

Midterm Exam II  
CIS 341: Introduction to Logic and Automata — **Spring 2004, evening**  
Prof. Marvin K. Nakayama

Print family (or last) name: \_\_\_\_\_

Print given (or first) name: \_\_\_\_\_

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

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Signature and Date

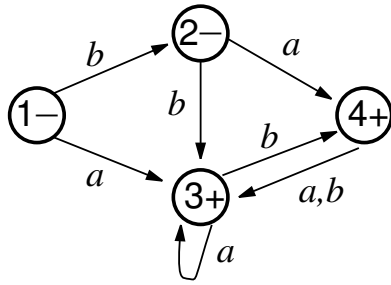
- This exam has 5 pages in total, numbered 1 to 5. Make sure your exam has all the pages.
- This exam will be 1 hour and 25 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:
  1. Show your work and give reasons (except for question 1).
  2. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the sheets to work out your answers before filling in the answer space.
  3. FA stands for finite automaton; TG stands for transition graph.
  4. For any proofs, be sure to provide a step-by-step argument, with justifications for every step.

Problem	1	2	3	4	Total
Points					

1. [30 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

- (a) TRUE FALSE — If  $L$  is a nonregular language, then  $L$  has a nondeterministic finite automaton.
- (b) TRUE FALSE — A regular expression for the language  $L = \{a^n b^n : n \geq 0\}$  is  $\mathbf{a^*b^*}$ .
- (c) TRUE FALSE — If  $L_1$  and  $L_2$  are languages such that  $L_1 + L_2$  is regular, then  $L_1$  and  $L_2$  must be regular.
- (d) TRUE FALSE — If language  $L_1$  has a regular expression and language  $L_2$  has a transition graph, then  $L_1 \cap L_2$  must be regular.
- (e) TRUE FALSE — If a language  $L$  has a nondeterministic finite automaton, then  $L'$  also does.
- (f) TRUE FALSE — If  $L$  is infinite, then  $L$  is a nonregular language.
- (g) TRUE FALSE — If  $L$  has a context-free grammar, then  $L$  is a nonregular language.
- (h) TRUE FALSE — If a language  $L$  has a transition graph, then  $L^*$  has a regular expression.
- (i) TRUE FALSE — If  $L_1$  and  $L_2$  are nonregular languages, then  $L_1 \cap L_2$  must be nonregular.
- (j) TRUE FALSE — An effective procedure to decide if a finite automaton with  $N$  states accepts any strings over an alphabet  $\Sigma$  is to keep testing strings over  $\Sigma$  on the FA until a string is accepted.

2. [30 points] Let  $L$  be the language of finite automata below.



Give a regular expression for  $L$ .

Answer: \_\_\_\_\_

**Scratch-work area**

3. [20 points] Suppose that a language  $L_1$  is accepted by a finite automaton  $FA_1 = (K_1, \Sigma, \pi_1, s_1, F_1)$ , and a language  $L_2$  is accepted by a finite automaton  $FA_2 = (K_2, \Sigma, \pi_2, s_2, F_2)$ , where for  $FA_i$ ,  $i = 1, 2$ ,  $K_i$  is the set of states,  $\Sigma$  is the input alphabet,  $\pi_i : K_i \times \Sigma \rightarrow K_i$  is the transition function,  $s_i$  is the initial state, and  $F_i$  is the set of final states. Define the language  $L_3 = L_1 \cap L_2$ , and for this question, you are to define a finite automaton  $FA_3 = (K_3, \Sigma, \pi_3, s_3, F_3)$  for  $L_3$  in terms of  $FA_1$  and  $FA_2$ , as we covered in class.

(a) Define the set  $K_3$  of states of  $FA_3$  in terms of  $FA_1$  and  $FA_2$ .

(b) Define the transition function  $\pi_3$  of  $FA_3$  in terms of  $FA_1$  and  $FA_2$ .

(c) Define the initial state  $s_3$  of  $FA_3$  in terms of  $FA_1$  and  $FA_2$ .

(d) Define the set  $F_3$  of final states of  $FA_3$  in terms of  $FA_1$  and  $FA_2$ .

4. [20 points] Recall the pumping lemma:

**Theorem 14** *Let  $L$  be a language accepted by a finite automaton having  $N$  states, and let  $w \in L$  with  $\text{length}(w) \geq N$ . Then there exists strings  $x$ ,  $y$ , and  $z$  such that*

- (i)  $w = xyz$ ,
- (ii)  $y \neq \Lambda$ ,
- (iii)  $\text{length}(x) + \text{length}(y) \leq N$ ,
- (iv)  $xy^kz \in L$  for all  $k = 0, 1, 2, \dots$

Let  $L$  be the language defined by the context-free grammar with  $\Omega = \{S\}$ ,  $\Sigma = \{a, b\}$ , and productions

$$S \rightarrow aSa \mid bSb \mid a \mid b$$

Prove that  $L$  is a nonregular language.