

# Energy efficient rate control mechanism for multimedia delivery over wireless Internet



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The  
Economist

A very Japanese revolution

PAGES 17 AND 22-24

After Katrina

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Tackling Latin American poverty

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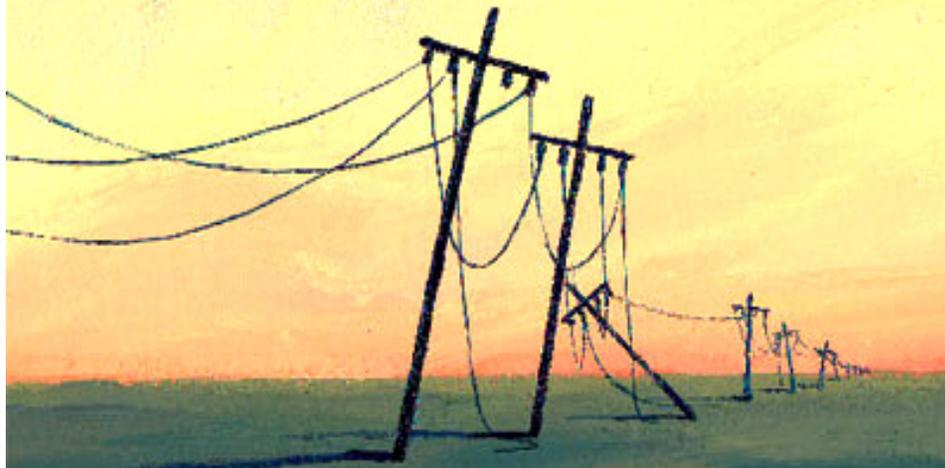
Steve Jobs, resurrection man

PAGE 48

SEPTEMBER 27TH - 23RD 2005

www.economist.com

# How the internet killed the phone business



Argentina	\$6.00	Canada	US\$ 9.00	Japan	¥120.00	Spain	€6.75
Australia	A\$10.00	China	¥45.00	India	₹100.00	Sweden	SEK 60.00
Brazil	R\$12.00	France	€7.50	Italy	₹100.00	Switzerland	CHF 6.00
India	₹100.00	Germany	€7.50	South Korea	₩100.00	UK	£5.00
Indonesia	₹100.00	USA	US\$ 9.00	Taiwan	NT\$ 300.00	USA	US\$ 9.00

THE Internet touched our lives in terms of

2

- ✓ How we communicate
- ✓ How we promote our products
- ✓ How we teach our children
- ✓ How we invest our time

Growth.....50 million

Radio..... 40 years

TV .....15 years

Internet .....5 years

[*Economist, September 2005*]



# Outline

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- 
- Wireless Internet (WI) and Multimedia Delivery
  - Major Challenges
  - Need for Loss Differentiation
  - RCM-BR Protocol
  - Simulation study
  - Conclusion



# Wireless Internet (WI)

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A service granting access to the World Wide Web or Internet via Wireless Networks. It consists of

- ✓ **Fixed internet hosts.**
- ✓ **Mobile nodes.**
- ✓ **Connectors like access points in between.**

They host information resources as well as application software for providing network services.

They have wireless access card to access Internet. Ex. Wi-Fi card

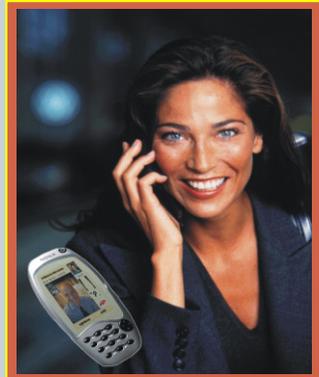
To provide wireless Access to the devices





# Multimedia over IP Networks

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Wireless Browsing



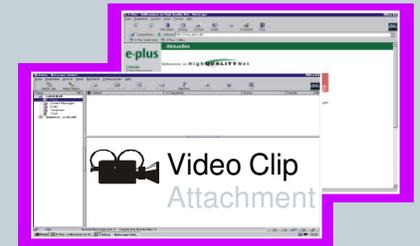
VoIP



Video Conference



Internet



E-mail



Music Streaming



Information Search



Movies Streaming



Finance, Brokerage



Digital Photos



# MM Networking Applications

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## Classes of MM applications:

- 1) Streaming stored audio and video
- 2) Streaming live audio and video
- 3) Real-time interactive audio and video

## Fundamental characteristics:

- Typically **delay sensitive**
  - End-to-end delay
  - Jitter – Delay variation
- But **loss tolerant**: infrequent losses cause minor glitches
- Normal data can tolerate delay tolerant, but not loss



# Why Is Real-Time Transport Hard?

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Internet is a best-effort network . . .

*Congestion*

Insufficient rate to communicate

*Packet loss*

Impairs perceptual quality

*Delay*

Impairs interactivity of services;

Telephony: one way delay < 150 ms

[ITU-T Rec. G.114]

*Delay jitter*

Obstructs continuous media playout



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# Challenges in Multimedia Delivery over Wireless Internet

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- Different QoS requirements for different types of media.
- High packet loss rate and bit error rate.
- Bandwidth limitation and fluctuation.
- Low performance in traditional transport protocols.
- Heterogeneity among users and networks.
- Limited battery life.

Ref: Qian Zhang, Fan Yang, Wenwu Zhu, "Cross-layer Qos Support for Multimedia delivery over Wireless Internet", EURASIP Journal on Applied signal processing 2005, pp207-219.



# Audio & Video Quality Requirements

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Medium	Application	Degree of symmetry	Typical data rates	Key performance parameters and target values			
				One-way Delay	Delay variation	Information loss	Other
Audio	Conversational voice	Two-way	4-13 kb/s	< 150 ms preferred* < 400 ms limit*	< 1 ms	< 3% packet loss ratio (PLR)	
Audio	Voice messaging	Primarily one-way	4-13 kb/s	< 1 s for playback < 2 s for record	< 1 ms	< 3% PLR	
Audio	High quality streaming audio	Primarily one-way	32-128 kb/s	< 10 s	< 1 ms	< 1% PLR	
Video	Videophone	Two-way	32-384 kb/s	< 150 ms preferred < 400 ms limit		< 1% PLR	Lip-synch < 80 ms
Video	One-way	One-way	32-384 kb/s	< 10 s		< 1% PLR	

\* Assumes adequate echo control



# Outline

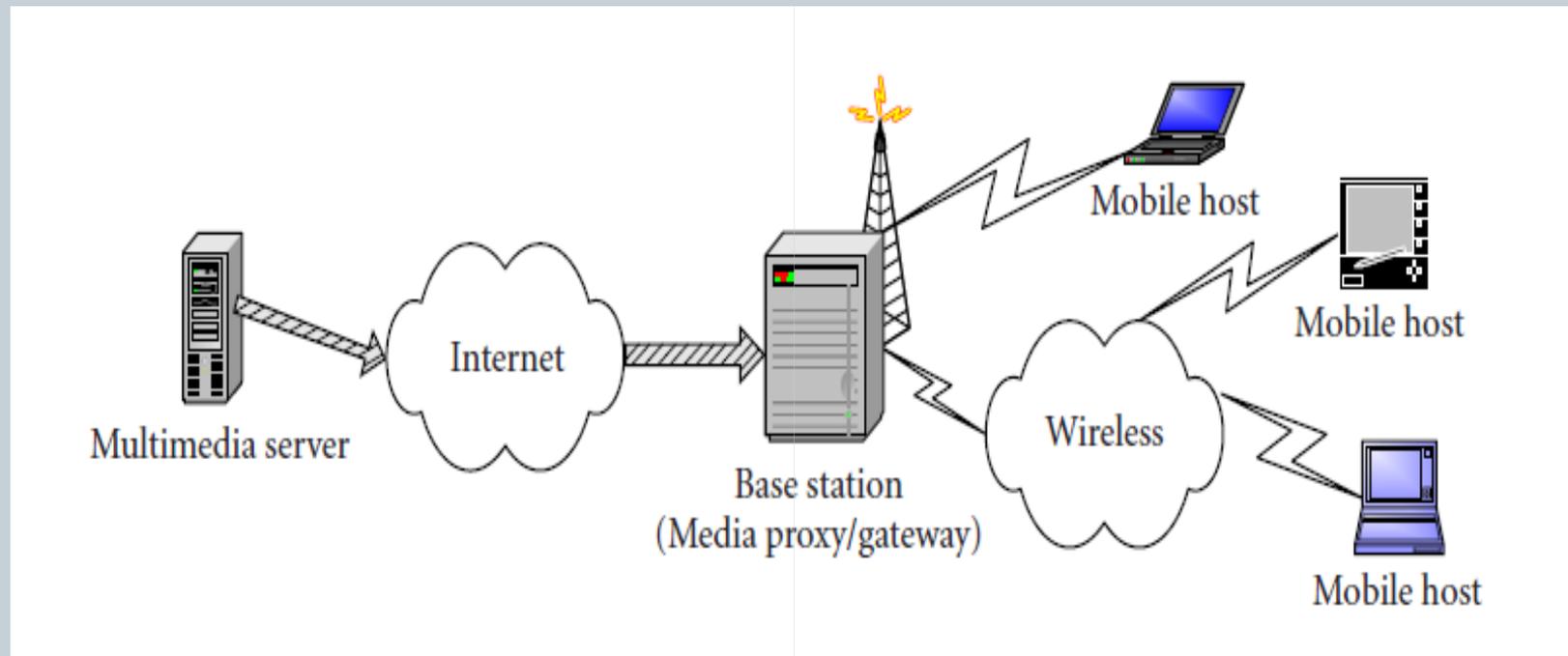
11

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# Wireless Multimedia Streaming Architecture

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Ref: Qian Zhang, Fan Yang, Wenwu Zhu, "Cross-layer Qos Support for Multimedia delivery over Wireless Internet", EURASIP Journal on Applied signal processing 2005, pp207-219.



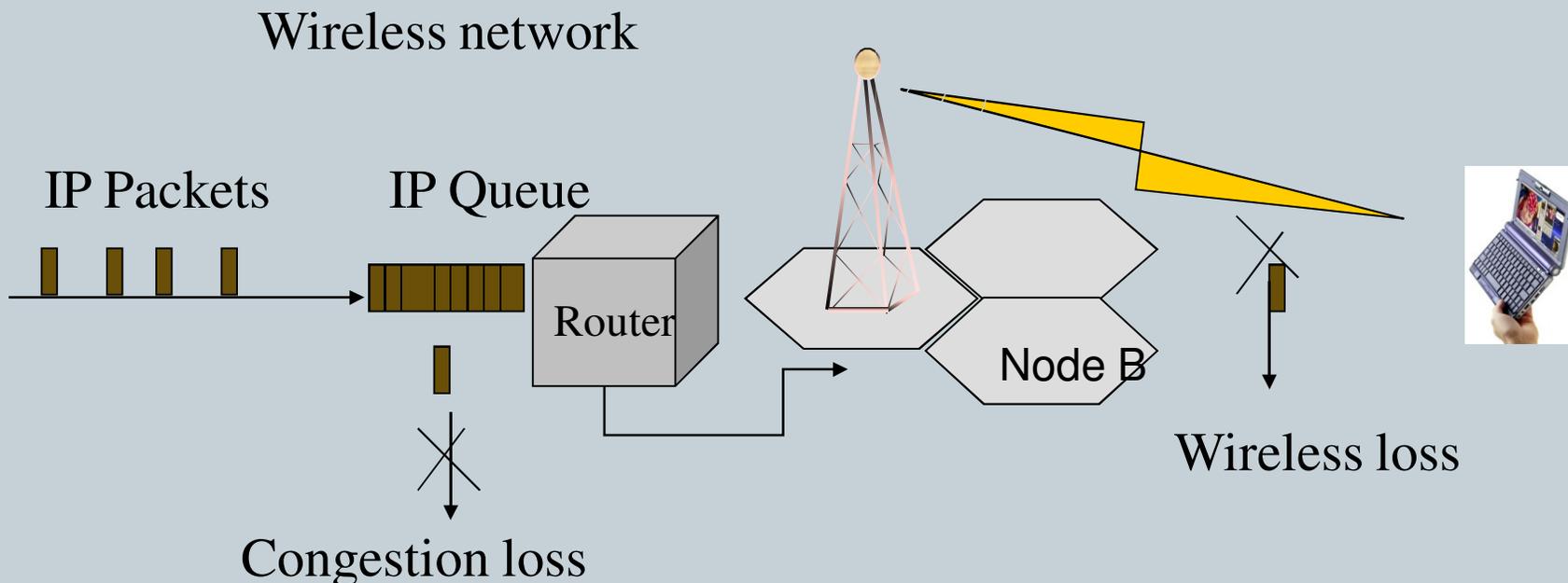
# Problems: Wireless Losses

13

- Problem: Two types of losses in wireless Networks

Packet drops due to congestion

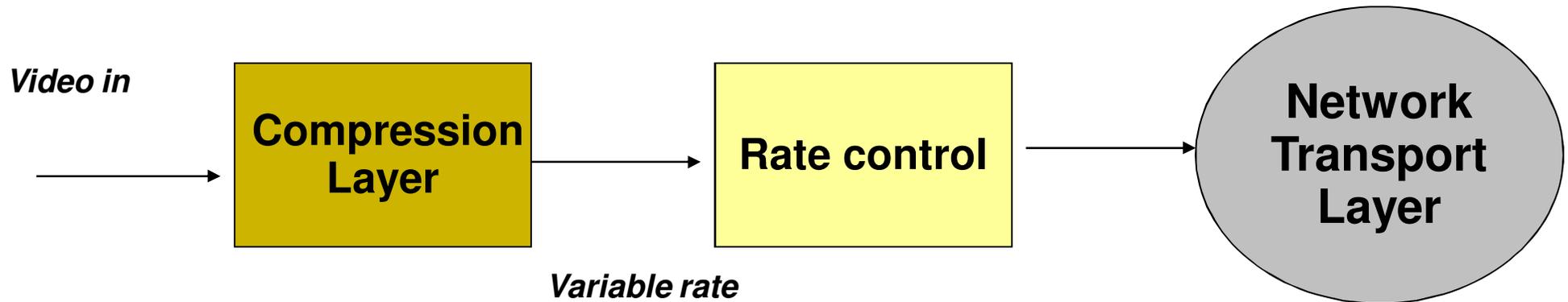
Packet drops due to bad channel conditions





# Need For Rate Control

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*“Multimedia traffic transport requires rate control deployment to protect the shared Internet from unfairness and further congestion collapse.”*

Ref: O. B. Akan and I. F. Akyildiz, “ARC: the analytical rate control scheme for real-time traffic in wireless networks,” *IEEE/ACM Transactions on Networking*, vol. 12, no. 4, pp. 634–644, 2004



# Desired TCP Behavior in Wireless Internet

15

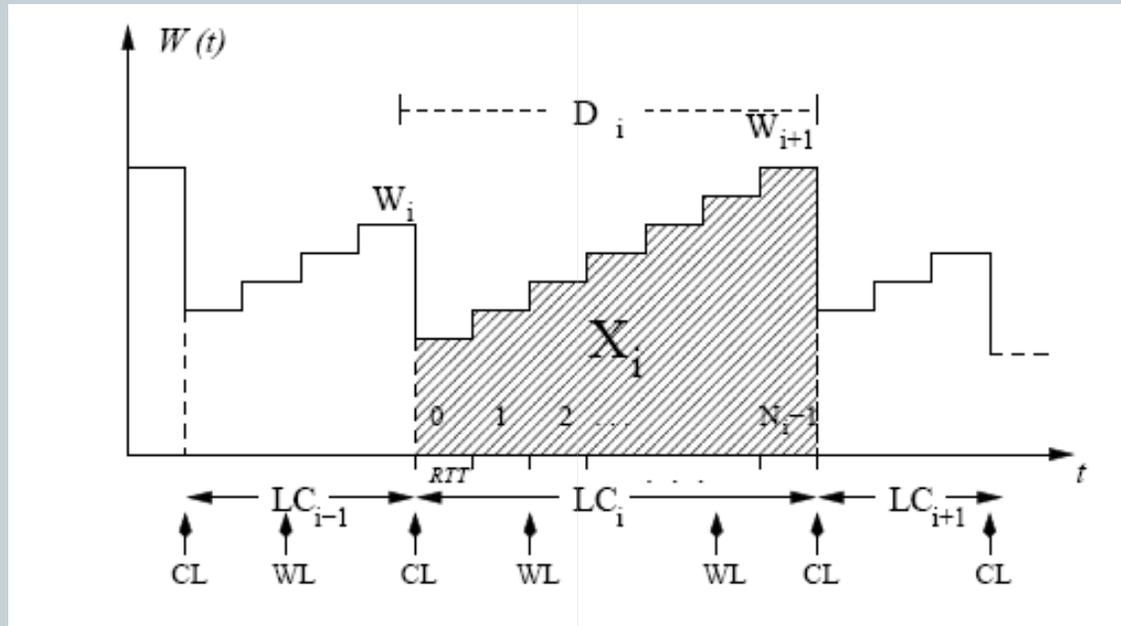


Fig:1 Steady state desired behavior of a TCP source in wireless link



# Equation Based Rate Control

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*“Equation based rate control has lower variation in throughput over time compared with TCP. That’s why it is more suitable for applications where smooth sending rate is required”.*

Ref: RFC 3448 bis



# Existing Schemes

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➤ TFRC [2000]

$$T = \frac{s}{RTT \sqrt{\frac{2p}{3}} + t_{RTO} \left( 3 \sqrt{\frac{3p}{8}} \right) p (1 + 32 p^2)}$$

➤ ARC [2004]

$$T = \frac{s}{4 \cdot RTT \left( 3 + \sqrt{25 + 24 \left( \frac{1 - \omega}{\pi - \omega} \right)} \right)}$$

➤ TFRC-JR[2007]

➤ Some Other Schemes:

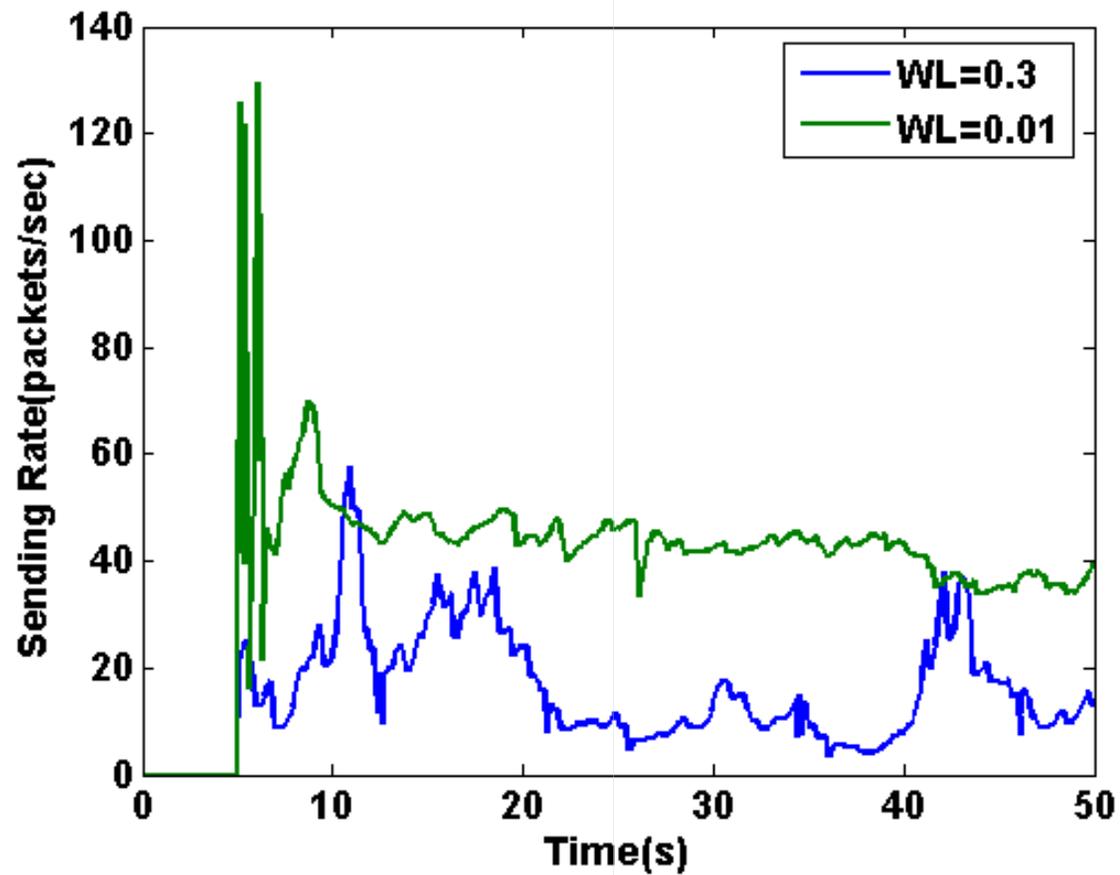
RCS[2001], WMSTFP[2003], WLED[2006], DCCP[2006],

Equation Based Rate control in Cellular Networks[2009];



# TFRC Rate Variation

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# Energy potential

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- Congestion losses
  - Burst congestion losses
  - Transient congestion losses
- Wireless losses
  - Transient/Random losses
  - Burst wireless losses

In case of burst wireless losses

$n_{\text{lost}} \propto \text{rate}$

$\text{energy wastage} \propto n_{\text{lost}}$



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# RCM-BR protocol overview

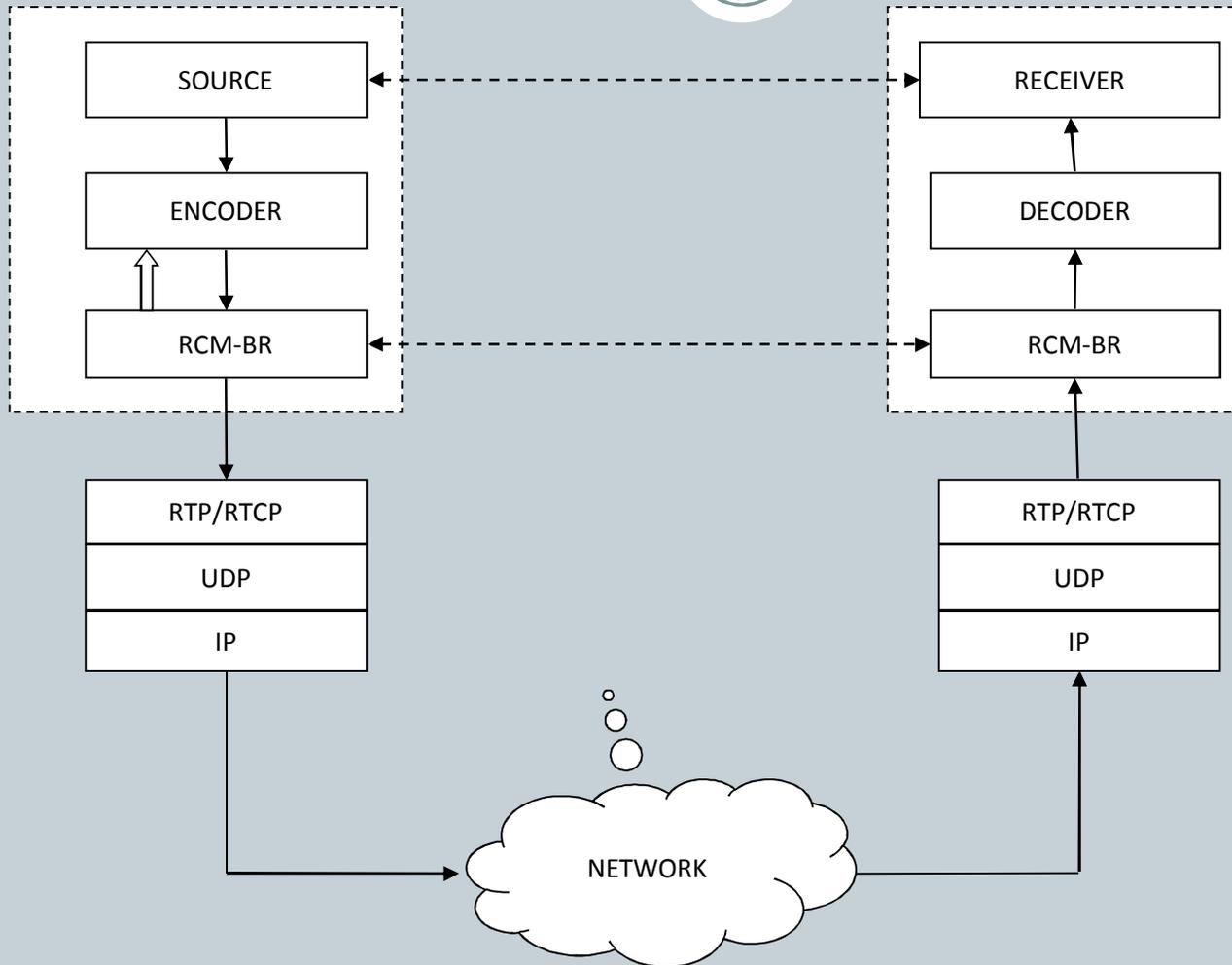
21

- Estimate RTT
- Differentiate wireless and congestion losses
- Detect burst wireless loss
- Estimate loss probability
- Estimate available B.W.
- Adjust the rate



# The Protocol Structure

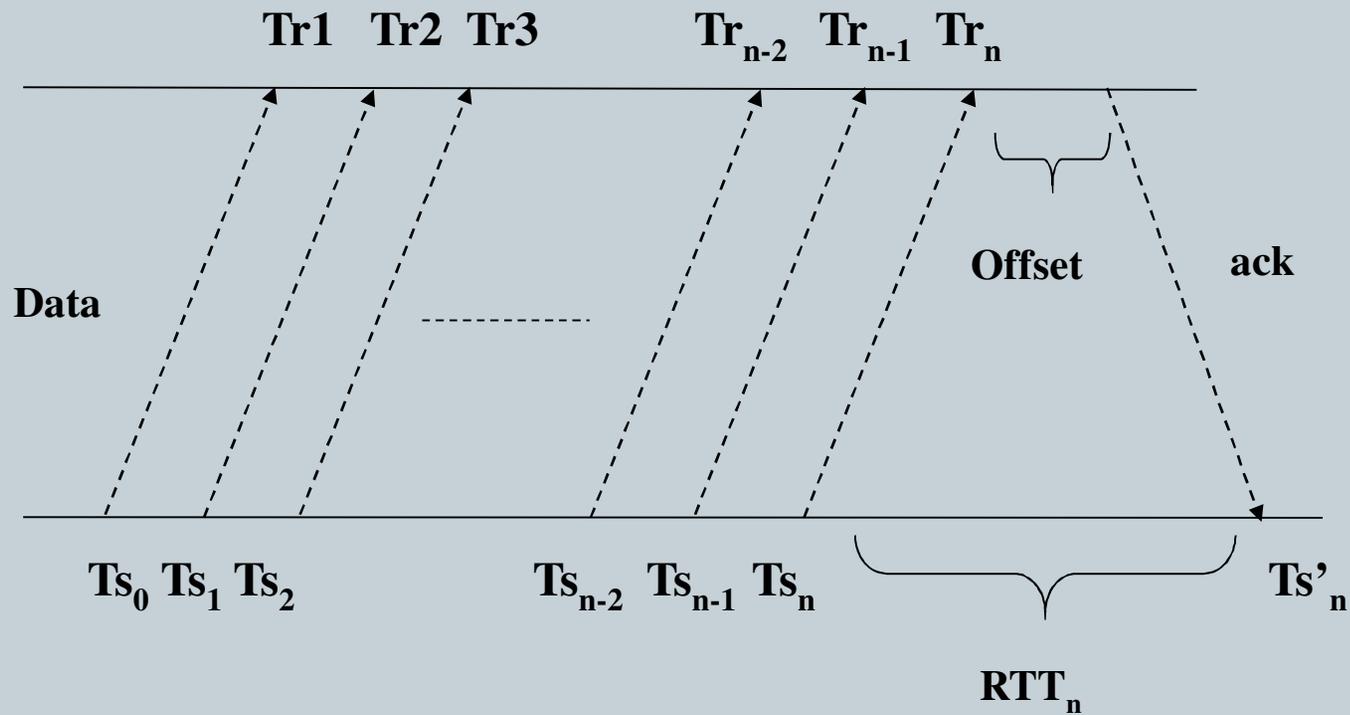
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# Estimating RTT

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# Adaptive Loss Differentiation

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- Infers the cause of loss: The threshold is a function of the minimum RTT, and current sample RTT so that it may automatically adapt itself to current congestion level.

$$T > \bar{T} + T_{dev} \cdot \left( 2 \cdot \left( \frac{T_p}{T} \right)^K - 1 \right)$$

$$\bar{T} = \frac{7}{8} \bar{T} + \frac{1}{8} T$$

$$T_{dev} = \frac{3}{4} T_{dev} + \frac{1}{4} |T - \bar{T}|$$



# Estimating loss probabilities

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- Wireless loss

$$A = \begin{bmatrix} 1-p & p \\ q & 1-q \end{bmatrix}$$

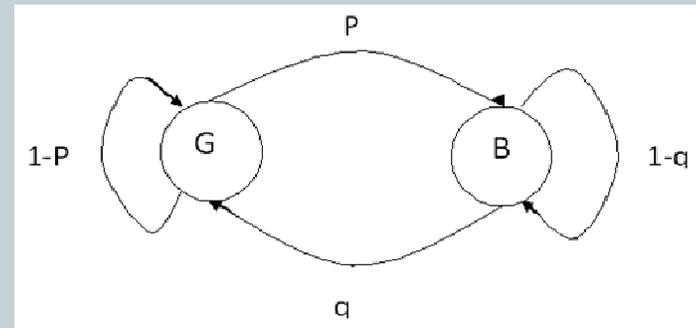
$$p = P(q_t = B | q_{t-1} = G);$$

$$q = P(q_t = G | q_{t-1} = B);$$

$$\pi_B = \frac{p}{p + q}$$

$$\pi_G = \frac{q}{p + q}$$

$$p_E = (1-P_G) \pi_G + (1-P_B) \pi_B;$$



- Congestion loss

$$p_c = (1-P_r) \pi_r + (1-P_l) \pi_l;$$

- Total loss

$$\Pi = 1 - (1 - p_E)(1 - p_c)$$



# Accuracy of Classification

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Type of algorithm	C/C	W/W	C/W	W/C
Spike	100%	0%	100%	0%
Zig-Zag	85%	60%	40%	15%
Adaptive LDA	82.5%	85.38%	14.62%	17.5%
SVM	88%	84%	16%	12%



# Detecting burst wireless loss

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burst-thrust function

$$g(k) = \begin{cases} 1, & k \leq \theta \\ \sqrt{k - \theta}, & k > \theta \end{cases}$$

Where,  $\theta = \max(4, \text{previous rate} * \text{rtt} / 2)$

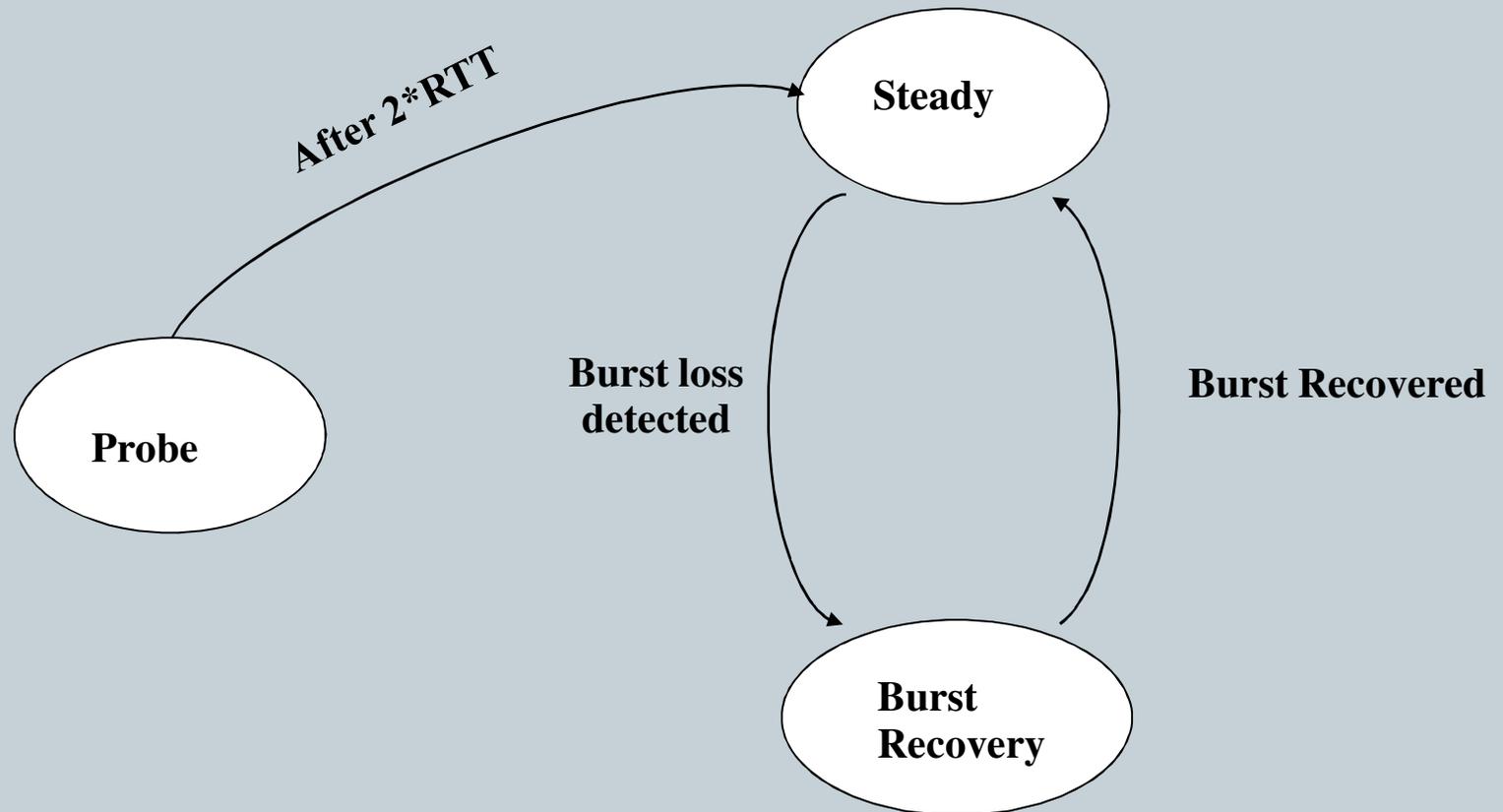
k is the no .of packets lost due to wireless loss

$g(k) > 1 \Rightarrow$  Burst wireless loss



# Finite state machine model of RCM-BR source

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# Algorithm

29

**Probe()**

While( $t \leq RTT$ )

    Send PROBE\_PACKET with  $S_{target}$

End;

Estimate initial rate

**end;**

**STEADY()**

    Calculate  $RTT$ ;

    Calculate  $g(wlost)$

    If ( $g(wlost) = 1$ )

        Calculate  $\pi, \omega$ ;

        Calculate  $T$ ;

$R_f = (now - lastchange) / RTT$

        If ( $T > rate\_$ )

$rate = rate + R_f * (1 - \pi) * T$ ;

        else

$rate = a * T + (1 - a) * rate$ ;

    else

$rate = (1 - w) * rate$

        burst recovery()

**end;**

**Burst recovery()**

    calculate  $g(k)$

    if ( $g(k) > 1$ )

$rate = (1 - w) * rate$ ;

    else

$rate = prev\_g(k) * rate$

        steady()

**end;**

Where  $R_f$  is the rate change factor that detects how many times the sending rate is changed with respect to  $RTT$ . 'now' represent the current time and 'last change' represents the last time when the rate was changed.



# Probe period

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- Initially sender sends probe packets with target rate till RTT.
  - Low priority dummy packets
- Receiver sends acks to received probe packets
- At  $2 * RTT$ , sender sets initial transmission rate

$$S_{init} = \frac{\max\{1, n_{ack}\}}{RTT}$$



# Steady period

31

calculate RTT

Calculate  $g(k)$

If Burst wireless loss is not detected

Calculate  $\Pi, \omega$

Estimate available B.W

$$T = \frac{s}{4.RTT} \left( 3 + \sqrt{25 + 24 \left( \frac{1-\omega}{\pi-\omega} \right)} \right)$$

$Rf = (now - lastchange) / RTT$

If  $(T > rate)$

$$rate = rate + Rf * (1 - \pi) * T;$$

else

$$rate = a * T + (1-a) * rate$$



# Burst recovery

32

calculate  $g(k)$

if ( $g(k) > 1$ )

$\text{rate} = \max(\text{minrate}, (1-w) * \text{rate});$

else

$\text{rate} = \text{prevg}(k) * \text{rate}$



# Outline

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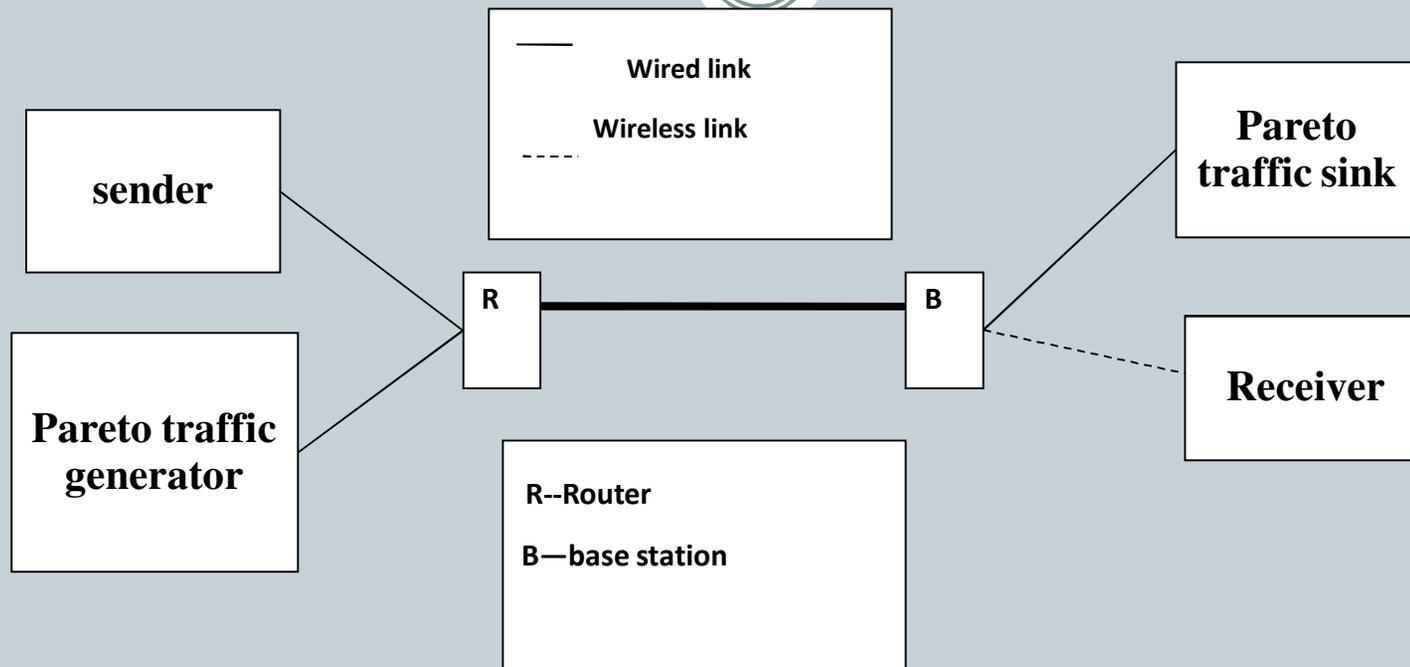
- Wireless Internet (WI) and Multimedia Delivery
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# Simulation model (NS-2)

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- Pareto traffic to generate congestion losses in wired bottleneck
- Three state markov model for wireless losses



# Performance Metrics

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- Throughput or Goodput.
- Delay or latency.
- Delay Variation - Jitter.
- Packet loss rate.

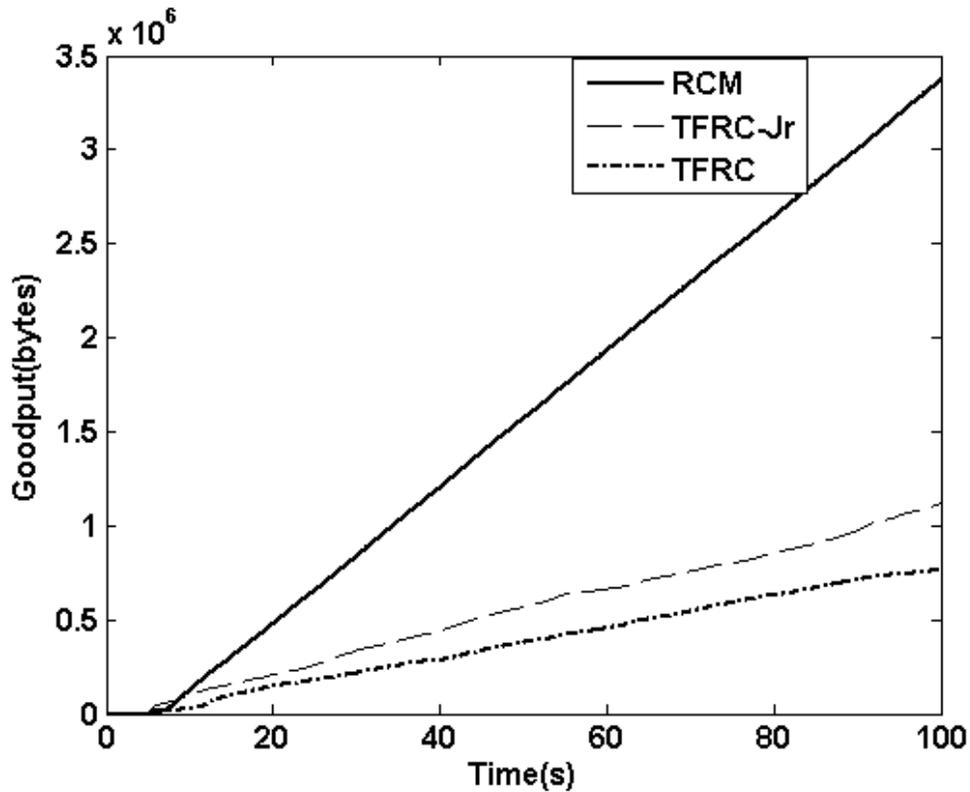
- Fairness Index, 
$$F.I = \frac{\left( \sum_{i=0}^n Throughput_i \right)^2}{n \left( \sum_{i=0}^n Throughput_i^2 \right)}$$

Ref: RFC5166,S.Floyd,March 2008.

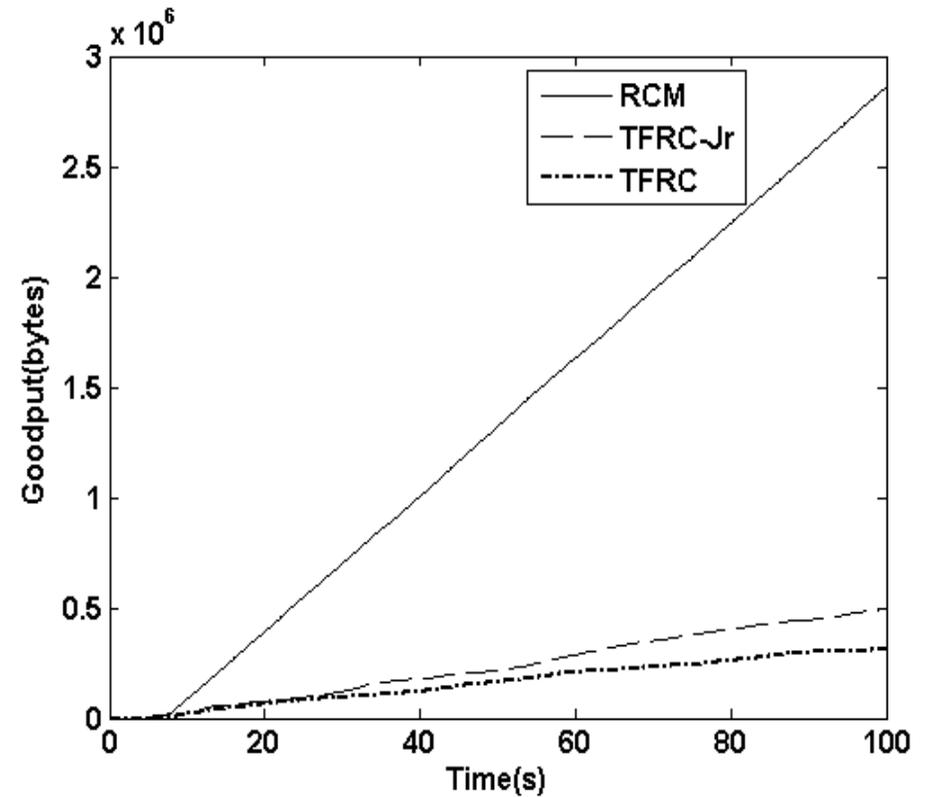


# Goodput Performance of RCM

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GoodPut Vs Time for WP=0.1.

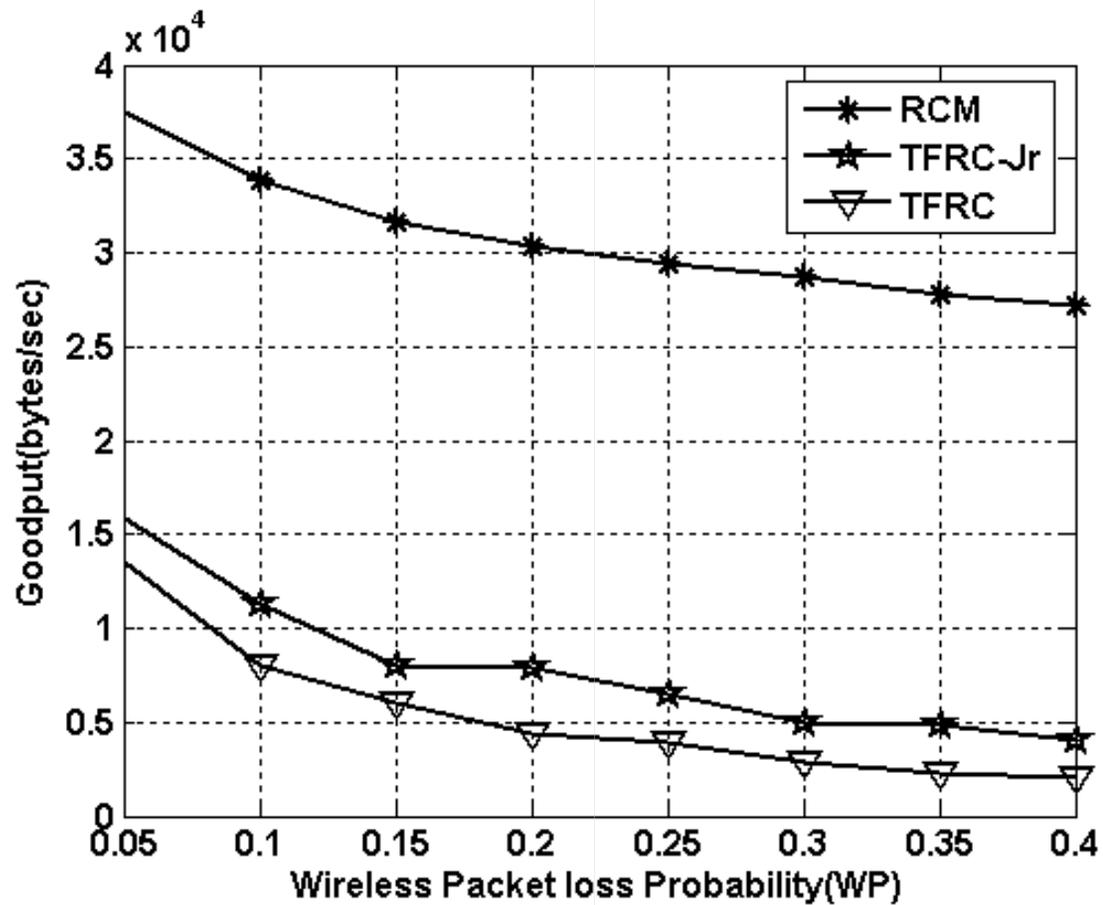


GoodPut Vs Time for WP=0.3.



# Goodput Performance of RCM

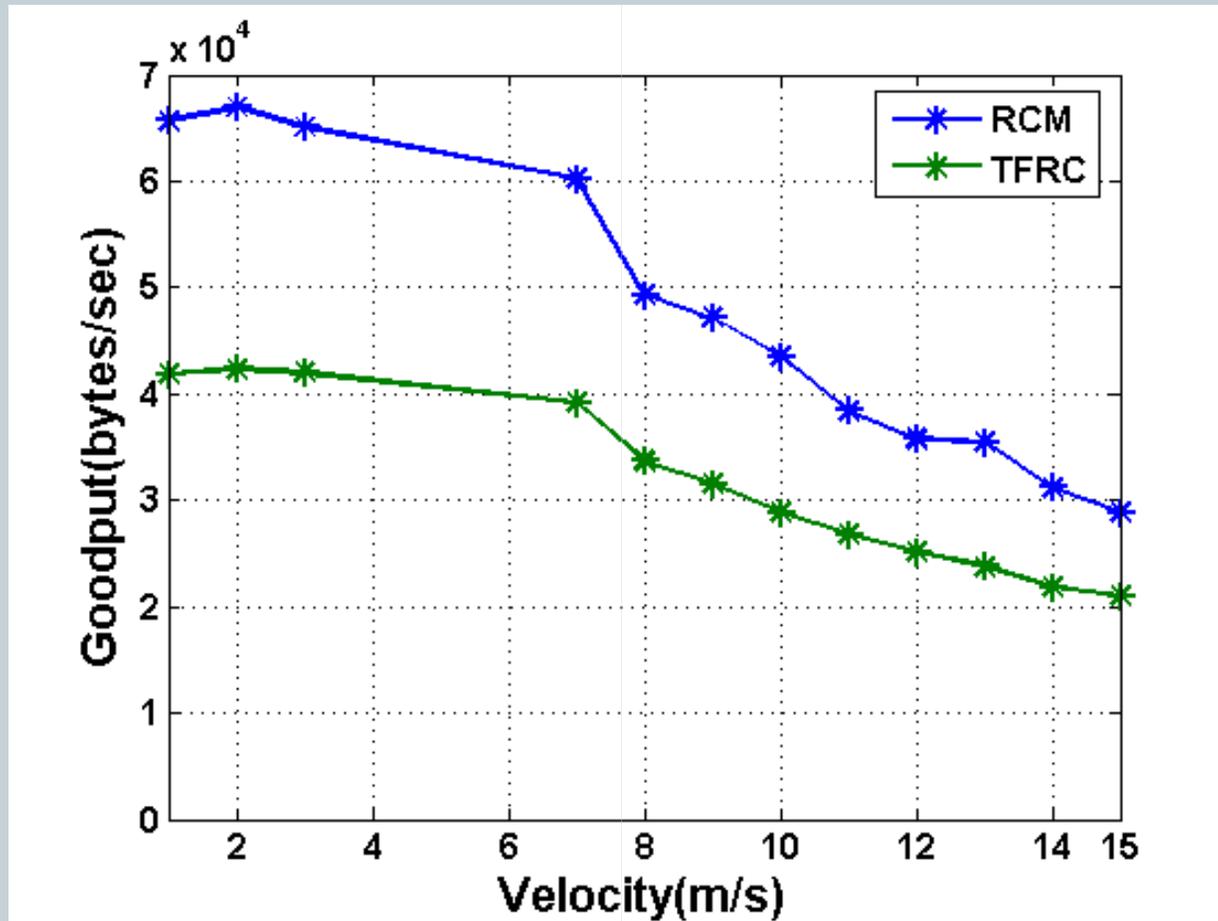
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# Good Put Vs Velocity of Mobile Node

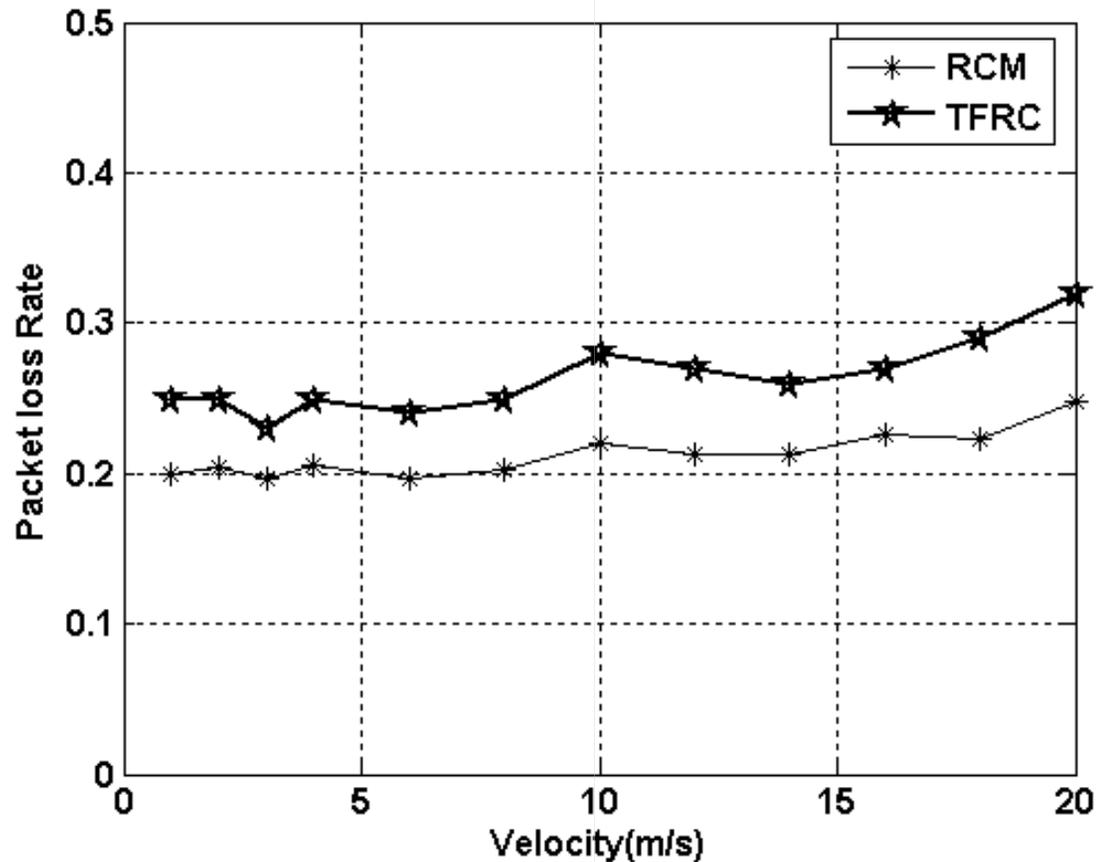
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# Packet loss rate Vs. Velocity of mobile node

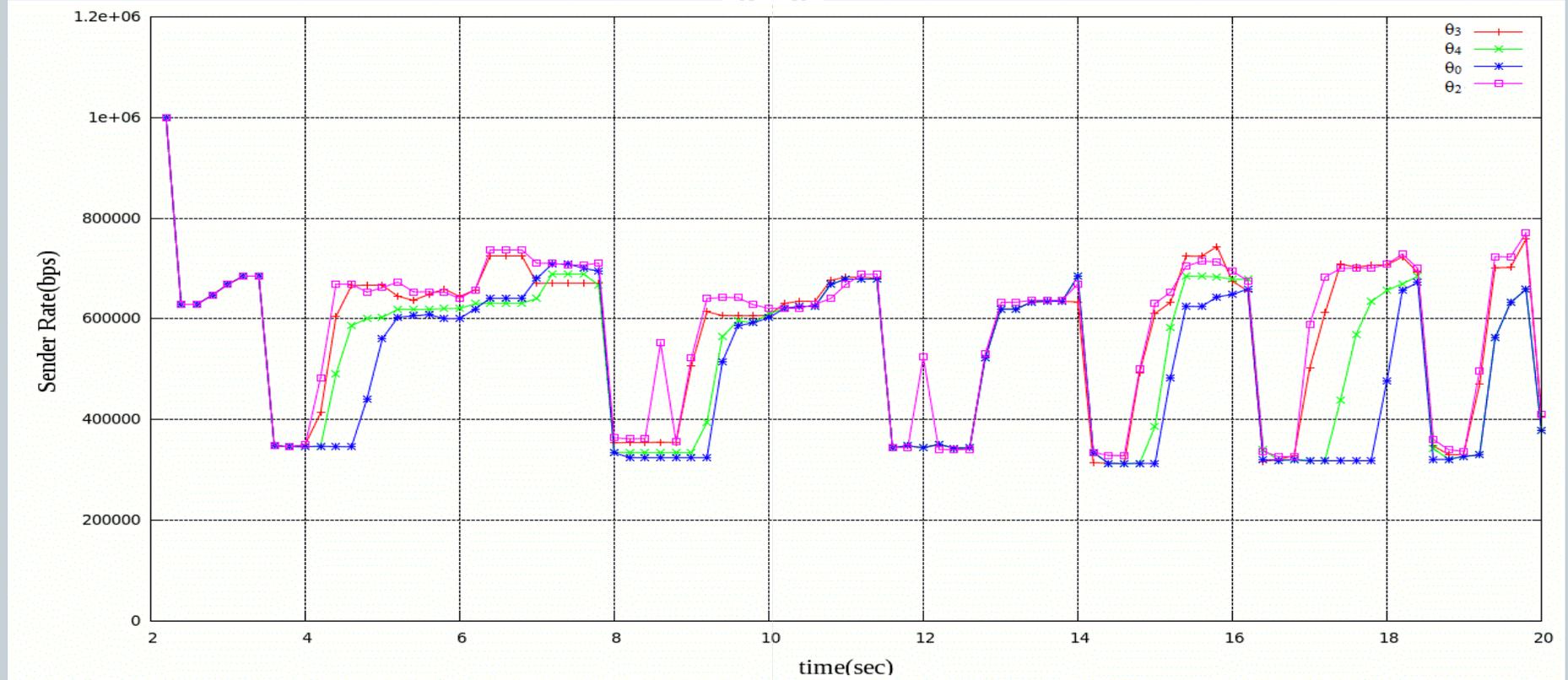
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# Analysis on decision factor( $\theta$ )

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Where,  $\theta_0 = 0$ ,

$$\theta_4 = \max(2, \text{prevrate} * \text{rtt} / 4)$$

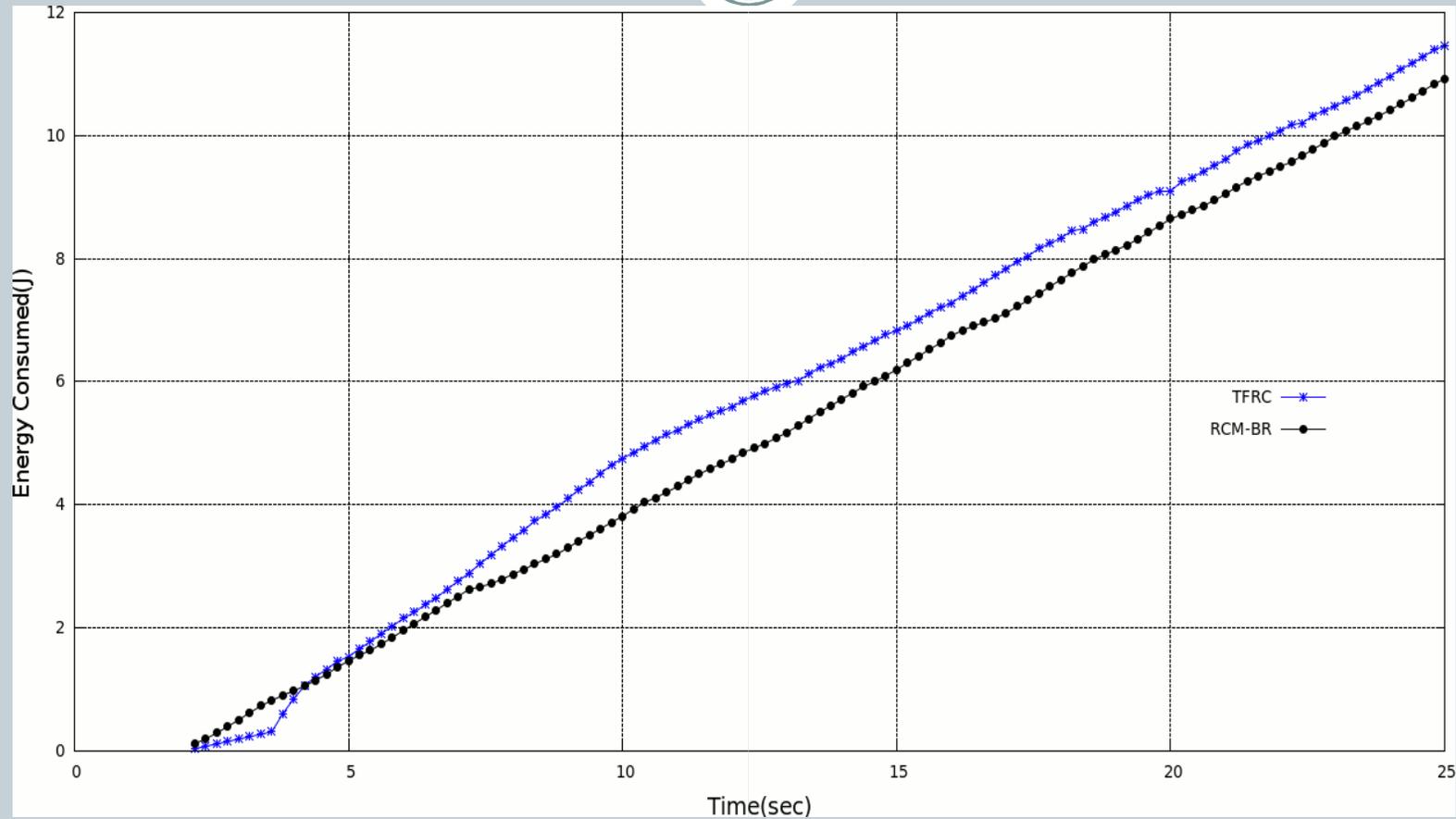
$$\theta_2 = \max(2, \text{prevrate} * \text{rtt} / 2) \text{ and}$$

$$\theta_3 = \max(2, \text{prevrate} * \text{rtt} / 3)$$



# Total Energy consumed

41

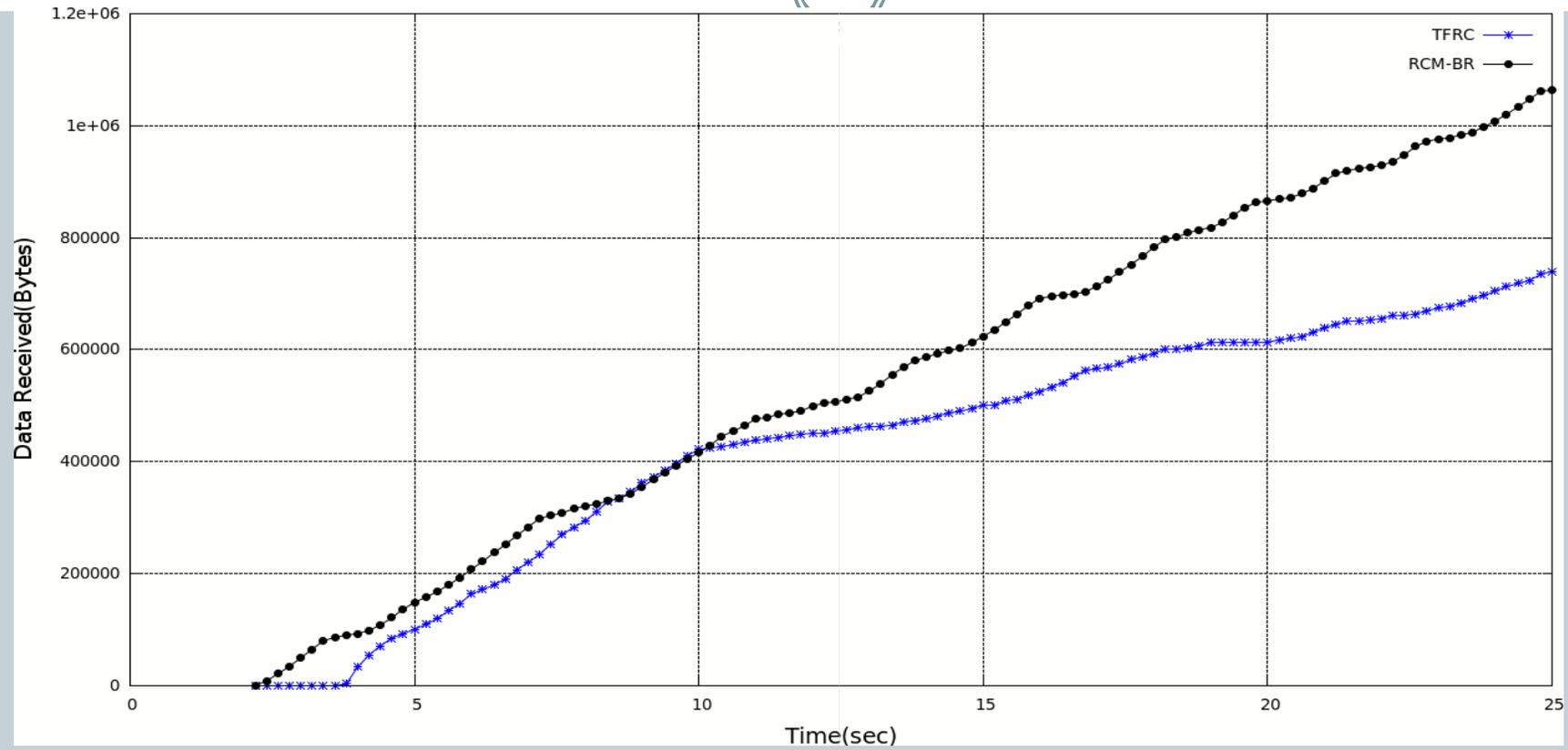


- With  $w_p=0.2$



# Data received

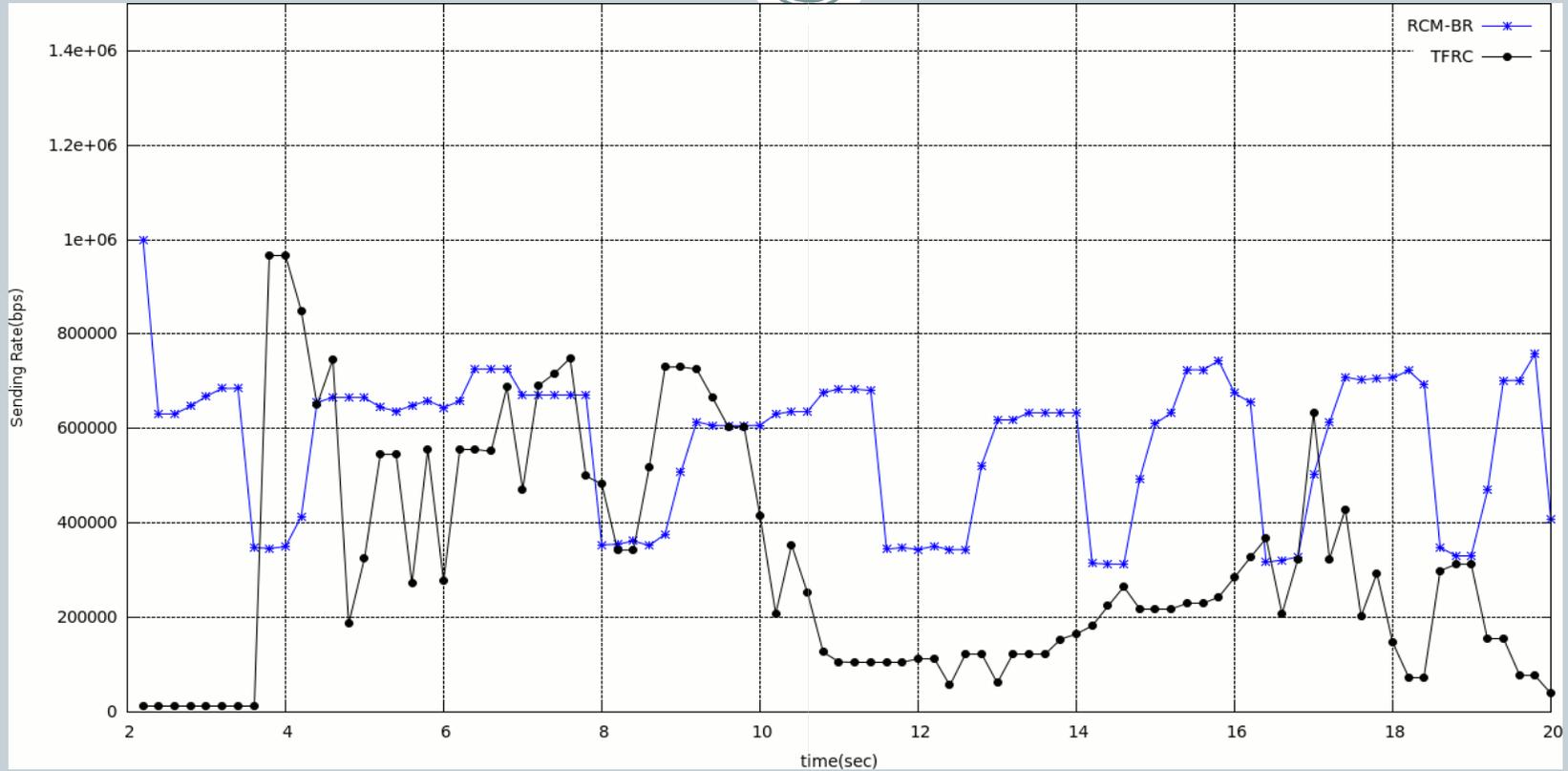
(42)





# Sender rate

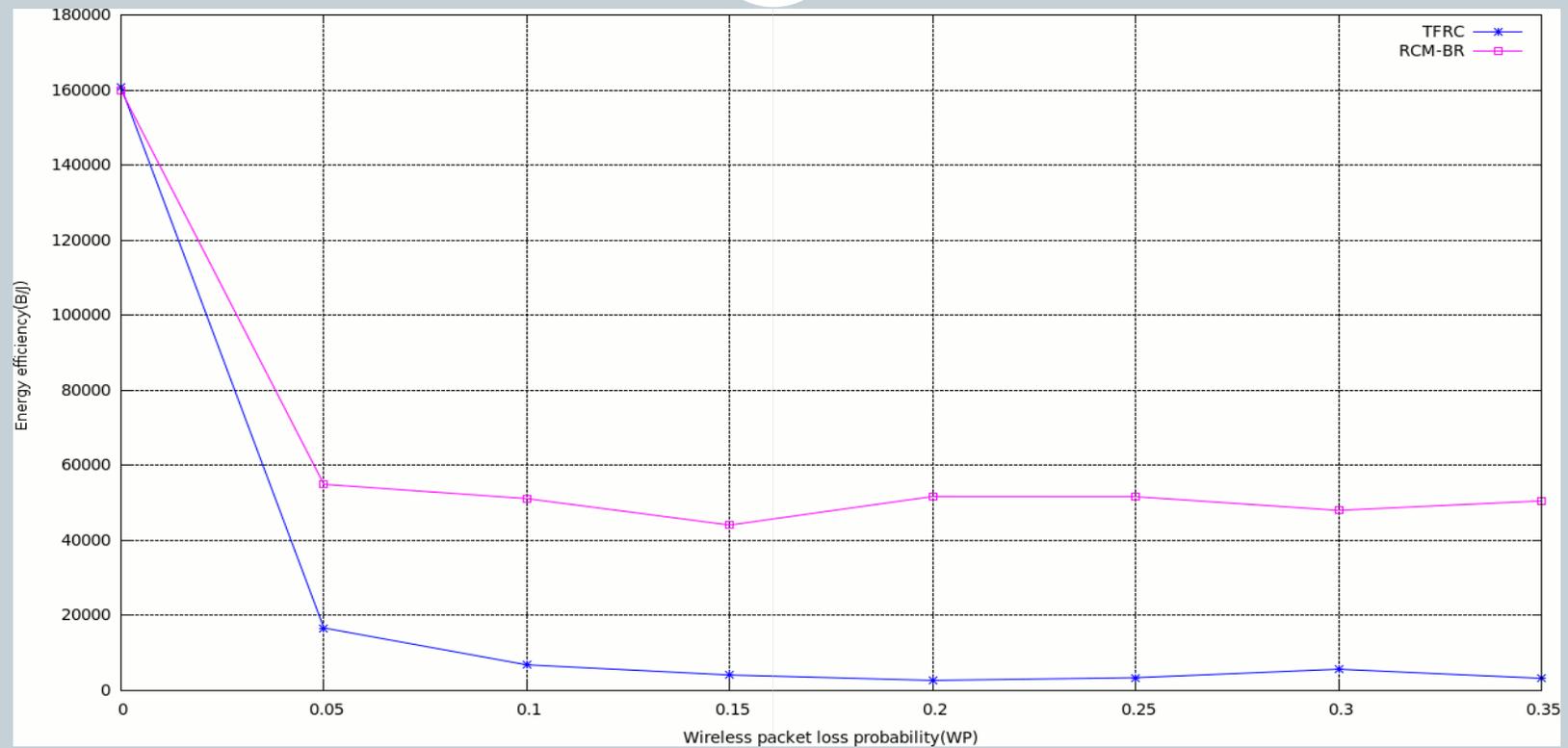
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# Energy efficiency

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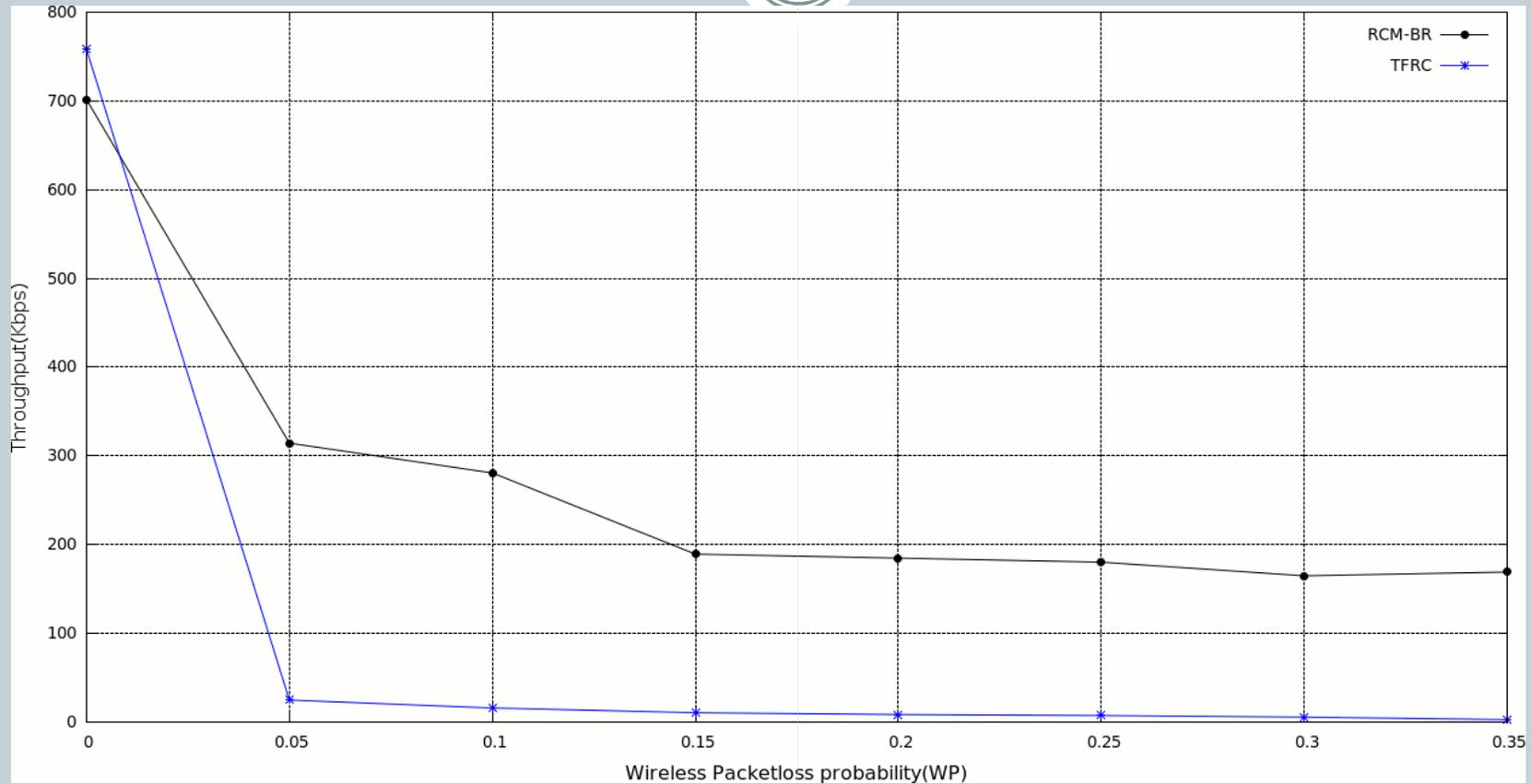


- Energy efficiency Vs Wireless loss probability



# Throughput

45

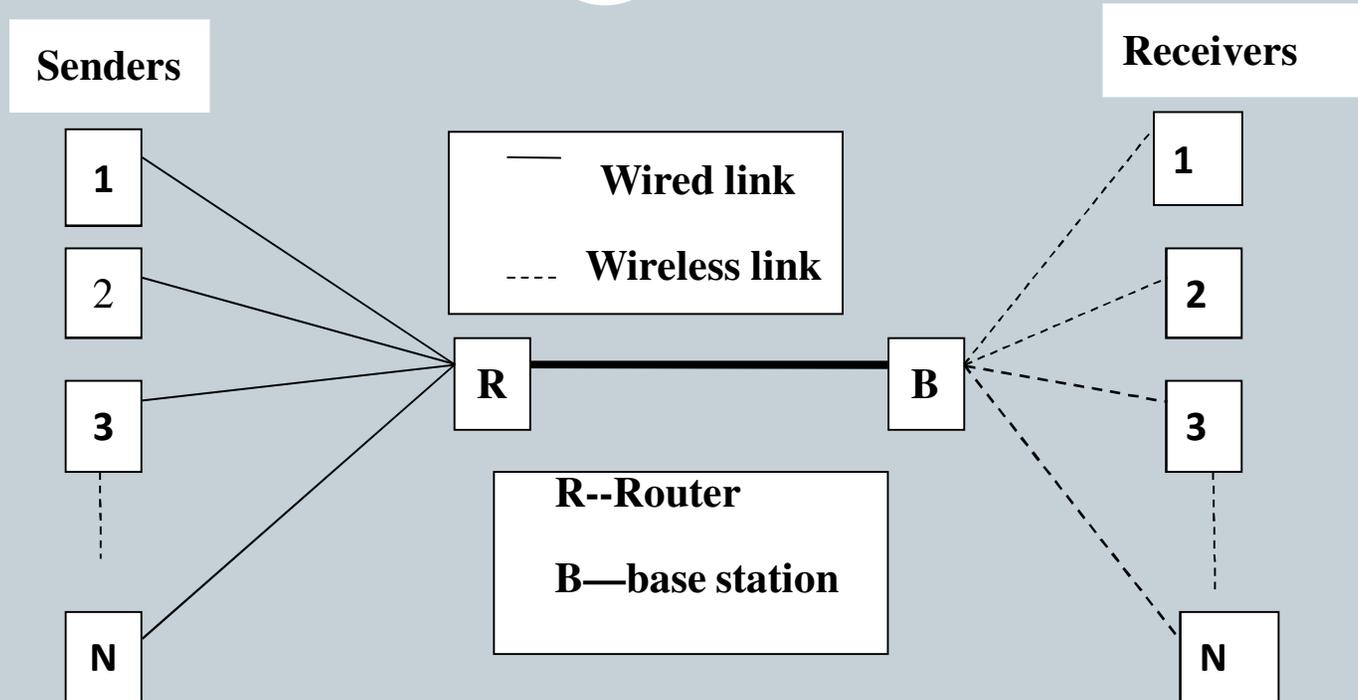


- Throughput Vs Wireless loss probability



# Fairness

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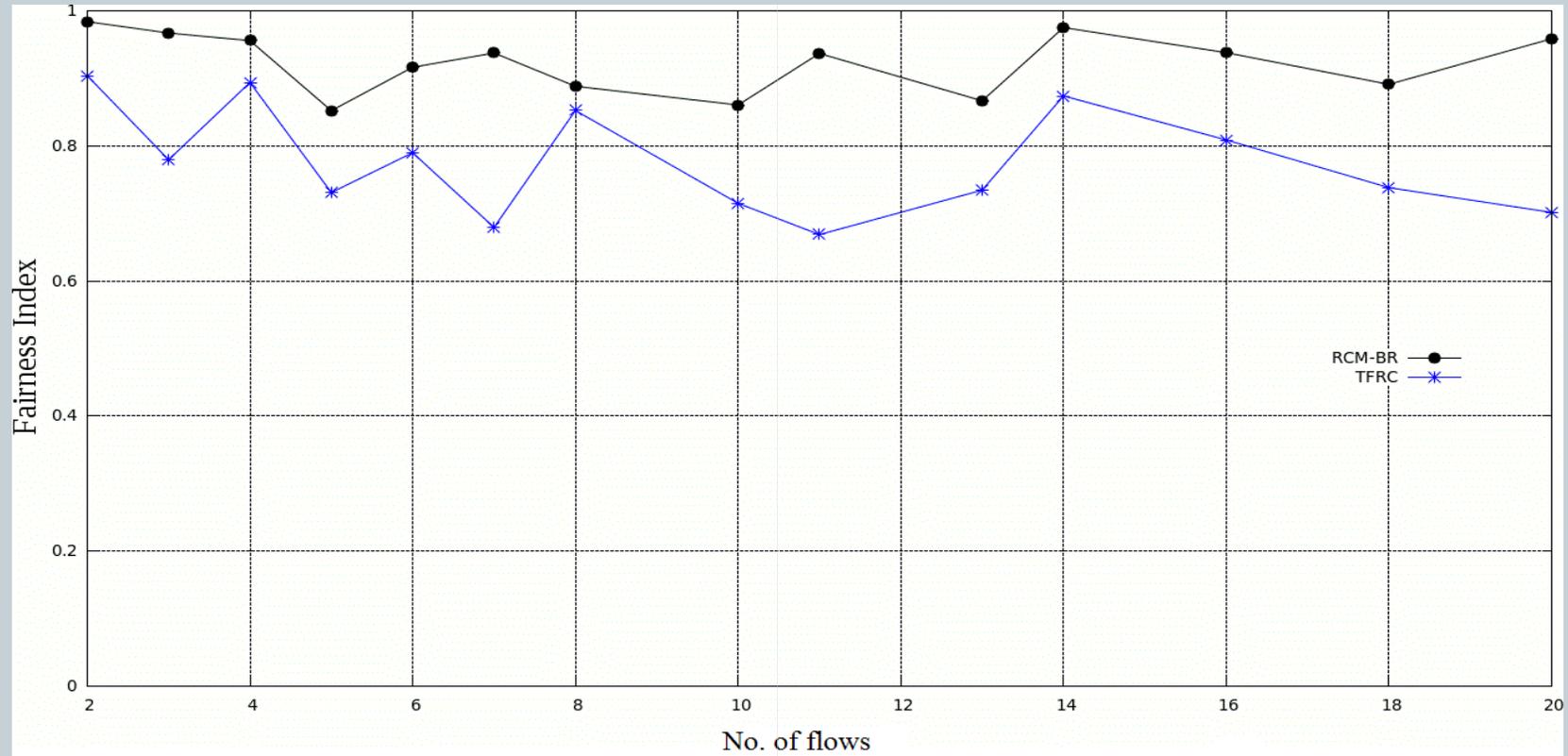
- With  $w_p=0.1$
- Starts between 1 and 3 sec



# Fairness Index

$$F.I = \frac{\left( \sum_{i=0}^n \text{Throughput}_i \right)^2}{n \left( \sum_{i=0}^n \text{Throughput}_i^2 \right)}$$

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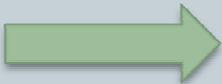
- FI Vs No. of flows



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# Conclusion

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- RCM protocol augmented with burst recovery(RCM-BR) improves the utilization and fairness of the transport protocol, while achieving high energy efficiency.
- Improve the end-to-end loss differentiation algorithm
- Improve burst detection mechanism



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**QUESTIONS?**



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# Back-up slides

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# Goodput?

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- **Goodput** is the application level throughput, i.e. the number of useful data bits delivered by the network per unit of time.
- Excludes protocol overhead bits as well as retransmitted data packets.

# RCM

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- "now" is the current time and "lastchange" is the time when RTT was changed last.
- RTT is round trip time calculated using exponential weighted moving average.
- In the equation  $S_{init} = \max(1, n_{ack}) / RTT$ ,
- "n\_ack" is the number of acknowledgements received by the sender for the probe packets sent in second RTT time.
- Initially the probe packets will be sent with a rate of "S\_target", which will be given by the application to achieve highest quality (maximum rate needed). In the simulation it is given manually in the TCL file with the variable like target\_rate.

# Goodput?

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- For example, if a file is transferred, the goodput that the user experiences corresponds to the file size in bits divided by the file transfer time. It is lower than the throughput.
- Examples of factors that cause lower goodput than throughput are:
  - *Protocol overhead*, such as transport layer, network layer and sometimes datalink layer protocol overhead is included in the throughput, but is excluded from the goodput.
  - *Transport layer flow control and congestion avoidance*, e.g. TCP slow start, may cause a lower goodput than the maximum throughput.
  - Retransmission of lost or corrupt packets due to transport layer automatic repeat request (ARQ), caused by bit errors or packet dropping in congested switches and routers, is included in the datalink layer or network layer throughput but not in the goodput.

# Performance of Loss differentiation algorithm

- **Accuracy of Adaptive loss differentiation without link layer ARQ**

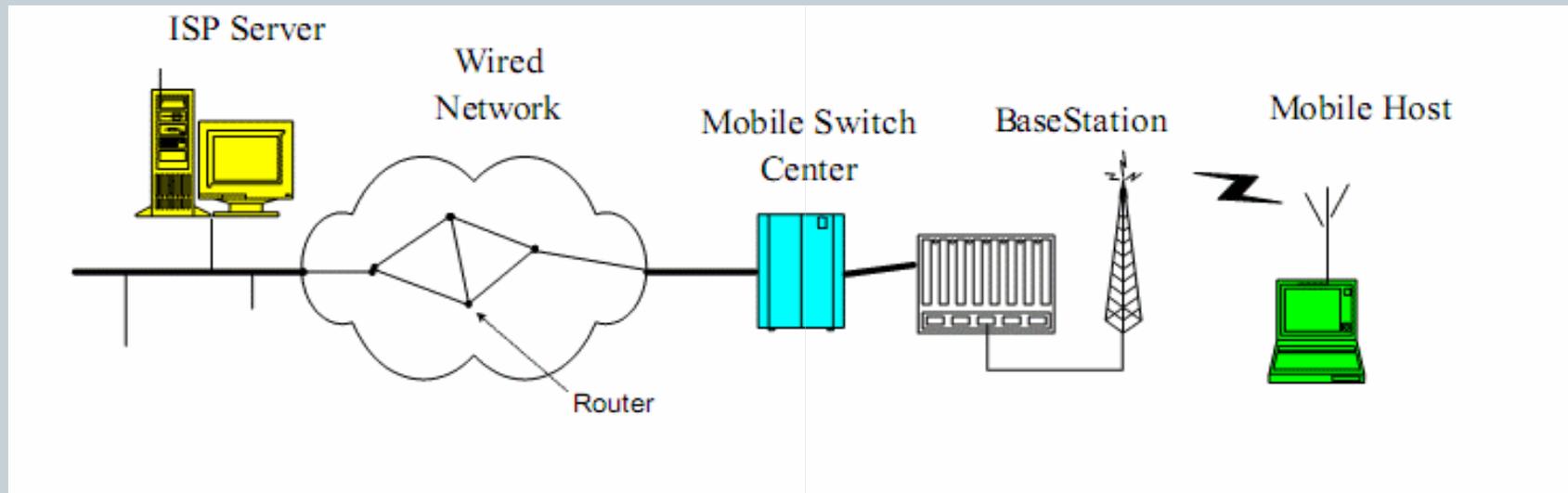
C/C	W/W	C/W	W/C
82.5%	85.38%	14.62%	17.5%

## **Accuracy of Adaptive loss differentiation with link layer ARQ**

C/C	W/W	C/W	W/C
95.3	0%	0%	4.7

# Wireless internet

60



Combination of wired and wireless networks