

Midterm Exam I  
CIS 341: Introduction to Logic and Automata — Fall 2004, day  
Prof. Marvin K. Nakayama

Print Family (i.e., Last) Name: \_\_\_\_\_

Print Given (i.e., First) Name: \_\_\_\_\_

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

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Signature and Date

- This exam has 6 pages in total, numbered 1 to 6. Make sure your exam has all the pages.
- This exam will be 1 hour and 25 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:
  1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the sheets to work out your answers before filling in the answer space.
  2. FA stands for finite automaton; TG stands for transition graph.
  3. For any proofs, be sure to provide a step-by-step argument, with justifications for every step.

Problem	1	2	3	4	Total
Points					

1. [30 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

- (a) TRUE FALSE — All finite automata are deterministic.
- (b) TRUE FALSE — All transition graphs are nondeterministic.
- (c) TRUE FALSE — A finite automaton must have exactly one initial state.
- (d) TRUE FALSE — The language  $L = \{a^n b^n : n \geq 0\}$  has regular expression  $\mathbf{a^*b^*}$ .
- (e) TRUE FALSE — A transition graph  $M = (K, \Sigma, \Pi, S, F)$  accepts a string  $w \in \Sigma^*$  if and only if every possible way of processing  $w$  on  $M$  results in ending in a final state with no unread letters left and without crashing.
- (f) TRUE FALSE —  $\Lambda \in \{\Lambda\}$ .
- (g) TRUE FALSE —  $L$  has a regular expression only if  $L$  is finite.
- (h) TRUE FALSE — If  $L$  has regular expression  $\mathbf{a^*}$ , then  $\{a\} \in L$ .
- (i) TRUE FALSE — If two sets  $A$  and  $B$  satisfy  $A = B$ , then  $A \subset B$ .
- (j) TRUE FALSE — A transition graph may have no final states.

2. [20 points] For each of the following languages  $L$  over the alphabet  $\Sigma = \{0, 1\}$ , give a regular expression for  $L$ .

(a)  $L = \{w \in \Sigma^* : w = s1 \text{ for some } s \in \Sigma^*\}$ .

**Regular Expression:** \_\_\_\_\_

(b)  $L$  consists of all strings over  $\Sigma$  that contain neither 00 nor 11 as a substring.

**Regular Expression:** \_\_\_\_\_

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Scratch-work area

3. [30 points] For each of the following languages  $L$  over the alphabet  $\Sigma = \{a, b\}$ , give a finite automaton that accepts exactly  $L$ . **For each part below, give two representations of the finite automaton: first as a picture and then the formal definition of the finite automaton as  $(K, \Sigma, \pi, s, F)$ ; i.e., specify each of  $K, \Sigma, \pi, s,$  and  $F$  for each of your finite automata.**

(a)  $L = \{w \in \Sigma^* : |w| \leq 3\}$ .

**Draw picture of finite automaton here:**

**Give formal specification of finite automaton as  $(K, \Sigma, \pi, s, F)$  here:**

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**Scratch-work area**

- (b)  $L$  exactly consists of all strings  $w$  with  $|w| \geq 2$  and whose second-to-last letter is  $b$ .

**Draw picture of finite automaton here:**

**Give formal specification of finite automaton as  $(K, \Sigma, \pi, s, F)$  here:**

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**Scratch-work area**

4. [20 points] Prove that  $S^{++} = S^+$  for any set of strings  $S$ .