

Print family (or last) name: _____

Print given (or first) name: _____

I have read and understand all of the instructions below, and I will obey the Academic Honor Code.

Signature and Date: _____

- This exam has 7 pages in total, numbered 1 to 7. Make sure your exam has all the pages.
- This exam will be 1 hour and 25 minutes in length.
- This is a closed-book, closed-note exam.
- For all problems, follow these instructions:
 1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the sheets to work out your answers before filling in the answer space.
 2. DFA stands for deterministic finite automaton; NFA stands for nondeterministic finite automaton; CFG stands for context-free grammar; PDA stands for pushdown automaton.
 3. For any proofs, be sure to provide a step-by-step argument, with justifications for every step. Unless you are specifically asked to prove a theorem from the book, you may assume that the theorems in the textbook hold; i.e., you do not have to reprove the theorems in the textbook. When using a theorem from the textbook, make sure you provide enough detail so that it is clear which result you are using; e.g., say something like, “By the theorem that states $S^{**} = S^*$, it follows that ...”

Problem	1	2	3	4	5	6	Total
Points							

1. [20 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

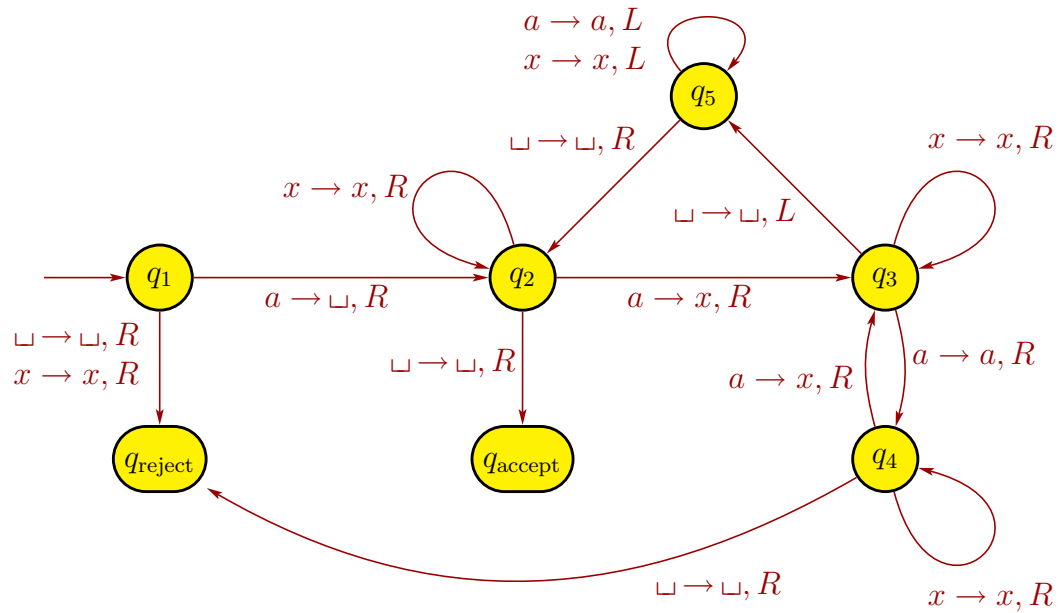
- (a) TRUE FALSE — If A is context-free, then A is Turing-recognizable.
- (b) TRUE FALSE — If A is context-free, then A is Turing-decidable.
- (c) TRUE FALSE — Every language is Turing-recognizable.
- (d) TRUE FALSE — For a Turing machine M and a string w , M either accepts or rejects w .
- (e) TRUE FALSE — The language $(0 \cup 1)^*$ is countable.
- (f) TRUE FALSE — The set \mathcal{B} of infinite binary sequences is countable.
- (g) TRUE FALSE — If language A is recognized by a 14-tape nondeterministic Turing machine, then there is a single-tape deterministic Turing machine that also recognizes A .
- (h) TRUE FALSE — The language EQ_{TM} is decidable, where
$$EQ_{\text{TM}} = \{ \langle M, N \rangle \mid M \text{ and } N \text{ are TMs with } L(M) = L(N) \}.$$
- (i) TRUE FALSE — The language EQ_{DFA} is decidable, where
$$EQ_{\text{DFA}} = \{ \langle C, D \rangle \mid C \text{ and } D \text{ are DFAs with } L(C) = L(D) \}.$$
- (j) TRUE FALSE — The language $\overline{A_{\text{TM}}}$ is decidable, where
$$A_{\text{TM}} = \{ \langle M, w \rangle \mid M \text{ is a TM that accepts input } w \}.$$

2. **[15 points]** Give a short answer (at most three sentences) for each part below. Be sure to define any notation that you use.

(a) What does it mean for a function $f : A \rightarrow B$ to be one-to-one?

(b) Explain the difference between Turing-recognizable and Turing-decidable.

3. [20 points] The Turing machine M below recognizes the language $A = \{ a^{2^n} \mid n \geq 0 \}$.



In each of the parts below, give the sequence of configurations that M enters when started on the indicated input string.

(a) aa

(b) $aaaaa$

Scratch-work area

Each of the following problems requires you to prove a result. In your proofs, you can apply any theorems that we went over in class without proving them, except for the result you are asked to prove in the problem. When citing a theorem, make sure that you give enough details so that it is clear what theorem you are using (e.g., say something like, “By the theorem that says every context-free language has a CFG in Chomsky normal form, we can show that . . .”)

4. **[15 points]** Let \mathcal{L} be the set of all languages having alphabet $\Sigma = \{0, 1\}$. Show that \mathcal{L} is uncountable.

5. **[15 points]** Consider the decision problem of testing whether two NFAs are equivalent. Express this problem as a language and show that it is decidable.

6. [15 points] Recall that

$$HALT_{TM} = \{ \langle M, w \rangle \mid M \text{ is a Turing machine that halts on input } w \}.$$

Prove that $HALT_{TM}$ is undecidable by showing that A_{TM} reduces to $HALT_{TM}$, where

$$A_{TM} = \{ \langle M, w \rangle \mid M \text{ is a Turing machine that accepts input } w \}.$$