

Final Exam

CIS 341: Introduction to Logic and Automata — Spring 2002, day sections

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Print Name: _____

Student Number: _____

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

Signature and Date

- This exam has 12 pages in total, numbered 1 to 12. Make sure your exam has all the pages.
- This exam will be 2.5 hours in length.
- This is an open-book, open-note exam.
- For all problems, follow these instructions:
 1. Show your work and give reasons (except for question 1).
 2. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area to work out your answers before filling in the answer space.
 3. FA stands for finite automaton; TG stands for transition graph; CFG stands for context-free grammar; CFL stands for context-free language; PDA stands for push-down automaton; TM stands for Turing machine.
 4. For any proofs, be sure to provide a step-by-step argument, with justifications (e.g., cite a theorem or definition in the textbook) for every step. You may assume that the theorems in the textbook hold; i.e., you do not have to reprove the theorems in the textbook. When using a theorem or result from the textbook, make sure you refer to it by number (e.g., Theorem 3) or page number.

Problem	1	2	3	4	5	6	Total
Points							

1. [20 points] For each of the following, circle TRUE if the statement is always correct. Otherwise, circle FALSE

- (a) TRUE FALSE — If a language L is accepted by a nondeterministic finite automaton, then there must be some PDA that also accepts L .
- (b) TRUE FALSE — If L is generated by a context-free grammar that is *not* a regular grammar, then L must *not* be a regular language.
- (c) TRUE FALSE — If L_1 is a context-free language and L_2 is not context-free, then L_1L_2 must be context-free.
- (d) TRUE FALSE — If L_1 is a context-free language and L_2 is not context-free, then L_1L_2 must not be context-free.
- (e) TRUE FALSE — Every context-free language is a nonregular language.
- (f) TRUE FALSE — Every nondeterministic PDA can be transformed into an equivalent deterministic PDA.
- (g) TRUE FALSE — We can construct a PDA for the language $L = \{a^{2n}b^n : n = 1, 2, 3, \dots\}$ by first constructing an FA for L , and then converting the FA into a PDA for L .
- (h) TRUE FALSE — If w is an input string and T is a transition graph, then T must either accept or reject w .
- (i) TRUE FALSE — If w is an input string and T is a Turing machine, then T must either accept or reject w .
- (j) TRUE FALSE — There is a Turing machine that can take as input any encoded Turing machine P and any input w for P and decide if P halts on input w .

2. [20 points] For each of the following multiple-choice questions, circle the letter of the correct answer.

2.1. If a language L is finite, then

- (a) there must not exist a pushdown automaton for L .
- (b) L must have a transition graph.
- (c) L must be a non-regular language.

2.2. If L_1 and L_2 are context-free languages, then $L_1 \cap L_2$

- (a) must be context-free.
- (b) must not be context-free.
- (c) may or may not be context-free.

2.3. If L is a language accepted by a nondeterministic finite automaton, then L

- (a) must be context-free.
- (b) must not be context-free.
- (c) may or may not be context-free.

2.4. If L is not a context-free language, then

- (a) L must be finite.
- (b) L must be a non-regular language.
- (c) L must be a regular language.

2.5. If L is generated by a context-free language in Chomsky normal form, then

- (a) L must be regular.
- (b) L must be non-regular.
- (c) L can be either regular or non-regular.

2.6. If L is accepted by a finite automaton, then

- (a) there must exist some Turing machine that accepts L .
- (b) there must not exist some Turing machine that accepts L .
- (c) there may or may not exist some Turing machine that accepts L .

2.7. If L is a language accepted by some Turing machine M , then

- (a) $w \in L'$ must be rejected by M .
- (b) running M with $w \in L$ initially loaded on the TAPE will result in M ending in a HALT state.
- (c) L must be a nonregular language.

2.8. Let $\Sigma = \{a, b\}$, and let L be the language over Σ exactly consisting of all strings that do not contain the substring aaa . Consider the following regular expressions:

- (i) $(a + b)^*(a + aa)^*(a + b)^*$
- (ii) $(b + ab + aab)^*(\Lambda + a + aa)$
- (iii) $(\Lambda + a + aa)(b + ba + baa)^*$

Which of the following is correct?

- (a) Only regular expression (i) generates L .
- (b) Only regular expression (ii) generates L .
- (c) Only regular expression (iii) generates L .
- (d) Only regular expressions (i) and (ii) generate L .
- (e) Only regular expressions (i) and (iii) generate L .
- (f) Only regular expressions (ii) and (iii) generate L .
- (g) All 3 regular expressions generate L .
- (h) None of the 3 regular expressions generates L .

2.9. Let $\Sigma = \{a, b\}$, and let $L = \{a^n w a^n : n \geq 1, w \in \Sigma^*\}$. Consider the following statements:

- (i) L has regular expression $a^*(a + b)^*a^*$.
- (ii) L is a non-regular language.
- (iii) L has context-free grammar

$$S \rightarrow aSa \mid aS \mid bS \mid aa$$

where S, X are nonterminals.

- (iv) L has context-free grammar

$$\begin{aligned} S &\rightarrow aSa \mid aXa \\ X &\rightarrow aX \mid bX \mid \Lambda \end{aligned}$$

where S, X are nonterminals.

Which of the following is true?

- (a) Only statement (i) is correct.
- (b) Only statement (ii) is correct.
- (c) Only statement (iii) is correct.
- (d) Only statement (iv) is correct.
- (e) Only statements (i) and (iii) are correct.
- (f) Only statements (i) and (iv) are correct.
- (g) Only statements (ii) and (iii) are correct.
- (h) Only statements (ii) and (iv) are correct.
- (i) Only statements (ii), (iii) and (iv) are correct.
- (j) All the 4 statements are correct.
- (k) None of the 4 statements is correct.

2.10. Suppose L is a language generated by a regular grammar, and consider the following statements:

- (i) L is a regular language.
- (ii) L is a context-free language.
- (iii) There is a Turing machine that accepts L .

Which of the following is correct?

- (a) Only statement (i) is true.
- (b) Only statement (ii) is true.
- (c) Only statement (iii) is true.
- (d) Only statements (i) and (ii) are true.
- (e) Only statements (i) and (iii) are true.
- (f) Only statements (ii) and (iii) are true.
- (g) All 3 statements are true.
- (h) None of the 3 statements is true.

3. [20 points] Let $\Sigma = \{a, b\}$. Each of the following languages L defined over Σ falls into one of the following categories:

- (i) L is a regular language.
- (ii) L is a context-free language, but not a regular language.
- (iii) L is recursively enumerable, but not a context-free language.

For each of the following languages, specify which category it is in. If a language L is in category (i), give a regular expression for L . If a language L is in category (ii), give a context-free grammar for L . If a language L is in category (iii), you only need to specify that it is in this category.

- (a) L exactly consists of all strings $w \in \Sigma^*$ such that $w = \text{reverse}(w)$ and the length of w is divisible by 4 (i.e., w has length $4n$ for some $n = 0, 1, 2, 3, \dots$).

Circle one: Category (i) Category (ii) Category (iii)

- (b) $L = \{b^n a^n b^n : n \geq 0\}$.

Circle one: Category (i) Category (ii) Category (iii)

(c) $L = \{a^{2n} : n \geq 0\}$.

Circle one: Category (i) Category (ii) Category (iii)

(d) $L = \{a^m b^m a^n b^n : m \geq 0, n \geq 0\}$.

Circle one: Category (i) Category (ii) Category (iii)

Scratch-work area for problem 3.

4. [10 points] Let $L = \{a^{3n}b^{2n} : n = 0, 1, 2, \dots\}$.

(a) [5 points] Give a pushdown automaton that accepts L .

Draw PDA here:

(b) [5 points] Give a Turing machine that accepts L .

Draw Turing machine here:

Scratch-work area for problem 4.

5. **[10 points]** Let L_1, L_2, L_3, \dots be an infinite sequence of context-free languages, each of which is defined over a common alphabet Σ . Let L be the infinite union of L_1, L_2, L_3, \dots ; i.e., $L = L_1 + L_2 + L_3 + \dots$. Is it always the case that L is a context-free language?

YES NO (Circle one)

If your answer is YES, give a proof. If your answer is NO, give a counterexample. Explain your answer.

6. [20 points] Let L_1 , L_2 , and L_3 be languages defined over the alphabet $\Sigma = \{a, b\}$, where

- L_1 consists of all possible words over Σ except the words w_1, w_2, \dots, w_{100} ; i.e., start with all possible words over the alphabet, take out 100 particular words, and the remaining words form the language L_1 ;
- L_2 can be generated by a context-free grammar that is a regular grammar; and
- L_3 is accepted by some push-down automaton.

Prove that $(L_1 \cap L_2)L_3$ is a context-free language.