## Midterm Exam I

CIS 341: Introduction to Logic and Automata - Spring 2002, evening Prof. Marvin K. Nakayama

Print Name (last name first): $\qquad$

Student Number: $\qquad$

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

Signature and Date

- This exam has 8 pages in total, numbered 1 to 8 . Make sure your exam has all the pages.
- This exam will be 3 hour and 5 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:

1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area to work out your answers before filling in the answer space.
2. FA stands for finite automaton; TG stands for transition graph.
3. For any proofs, be sure to provide a step-by-step argument, with justifications for every step.

| Problem | 1 | 2 | 3 | 4 | 5 | 6 | 7 | Total |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Points |  |  |  |  |  |  |  |  |

1. [10 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE
(a) TRUE FALSE - All languages are infinite.
(b) TRUE FALSE - All regular languages are infinite.
(c) TRUE FALSE - All nonregular languages are infinite.
(d) TRUE FALSE - All regular languages are finite.
(e) TRUE FALSE - If a transition graph accepts some language $L$, then there exists some nondeterministic finite automaton for the language $L^{\prime}$.
(f) TRUE FALSE - If $M$ is a finite automaton, then $M$ is also a transition graph.
(g) TRUE FALSE - If $L_{1}$ is a regular language, then so is $L_{1}^{*}$.
(h) TRUE FALSE - If $L$ is a nonregular language, then there exists a nondeterministic finite automaton that accepts $L$.
(i) TRUE FALSE - If $L_{1}$ and $L_{2}$ are languages such that $L_{1} \subset L_{2}$ and $L_{2}$ is nonregular, then $L_{1}$ must be nonregular.
(j) TRUE FALSE - The regular expressions $(\mathbf{a}+\mathbf{b})^{*}$ and $\left(\mathbf{a}^{*} \mathbf{b}^{*}\right)$ generate the same language.
2. [15 points] For each of the following languages $L$ over the alphabet $\Sigma=\{a, b\}$, give a regular expression for $L$.
(a) $L$ exactly consists of all strings containing exactly $2 a$ 's.

## Regular Expression:

(b) $L$ exactly consists of all strings with an even number of $a$ 's or an even number of $b$ 's. (Note that this says "or", not "and".)

Regular Expression:

## Scratch-work area

3. [15 points] For each of the following languages $L$ over the alphabet $\Sigma=\{a, b\}$, give a finite automaton that accepts exactly $L$.
(a) $L$ exactly consists of all strings whose first letter is $a$.

## Draw finite automaton here:

(b) $L$ exactly consists of all strings that do not end in $b b b$.

Draw finite automaton here:

Scratch-work area
4. [15 points] Suppose that $T$ is a transition graph with language $L$ defined over an alphabet $\Sigma$. Suppose that a friend claims that you can create a transition graph $T^{\prime}$ for the language $L^{\prime}$ as follows:

- $T^{\prime}$ has alphabet $\Sigma$.
- The states and arcs in $T^{\prime}$ are the same as those in $T$, and the arcs have the same labels.
- The start states of $T^{\prime}$ are the same as the start states of $T$.
- Every final state in $T$ is a non-final state in $T^{\prime}$, and every non-final state in $T$ is a final state in $T^{\prime}$.

Show that your friend is wrong by giving an example of a transition graph $T$ having a language $L$ such that if we construct $T^{\prime}$ using the above rules, the language of $T^{\prime}$ is not $L^{\prime}$. Be sure to explain your answer.

## Scratch-work area

5. [15 points] For each of the following parts, provide an example satisfying the given conditions. Give a brief explanation for each of your examples.
(a) Give an example of a set $S$ such that $S=S^{+}$.
(b) Give an example of sets $S$ and $T$ such that $S T=S$.
(c) Give an example of sets $S$ and $T$ such that $S T=\emptyset$.

## Scratch-work area

6. [15 points] Assume the following hold:

- $\Sigma$ is an alphabet consisting of $t$ letters.
- $L_{1}$ and $L_{2}$ are languages defined over $\Sigma$.
- There is a finite automaton $F A_{1}$ that accepts $L_{1}$, and $F A_{1}$ has $m$ states.
- There is a finite automaton $F A_{2}$ that accepts $L_{2}$, and $F A_{2}$ has $n$ states.

In class, we went over an algorithm to construct a finite automaton $F A_{3}$ for language $L_{1} L_{2}$ from $F A_{1}$ and $F A_{2}$. What is the maximum number of states in $F A_{3}$ constructed using this algorithm? Show your work.

## Answer:

7. [15 points] Recall the pumping lemma:

Theorem 14 Let $L$ be a language accepted by a finite automaton having $N$ states, and let $w \in L$ with length $(w) \geq N$. Then there exists strings $x, y$, and $z$ such that
(i) $w=x y z$
(ii) $y \neq \Lambda$
(iii) length $(x)+$ length $(y) \leq N$
(iv) $x y^{k} z \in L$ for all $k=0,1,2, \ldots$.

Prove that the language $L=\left\{b^{m} a b^{m}: m=0,1,2, \ldots\right\}$ is nonregular.

