## Midterm Exam I

CIS 341: Introduction to Logic and Automata - Spring 2004, evening
Prof. Marvin K. Nakayama

Print Family (i.e., Last) Name: $\qquad$

Print Given (i.e., First) Name: $\qquad$

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

Signature and Date

- This exam has 5 pages in total, numbered 1 to 5 . Make sure your exam has all the pages.
- This exam will be 1 hour and 25 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:

1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the sheets to work out your answers before filling in the answer space.
2. FA stands for finite automaton; TG stands for transition graph.
3. For any proofs, be sure to provide a step-by-step argument, with justifications for every step.

| Problem | 1 | 2 | 3 | 4 | Total |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Points |  |  |  |  |  |

1. [30 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE
(a) TRUE FALSE - Suppose that $T$ is a transition graph defined with an alphabet $\Sigma$. If every state in $T$ is a final state, then $T$ accepts $\Sigma^{*}$.
(b) TRUE FALSE - A transition graph may crash when processing a string.
(c) TRUE FALSE - If $L$ is an infinite language, then $L=L^{*}$.
(d) TRUE FALSE - The regular expressions $\mathbf{a}^{*} \mathbf{b}^{*}$ and ( $\left.\mathbf{a b}\right)^{*}$ generate the same language.
(e) TRUE FALSE - If $\Lambda \in L$, then $L^{+}=L^{*}$.
(f) TRUE FALSE - For any language $L, \Lambda \in L^{+}$.
(g) TRUE FALSE - A transition graph may have no final states.
(h) TRUE FALSE - A finite automaton can have more than one final state.
(i) TRUE FALSE - If $L=\emptyset$, then $\Lambda \in L$.
(j) TRUE FALSE - A regular expression for the language $L=\left\{b^{n}: n \geq 0\right\}$ is $\boldsymbol{\Lambda}+\mathbf{b}+\mathbf{b b}+\mathbf{b b b}+\cdots$.
2. [25 points] For each of the following languages $L$ over the alphabet $\Sigma=\{0,1\}$, give a regular expression for $L$.
(a) $L$ exactly consists of all strings that end in 110 .

## Regular Expression:

(b) $L$ exactly consists of all strings that do not contain 00 as a substring.

## Regular Expression:

## Scratch-work area

3. [25 points] For each of the following languages $L$ over the alphabet $\Sigma=\{a, b\}$, give a finite automaton that accepts exactly $L$.
(a) $L$ exactly consists of all strings that begin with $b$.

## Draw finite automaton here:

(b) $L$ exactly consists of all strings in which the number of $a$ 's is divisible by 3 ; i.e., the number of $a$ 's in a string $w \in L$ is $3 n$ for some integer $n \geq 0$.

## Draw finite automaton here:

## Scratch-work area

4. [20 points] For any language $L$, define the transpose of $L$, denoted by $L^{t}$, to be the language of exactly those words that are words in $L$ spelled backward; i.e., $L^{t}=\{\operatorname{reverse}(w): w \in L\}$. For example, if $L=\{a, a b b, b b a a b, b b b a a\}$, then $L^{t}=\{a, b b a, b a a b b, a a b b b\}$. Prove that $\left(L_{1} L_{2}\right)^{t}=L_{2}^{t} L_{1}^{t}$ for languages $L_{1}$ and $L_{2}$.
