL) and Reverse-Sort DoublyLinkedList defined similarly as before; no tail pointer again.

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<th>SLL</th>
<th>SDL</th>
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