

Midterm Exam I  
CIS 341: Introduction to Logic and Automata — Spring 2002  
Prof. Marvin K. Nakayama

Print Name (last name first): \_\_\_\_\_

Student Number: \_\_\_\_\_

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

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Signature and Date

- This exam has 6 pages in total, numbered 1 to 6. Make sure your exam has all the pages.
- This exam will be 1 hour and 25 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:
  1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area to work out your answers before filling in the answer space.
  2. FA stands for finite automaton; TG stands for transition graph.
  3. For any proofs, be sure to provide a step-by-step argument, with justifications for every step.

Problem	1	2	3	4	5	Total
Points						

1. [20 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

- (a) TRUE FALSE — If  $L$  is any language, then  $L^*$  must be infinite.
- (b) TRUE FALSE — If  $L$  is any language, then  $L^*$  must be finite.
- (c) TRUE FALSE — If  $L_1$  and  $L_2$  are languages such that  $L_1^* = L_2^*$ , then  $L_1 = L_2$ .
- (d) TRUE FALSE — A regular expression for the language  $L = \{a^n : n = 1, 2, 3, \dots\}$  is  $\mathbf{a + aa + aaa + \dots}$ .
- (e) TRUE FALSE — If  $L_1$  is language having regular expression  $r_1$  and  $L_2$  is language having regular expression  $r_2$ , then the language  $L_1 + L_2$  has regular expression  $r_1 + r_2$ .
- (f) TRUE FALSE — Let  $L$  be the language over  $\Sigma = \{a, b\}$  consisting of exactly all words that have either an even number of  $a$ 's or an even number of  $b$ 's. Then a regular expression for  $L$  is  $\mathbf{(aa + bb + (ab + ba)(aa + bb)^*(ab + ba))^*}$ .
- (g) TRUE FALSE — All transition graphs are non-deterministic.
- (h) TRUE FALSE — If a transition graph accepts the string  $\Lambda$ , then some start state must also be a final state.
- (i) TRUE FALSE — If a finite automaton accepts the string  $\Lambda$ , then the start state must also be a final state.
- (j) TRUE FALSE — If  $L_1 = \{\Lambda\}$  and  $L_2 = \emptyset$ , then  $L_1 = L_2$ .

2. [20 points] For each of the following languages  $L$  over the alphabet  $\Sigma = \{a, b\}$ , give a regular expression for  $L$ .

(a)  $L$  exactly consists of all strings that begin with  $a$  and end with  $b$ .

**Regular Expression:** \_\_\_\_\_

(b)  $L$  exactly consists of all strings that have an odd number of  $a$ 's.

**Regular Expression:** \_\_\_\_\_

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Scratch-work area

3. [20 points] For each of the following languages  $L$  over the alphabet  $\Sigma = \{a, b\}$ , give a finite automaton that accepts exactly  $L$ .

(a)  $L$  exactly consists of all strings that have length of either 1 or 2.

**Draw finite automaton here:**

(b)  $L$  exactly consists of all strings that do not contain the substring  $aba$ .

**Draw finite automaton here:**

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Scratch-work area

4. [20 points] For each of the following parts, provide an example satisfying the given conditions. Give a brief explanation for each of your examples.

(a) Give an example of a set  $S$  such that  $S^* = S^+$ .

(b) Give an example of a set  $S$  such that  $S = S^*$ .

(c) Give an example of a set  $S$  such that  $S \neq S^+$ .

(d) Give an example of a set  $S$  such that  $S^*$  is finite.

5. [20 points] Suppose we define a restricted version of the C++ programming language in which variable names must satisfy all of the following conditions:

- (i) A variable name can only use Roman letters (i.e., a, b, ..., z, A, B, ..., Z) or Arabic numerals (i.e., 0, 1, 2, ..., 9); i.e., underscore is not allowed.
- (ii) A variable name must start with a Roman letter: a, b, ..., z, A, B, ..., Z
- (iii) The length of a variable name must be no greater than 8.
- (iv) A variable name cannot be a keyword (e.g., if). The set of keywords is finite.

Let  $L$  be the set of all valid variable names in our restricted version of C++.

- (a) Let  $L_0$  be the set of strings satisfying conditions (i), (ii), and (iii) above; i.e., we do not require condition (iv). How many strings are in  $L_0$ ?

- (b) Prove that  $L$  has a regular expression.