

Print family (or last) name: _____

Print given (or first) name: _____

I have read and understand all of the instructions below, and I will obey the University Policy on Academic Integrity.

Signature and Date

- This exam has 7 pages in total, numbered 1 to 7. Make sure your exam has all the pages.
- Note the number written on the upper right-hand corner of the first page. On the sign-up sheet being passed around, print your name next to this number.
- This exam will be 1 hour and 20 minutes in length.
- This is a closed-book, closed-note exam. Electronic devices (e.g., cellphone, smart watch, calculator) are not allowed.
- For all problems, follow these instructions:
 1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the exam sheets to work out your answers before filling in the answer space.
 2. DFA stands for deterministic finite automaton; NFA stands for nondeterministic finite automaton; PDA stands for push-down automaton; CFG stands for context-free grammar.
 3. For any state machines that you draw, you must include **all states and transitions**.
 4. For any proofs, be sure to provide a step-by-step argument, with justifications for every step. If you are asked to prove a result X, you may use in your proof of X any other result Y without proving Y. However, make it clear what the other result Y is that you are using; e.g., write something like, “By the result that $A^{**} = A^*$, we know that”

Problem	1	2	3	4	5	Total
Points						

1. [20 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

- (a) TRUE FALSE — Every context-free language is also regular.
- (b) TRUE FALSE — The class of regular languages is closed under intersection.
- (c) TRUE FALSE — A regular expression for $A = \{0^n 1^n 0^n \mid n \geq 0\}$ is $0^* 1^* 0^*$.
- (d) TRUE FALSE — If A has a regular expression, then A must have a PDA.
- (e) TRUE FALSE — If a language A has a PDA, then A must have a context-free grammar in Chomsky normal form.
- (f) TRUE FALSE — For an NFA $N = (Q, \Sigma, \delta, q_0, F)$, its transition function has the form $\delta : Q \times \Sigma_\epsilon \rightarrow Q$.
- (g) TRUE FALSE — If a language A is nonregular, then A has an NFA.
- (h) TRUE FALSE — If $A \subseteq B$ and B is a context-free language, then A must be context-free.
- (i) TRUE FALSE — If $A \subseteq B$ and A is a regular language, then B must be regular.
- (j) TRUE FALSE — If A is regular, then A must be finite.

2. [20 points] Give short answers to each of the following parts. Each answer should be at most a few sentences. Be sure to define any notation that you use.

- (a) Let $\Sigma = \{0, 1\}$. For any string $w \in \Sigma^*$, let $n_0(w)$ denote the number of 0's in w , and let $n_1(w)$ denote the number of 1's in w . Define the language $A = \{w \in \Sigma^* \mid n_0(w) \leq 1, n_1(w) \geq 2\}$. Give a DFA for A .

Answer:

- (b) Give a regular expression for the language $A = \{w \in \Sigma^* \mid n_0(w) \leq 1, n_1(w) \geq 2\}$, where $\Sigma = \{0, 1\}$.

Answer: _____

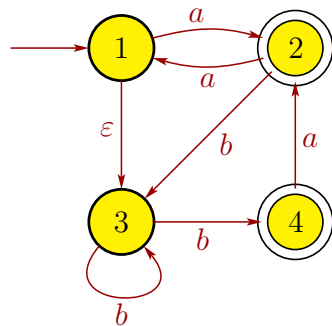
(c) Suppose a language A is generated by a context-free grammar $G = (V, \Sigma, R, S)$. Give a context-free grammar G' for A^* in terms of G . You do not have to prove the correctness of your CFG G' , but do not just give an example.

(d) Suppose that we are in the process of converting a CFG G with $\Sigma = \{0, 1\}$ into Chomsky normal form. We have already applied some steps in the process, and we currently have the following CFG:

$$\begin{aligned} S_0 &\rightarrow S \\ S &\rightarrow 1SA0A \mid 0AS1S \mid \varepsilon \\ A &\rightarrow 10S1 \mid \varepsilon \end{aligned}$$

In the next step, we want to remove the ε -rule $A \rightarrow \varepsilon$. Give the CFG after carrying out just this one step.

3. [15 points] Let N be the following NFA with $\Sigma = \{a, b\}$, and let $C = L(N)$.



Give a DFA for C .

Scratch-work area

4. [30 points] Consider the alphabet $\Sigma = \{a, b\}$ and the language

$$L = \{b^i a^j \mid i \geq j \geq 0\}.$$

(a) Give a context-free grammar G for L . Be sure to specify G as a 4-tuple $G = (V, \Sigma, R, S)$.

(b) Give a PDA for L . You only need to draw the graph.

Scratch-work area

5. [15 points] Recall the pumping lemma for regular languages:

Theorem: If L is a regular language, then there exists a pumping length p where, if $s \in L$ with $|s| \geq p$, then s can be split into three pieces $s = xyz$ such that (i) $xy^iz \in L$ for each $i \geq 0$, (ii) $|y| \geq 1$, and (iii) $|xy| \leq p$.

Let $A = \{ b^i a^j \mid i \geq j \geq 0 \}$. Is A a regular or nonregular language? If A is regular, give a regular expression for A . If A is not regular, prove that it is a nonregular language.

Circle one: **Regular Language** **Nonregular Language**