

Midterm Exam 1

CS 341-008: Foundations of Computer Science II — **Spring 2021, hybrid section**

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Print family (or last) name: \_\_\_\_\_

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I have read and understand all of the instructions below, and I will obey the University Policy on Academic Integrity.

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Signature and Date

- This exam has 7 pages in total, numbered 1 to 7. Make sure your exam has all the pages.
- Note the number written on the upper right-hand corner of the first page. On the sign-up sheet being passed around, print your name next to this number.
- This exam will be 1 hour and 20 minutes in length.
- This is a closed-book, closed-note exam. Electronic devices (e.g., cellphone, smart watch, calculator) are not allowed.
- For all problems, follow these instructions:
  1. Give only your answers in the spaces provided. I will only grade what you put in the answer space, and I will take off points for any scratch work in the answer space. Use the scratch-work area or the backs of the exam sheets to work out your answers before filling in the answer space.
  2. DFA stands for deterministic finite automaton; NFA stands for nondeterministic finite automaton; PDA stands for push-down automaton; CFG stands for context-free grammar.
  3. For any state machines that you draw, you must include **all states and transitions**.
  4. For any proofs, be sure to provide a step-by-step argument, with justifications for every step. If you are asked to prove a result X, you may use in your proof of X any other result Y without proving Y. However, make it clear what the other result Y is that you are using; e.g., write something like, “By the result that  $A^{**} = A^*$ , we know that . . . .”

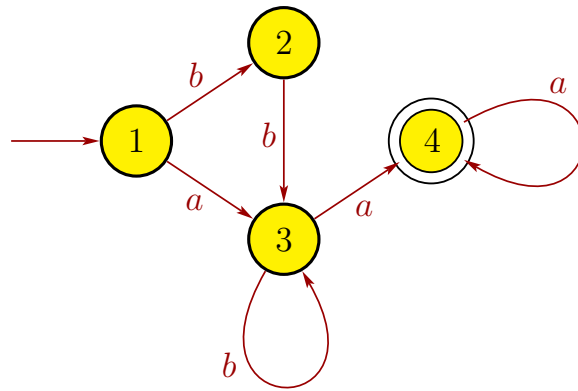
Problem	1	2	3	4	5	Total
Points						

1. [10 points] For each of the following, circle TRUE if the statement is correct. Otherwise, circle FALSE

- (a) TRUE FALSE — If  $A$  is a regular language, then  $A$  must be finite.
- (b) TRUE FALSE — The class of languages recognized by NFAs is closed under complementation.
- (c) TRUE FALSE — If  $A \subseteq B$  and  $A$  is a nonregular language, then  $B$  must be nonregular.
- (d) TRUE FALSE — If  $A$  and  $B$  are context-free languages, then  $A \cap B$  must also be context-free.
- (e) TRUE FALSE — Every NFA recognizes a nonregular language.
- (f) TRUE FALSE — If  $A$  is a context-free language, then  $A$  must also be regular.
- (g) TRUE FALSE — If a language  $A$  is recognized by an NFA, then  $A$  must have a context-free grammar in Chomsky normal form.
- (h) TRUE FALSE — If  $A \subseteq B$  and  $B$  is a context-free language, then  $A$  must be context-free.
- (i) TRUE FALSE — If  $A$  is a context-free language and  $B$  is a non-context-free language, then  $A \circ B$  must be non-context-free.
- (j) TRUE FALSE — A regular expression for  $A = \{ 1^n 0^n \mid n \geq 0 \}$  is  $1^*0^*$ .

2. [40 points] Give short answers to each of the following parts. Each answer should be at most a few sentences. Be sure to define any notation that you use.

(a) Give a regular expression for the language recognized by the NFA below.



(b) Suppose a language  $A_1$  is generated by a context-free grammar  $G_1 = (V_1, \Sigma, R_1, S_1)$ , where  $V_1$  is the set of variables,  $\Sigma$  is the alphabet of terminals,  $R_1$  is the set of rules, and  $S_1$  is the start variable. Give a context-free grammar  $G_2$  for  $A_1^*$  in terms of  $G_1$ . You do not have to prove the correctness of your CFG  $G_2$ , but do not just give an example.

- (c) Let  $M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$  be a DFA with language  $A_1$ , and  $M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$  be a DFA with language  $A_2$ . Consider the language  $A = A_1 \cup A_2$ . Give a DFA  $M_3$  for  $A$  in terms of  $M_1$  and  $M_2$ . Your DFA  $M_3$  must be completely general. Do not prove the correctness of your DFA  $M_3$ , but do not just give an example.

- (d) Suppose that we are in the process of converting a CFG  $G$  with  $\Sigma = \{0, 1\}$  into Chomsky normal form. We have already applied some steps in the process, and we currently have the following CFG:

$$\begin{aligned} S_0 &\rightarrow S \\ S &\rightarrow 10A1SA \mid A101 \mid \varepsilon \\ A &\rightarrow 110A0 \mid \varepsilon \end{aligned}$$

In the next step, we want to remove the  $\varepsilon$ -rule  $A \rightarrow \varepsilon$ . Give the CFG after carrying out just this one step.

3. [20 points] For  $\Sigma = \{a, b\}$ , we say that a string contains a *double letter* if it contains  $aa$  or  $bb$  as a substring. Consider the language

$$C = \{w \in \Sigma^* \mid w \text{ does not end in a double letter}\}.$$

- (a) Give a 5-tuple description for a DFA for  $C$ . Be sure to explicitly define each part of the 5-tuple for your DFA for  $C$ .

- (b) Give a regular expression for  $C$ .

4. [15 points] Consider the language

$$D = \{ b^i a^j b^k \mid i, j, k \geq 0, \text{ and } k = i + j \}.$$

Give a context-free grammar  $G$  for  $D$ . Be sure to specify  $G$  as a 4-tuple  $G = (V, \Sigma, R, S)$ .

5. [15 points] Recall the pumping lemma for regular languages:

**Theorem:** If  $L$  is a regular language, then there exists a pumping length  $p$  where, if  $s \in L$  with  $|s| \geq p$ , then  $s$  can be split into three pieces  $s = xyz$  such that (i)  $xy^iz \in L$  for each  $i \geq 0$ , (ii)  $|y| \geq 1$ , and (iii)  $|xy| \leq p$ .

Let  $D = \{b^i a^j b^k \mid i, j, k \geq 0, \text{ and } k = i + j\}$ . Is  $D$  a regular or nonregular language? If  $D$  is regular, give a regular expression for  $D$ . If  $D$  is not regular, prove that it is a nonregular language.

Circle one:

Regular Language

Nonregular Language