

**CS 341, Spring 2009**  
**Solutions for Midterm 2**

1. (a) False, e.g.,  $\overline{A_{TM}}$  is not Turing-recognizable.  
 (b) False, e.g., if  $A = \{00, 11, 111\}$  and  $B = \{00, 11\}$ , then  $\overline{A} \cap B = \emptyset$  but  $A \neq B$ .  
 For  $A$  and  $B$  to be equal, we instead need  $(\overline{A} \cap B) \cup (A \cap \overline{B}) = \emptyset$ .  
 (c) False. A TM  $M$  may loop on input  $w$ .  
 (d) True, by Theorem 4.9.  
 (e) False, by Theorem 4.8.  
 (f) True, by Theorem 5.2.  
 (g) False, by Theorem 4.11.  
 (h) True, by Theorem 4.5.  
 (i) False, by slide 4-39.  
 (j) False, by Corollary 4.23.
2. (a) No, because  $f(x) = f(z) = 2$ .  
 (b) Yes, because  $f(y) = 1$  and  $f(x) = 2$ , so all members of  $B$  are hit by  $f$ .  
 (c) No, because  $f$  is not one-to-one.  
 (d) An algorithm is a Turing machine that always halts.  
 (e) A language  $L_1$  that is Turing-recognizable has a Turing machine  $M_1$  that may loop forever on a string  $w \notin L_1$ . A language  $L_2$  that is Turing-decidable has a Turing machine  $M_2$  that always halts.
3. (a)  $q_1 0 1 0 \# 1 \quad x q_2 1 0 \# 1 \quad x 1 q_2 0 \# 1 \quad x 1 0 q_2 \# 1 \quad x 1 0 \# q_4 1 \quad x 1 0 \# 1 q_{\text{reject}}$   
 (b)  $q_1 1 \# 1 \quad x q_3 \# 1 \quad x \# q_5 1 \quad x q_6 \# x \quad q_7 x \# x \quad x q_1 \# x \quad x \# q_8 x \quad x \# x q_8$   
 $x \# x \sqcup q_{\text{accept}}$
4. Homework 9, problem 1.
5. Define the language as
 
$$C = \{ \langle N, R \rangle \mid N \text{ is an NFA and } R \text{ is a regular expression with } L(N) = L(R) \}.$$
 Recall that the proof of Theorem 4.5 defines a Turing machine  $F$  that decides the language  $EQ_{\text{DFA}} = \{ \langle A, B \rangle \mid A \text{ and } B \text{ are DFAs and } L(A) = L(B) \}$ . Then the following Turing machine  $T$  decides  $C$ :
 
$$T = \text{“On input } \langle N, R \rangle, \text{ where } N \text{ is an NFA and } R \text{ is a regular expression:}$$
  1. Convert  $N$  and  $R$  into equivalent DFAs  $D_1$  and  $D_2$   
 using the algorithms in the proof of Kleene’s Theorem.
  2. Run TM  $F$  from Theorem 4.5 on input  $\langle D_1, D_2 \rangle$ .
  3. If  $F$  accepts, *accept*. If  $F$  rejects, *reject*.”
6. This is Theorem 5.1, whose proof is given on slide 5-8.