More Test Input Generation

Martin Kellogg

Q1: Which of the following does EvoSuite have in common with AFL? (select all that apply)

- **A.** its core algorithm is a genetic algorithm
- **B.** it derives oracles using mutation testing
- **C.** it can operate directly on the bytecode of a project, without the need to recompile the code

Q2: **TRUE** or **FALSE**: EvoSuite treats the String class in Java specially

Q1: Which of the following does EvoSuite have in common with AFL? (select all that apply)

- A. its core algorithm is a genetic algorithm
- **B.** it derives oracles using mutation testing
- **C.** it can operate directly on the bytecode of a project, without the need to recompile the code

Q2: **TRUE** or **FALSE**: EvoSuite treats the String class in Java specially

Q1: Which of the following does EvoSuite have in common with AFL? (select all that apply)

- A. its core algorithm is a genetic algorithm
- **B.** it derives oracles using mutation testing
- **C.** it can operate directly on the bytecode of a project, without the need to recompile the code

Q2: **TRUE** or **FALSE**: EvoSuite treats the String class in Java specially

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 o concolic testing

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 concolic testing

• fuzzing isn't the only way to randomly generate tests

- fuzzing isn't the only way to randomly generate tests
- Randoop (Pacheco et al. 2007) is a good example of an alternative approach: *feedback-directed* random testing

- fuzzing isn't the only way to randomly generate tests
- Randoop (Pacheco et al. 2007) is a good example of an alternative approach: *feedback-directed* random testing
 - does **not** use a genetic algorithm

- fuzzing isn't the only way to randomly generate tests
- Randoop (Pacheco et al. 2007) is a good example of an alternative approach: *feedback-directed* random testing
 - does **not** use a genetic algorithm
 - but the core idea is similar: build test inputs incrementally
 - that is, new inputs extend existing inputs

- fuzzing isn't the only way to randomly generate tests
- Randoop (Pacheco et al. 2007) is a good example of an alternative approach: *feedback-directed* random testing
 - does **not** use a genetic algorithm
 - but the core idea is similar: build test inputs incrementally
 - that is, new inputs extend existing inputs
 - execute each new input **immediately** (but there is no explicit fitness function: it is not designed as a genetic algorithm)

- fuzzing isn't the only way to randomly generate tests
- Randoop (Pacheco et al. 2007) is a good example of an alternative approach: *feedback-directed* random testing
 - does **not** use a genetic algorithm
 - but the core idea is similar: build test inputs **incrementally**
 - that is, new inputs extend existing inputs
 - execute each new input **immediately** (but there is no explicit fitness function: it is not designed as a genetic algorithm)
 - tests are discarded if they do not discover new states

• inputs:

• output:

- inputs:
 - classes under test (this tool targets Java/OOP)

• output:

- inputs:
 - classes under test (this tool targets Java/OOP)
 - \circ time limit

• output:

- inputs:
 - classes under test (this tool targets Java/OOP)
 - time limit
 - a set of contracts to use as oracles
 - e.g., "o.equals(o) == true" or "o.hashCode() never throws"
- output:

- inputs:
 - classes under test (this tool targets Java/OOP)
 - time limit
 - a set of contracts to use as oracles

e.g., "o.equals(o) == true" or "o.hashCode() never throws"

- output:
 - sequences of method calls that cause a contract violation

- inputs:
 - classes under test (t
 - time limit
 - a set of contracts to■ e.g., "o.equals(o)
- output:
 - sequences of metho

```
For example:
Map h = new HashMap();
Collection c = h.values();
Object[] a = c.toArray();
List l = new LinkedList();
l.addFirst(a);
Set t = new TreeSet(1);
Set u =
Collections.unmodifiableSet(t);
assertTrue(u.equals(u));
```

- inputs:
 - classes under test (t
 - time limit
 - a set of contracts to■ e.g., "o.equals(o)
- output:
 - sequences of metho

```
For example:
```

```
Map h = new HashMap();
Collection c = h.values();
Object[] a = c.toArray();
```

```
List l = new LinkedList();
```

```
l.addFirst(a);
```

```
Set t = new TreeSet(1);
```

```
Set u =
```

```
Collections.unmodifiableSet(t);
```

assertTrue(u.equals(u));

fails when executed ²

• how does it work?

- how does it work?
 - o start with a set of seed sequences of size 1 (e.g., int i = 0;)

- how does it work?
 - o start with a set of seed sequences of size 1 (e.g., int i = 0;)
 - randomly select a method call $m(T_1, \ldots, T_k) / T_{ret}$ s.t. there is a sequence in the seed pool that ends in all T_i for $1 \le i \le k$

- how does it work?
 - o start with a set of seed sequences of size 1 (e.g., int i = 0;)
 - randomly select a method call $m(T_1, \ldots, T_k) / T_{ret}$ s.t. there is a sequence in the seed pool that ends in all T_i for $1 \le i \le k$
 - $\circ~$ for each ${\rm T_i}$, choose a sequence ${\rm S_i}$ that constructs an object ${\rm v_i}$ of type ${\rm T_i}$ from the pool

- how does it work?
 - o start with a set of seed sequences of size 1 (e.g., int i = 0;)
 - randomly select a method call $m(T_1, \ldots, T_k) / T_{ret}$ s.t. there is a sequence in the seed pool that ends in all T_i for $1 \le i \le k$
 - $\circ~$ for each ${\rm T}_{i}$, choose a sequence ${\rm S}_{i}$ that constructs an object ${\rm v}_{i}$ of type ${\rm T}_{i}$ from the pool
 - create a new sequence:

$$S_{new} = S_1; ...; S_k; T_{ret} v_{new} = m(v_1, ..., v_k);$$

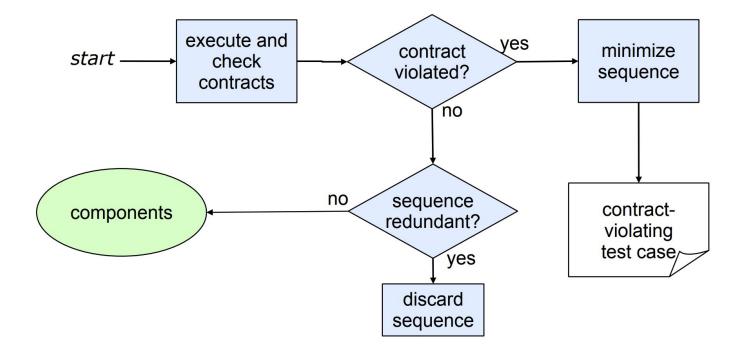
- how does it work?
 - o start with a set of seed sequences of size 1 (e.g., int i = 0;)
 - randomly select a method call $m(T_1, \ldots, T_k) / T_{ret}$ s.t. there is a sequence in the seed pool that ends in all T_i for $1 \le i \le k$
 - $\circ~$ for each ${\rm T}_{i}$, choose a sequence ${\rm S}_{i}$ that constructs an object ${\rm v}_{i}$ of type ${\rm T}_{i}$ from the pool
 - \circ create a new sequence:

$$\blacksquare S_{new} = S_1; ...; S_k; T_{ret} v_{new} = m(v_1, ..., v_k);$$

 classify the new sequence by executing it: may discard, output as a test case, or add it to the pool of sequences

Randoop: classifying sequences

Randoop: classifying sequences



• Randoop discards *redundant* sequences

- Randoop discards *redundant* sequences
 - during generation, it maintains a set O of all objects that it has ever created

- Randoop discards *redundant* sequences
 - during generation, it maintains a set O of all objects that it has ever created
 - a sequence is considered redundant if all of the objects created during its execution are members of O (using .equals)

- Randoop discards *redundant* sequences
 - during generation, it maintains a set O of all objects that it has ever created
 - a sequence is considered redundant if all of the objects created during its execution are members of O (using .equals)
 - Randoop would work with other reasonable definitions of redundant, too
 - e.g., heap canonicalization

• Randoop has been used to find real bugs in e.g., the JDK

- Randoop has been used to find real bugs in e.g., the JDK
- It has been **deployed at companies** (e.g., Microsoft)

- Randoop has been used to find real bugs in e.g., the JDK
- It has been deployed at companies (e.g., Microsoft)
- The tool is **still maintained** (so you could use it yourself)
 - <u>https://randoop.github.io/randoop/</u>

- Randoop has been used to find real bugs in e.g., the JDK
- It has been **deployed at companies** (e.g., Microsoft)
- The tool is still maintained (so you could use it yourself)
 https://randoop.github.io/randoop/
- It is commonly used in research papers as a *baseline*: that is, a method that any new technique is expected to outperform
 - Randoop is fast and easy enough to use that if a new technique cannot outperform it, it's probably not worth using!

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 o concolic testing

• Suppose you wanted to evaluate the quality of two truffle-sniffing pigs

- Suppose you wanted to evaluate the quality of two truffle-sniffing pigs
 - **Intuition**: test whether they can actually find truffles!

- Suppose you wanted to evaluate the quality of two truffle-sniffing pigs
 - Intuition: test whether they can actually find truffles!
- Test idea: hide some truffles in your backyard and see how many each pig finds

- Suppose you wanted to evaluate the quality of two truffle-sniffing pigs
 - Intuition: test whether they can actually find truffles!
- Test idea: hide some truffles in your backyard and see how many each pig finds
 - The pig that finds more of the hidden truffles in your backyard is assumed to find more real truffles in the wild

- Suppose you wanted to evaluate the quality of two truffle-sniffing pigs
 - Intuition: test whether they can actually find truffles!
- Test idea: hide some truffles in your backyard and see how many each pig finds
 - The pig that finds more of the hidden truffles in your backyard is assumed to find more real truffles in the wild
- Suppose you wanted to evaluate the quality of two bug-finding test suites ...

The Lens of Adversity: mutation testing

Definition: *Mutation testing* (or *mutation analysis*) is a test suite adequacy metric in which the quality of a test suite is related to the number of intentionally-added defects it finds

The Lens of Adversity: mutation testing

Definition: *Mutation testing* (or *mutation analysis*) is a test suite adequacy metric in which the quality of a test suite is related to the number of intentionally-added defects it finds

• Informally: "You claim your test suite is really great at finding security bugs? Well, I'll just **intentionally add a bug** to my source code and see if your test suite finds it!"

Definition: *Defect seeding* is the process of intentionally introducing a defect into a program.

Definition: *Defect seeding* is the process of intentionally introducing a defect into a program.

• The defect introduced is typically intentionally similar to defects introduced by real developers.

Definition: *Defect seeding* is the process of intentionally introducing a defect into a program.

- The defect introduced is typically intentionally similar to defects introduced by real developers.
- The seeding is typically done by changing the source code.

Definition: *Defect seeding* is the process of intentionally introducing a defect into a program.

- The defect introduced is typically intentionally similar to defects introduced by real developers.
- The seeding is typically done by changing the source code.
- For mutation testing, defect seeding is typically done automatically (given a model of what human bugs look like)

Mutation testing: mutation operators

Definition: A *mutation operator* systematically changes a program point. In mutation testing, the mutation operators are usually modeled on historical human defects.

Mutation testing: mutation operators

Definition: A *mutation operator* systematically changes a program point. In mutation testing, the mutation operators are usually modeled on historical human defects.

• Example mutations:

○ if (a < b)	\rightarrow if (a <= b)
○ if (a == b)	\rightarrow if (a != b)
\circ a = b + c	\rightarrow a = b - c
o f(); g();	→ g(); f();
о х = у	\rightarrow x = z

Mutation testing: mutants

Definition: A *mutant* (or *variant*) is a version of the original program produced by applying one or more mutation operators to one or more program locations.

Mutation testing: mutants

Definition: A *mutant* (or *variant*) is a version of the original program produced by applying one or more mutation operators to one or more program locations.

Definition: The *order* of a mutant is the number of mutation operators applied.

Mutation testing: mutants

Definition: A *mutant* (or *variant*) is a version of the original program produced by applying one or more mutation operators to one or more program locations.

Definition: The order of a mutant is the number of mutation operators applied.

Mutation testing: killing mutants

Definition: A test suite is said to *kill* (or *detect*, or *reveal*) a mutant if the mutant fails a test that the original passes.

Mutation testing: killing mutants

Definition: A test suite is said to *kill* (or *detect*, or *reveal*) a mutant if the mutant fails a test that the original passes.

• test suites that kill more mutants are generally considered better

Mutation testing: killing mutants

Definition: A test suite is said to *kill* (or *detect*, or *reveal*) a mutant if the mutant fails a test that the original passes.

- test suites that kill more mutants are generally considered better
- (sorry for all the vocabulary, but it's necessary to understand how EvoSuite works)

Mutation testing: more to come!

- My intention today is to give you a high-level idea of how mutation testing works
 - because EvoSuite (which you'll use for HW4) relies on it

Mutation testing: more to come!

- My intention today is to give you a high-level idea of how mutation testing works
 - because EvoSuite (which you'll use for HW4) relies on it
- We will discuss mutation testing in much more detail in two weeks
 - and you'll get a chance to try your hand at it in HW6

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 concolic testing

• much like AFL or other fuzzers, EvoSuite uses a **genetic algorithm** to evolve better tests

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the **test suite** as the individual

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the **test suite** as the individual
 - individual tests are themselves "chromosomes"

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the **test suite** as the individual
 - individual tests are themselves "chromosomes"
 - the whole population is made up of many test suites

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the **test suite** as the individual
 - individual tests are themselves "chromosomes"
 - the whole population is made up of many test suites
 - this simplifies crossover/parentage: just add/remove tests

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the test suite as the individual
 - individual tests are themselves "chromosomes"
 - the whole population is made up of many test suites
 - this simplifies crossover/parentage: just add/remove tests
- EvoSuite also uses mutation testing to produce oracles

- much like AFL or other fuzzers, EvoSuite uses a genetic algorithm to evolve better tests
 - however, EvoSuite views the test suite as the individual
 - individual tests are themselves "chromosomes"
 - the whole population is made up of many test suites
 - this simplifies crossover/parentage: just add/remove tests
- EvoSuite also uses mutation testing to produce oracles
 - **key idea**: assertions that **kill mutants** make good oracles
 - we'll come back to this idea next week

• EvoSuite emphasizes producing human-readable tests

- EvoSuite emphasizes producing human-readable tests
- EvoSuite's developers actually expect you to look at the tests that it produces, and to use them for regression testing

- EvoSuite emphasizes producing human-readable tests
- EvoSuite's developers actually expect you to look at the tests that it produces, and to use them for regression testing
- by contrast, AFL is looking to find bugs
 - leads to test inputs that aren't easy to understand!

EvoSuite: HW4 thoughts

- HW4 asks you to use EvoSuite to generate test suites for a Java library
- As you do, consider how EvoSuite:

EvoSuite: HW4 thoughts

- HW4 asks you to use EvoSuite to generate test suites for a Java library
- As you do, consider how EvoSuite:
 o differs from AFL as deployed in HW3

EvoSuite: HW4 thoughts

- HW4 asks you to use EvoSuite to generate test suites for a Java library
- As you do, consider how EvoSuite:
 - differs from AFL as deployed in HW3
 - compares to the sort of tests that you might write by hand

EvoSuite: HW4 thoughts

- HW4 asks you to use EvoSuite to generate test suites for a Java library
- As you do, consider how EvoSuite:
 - differs from AFL as deployed in HW3
 - compares to the sort of tests that you might write by hand
 - does it achieve its goal of creating useful regression suites?

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 o concolic testing

• we've seen coverage used as a fitness function for a fuzzer

 we've seen coverage used as a fitness function for a fuzzer
 but what if we just try to figure out which inputs would improve coverage directly?

we've seen coverage used as a fitness function for a fuzzer
 but what if we just try to figure out which inputs would

improve coverage directly?

• this is the key idea behind using **symbolic execution** to generate test inputs that improve coverage

Definition: *symbolic execution* abstractly runs the target program while computing a formula for each variable

Definition: *symbolic execution* abstractly runs the target program while computing a formula for each variable

• effectively, use math to figure out which values of each variable will cause the program to take particular paths

Definition: *symbolic execution* abstractly runs the target program while computing a formula for each variable

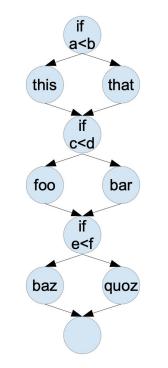
- effectively, use math to figure out which values of each variable will cause the program to take particular paths
- our plan: choose an uncovered bit of code, and then symbolically execute **backwards** from there to figure out what values the input variables would need to take on in order to cover the code

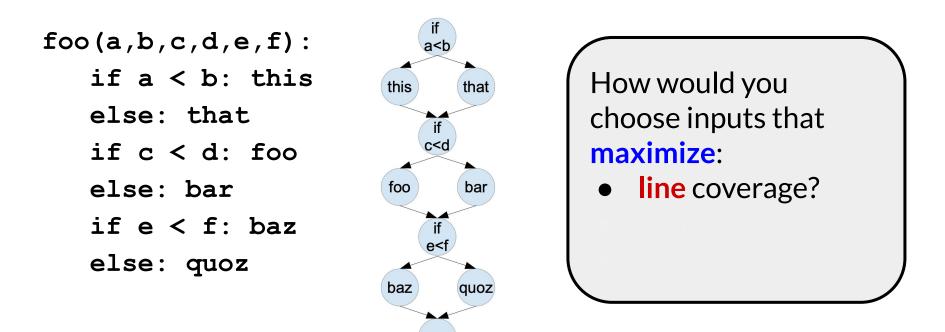
Definition: *symbolic execution* abstractly runs the target program while computing a formula for each variable

- effectively, use math to figure out which values of each variable will cause the program to take particular paths
- our plan: choose an uncovered bit of code, and then symbolically execute **backwards** from there to figure out what values the input variables would need to take on in order to cover the code
 - this is the Lens of Logic again, but applied in a different way

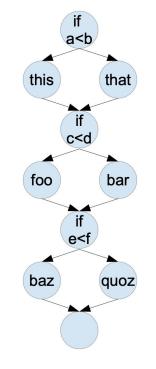
foo(a,b,c,d,e,f):
 if a < b: this
 else: that
 if c < d: foo
 else: bar
 if e < f: baz
 else: quoz</pre>

foo(a,b,c,d,e,f):
 if a < b: this
 else: that
 if c < d: foo
 else: bar
 if e < f: baz
 else: quoz</pre>





foo(a,b,c,d,e,f):
 if a < b: this
 else: that
 if c < d: foo
 else: bar
 if e < f: baz
 else: quoz</pre>

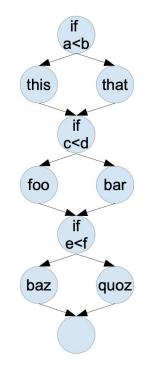


How would you choose inputs that **maximize**:

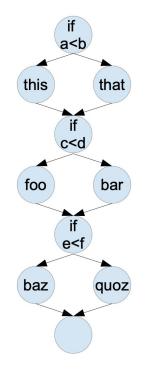
- line coverage?
- branch coverage?

if foo(a,b,c,d,e,f):a<b if a < b: this How would you this that choose inputs that else: that if c<d maximize: if c < d: foo line coverage? bar else: bar foo • • branch coverage? if e < f: baz if e<f • path coverage? else: quoz baz quoz

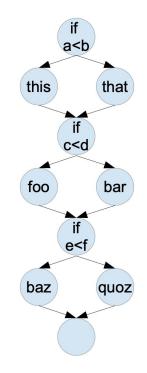
• If you have N sequential (or serial) if statements ...



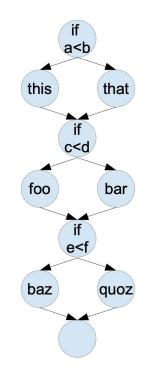
- If you have N sequential (or serial) if statements ...
- There are **2N** branch edges



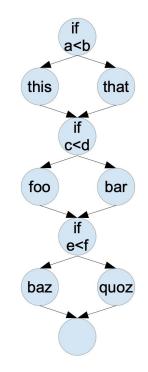
- If you have N sequential (or serial) if statements ...
- There are 2N branch edges
 - Which you could cover in 2 tests!



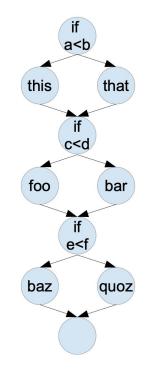
- If you have N sequential (or serial) if statements ...
- There are **2N** branch edges
 - Which you could cover in 2 tests!
 - One always goes left, one always right



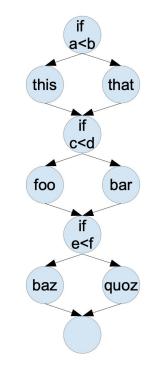
- If you have N sequential (or serial) if statements ...
- There are **2N** branch edges
 - Which you could cover in 2 tests!
 - One always goes left, one always right
- But there are 2^N paths



- If you have N sequential (or serial) if statements ...
- There are **2N** branch edges
 - Which you could cover in 2 tests!
 - One always goes left, one always right
- But there are 2^N paths
 - You need 2^N tests to cover them



- If you have N sequential (or serial) if statements ...
- There are **2N** branch edges
 - Which you could cover in 2 tests!
 - One always goes left, one always right
- But there are 2^N paths
 - You need 2^N tests to cover them
- Recall that path coverage subsumes branch coverage



• Consider generating test inputs to cover a path

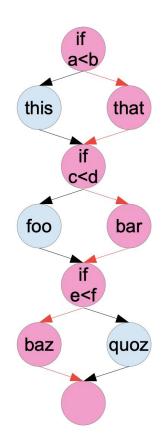
Consider generating test inputs to cover a path
 If we could do that, branch/statement/etc coverage is easy

- Consider generating test inputs to cover a path
 If we could do that, branch/statement/etc coverage is easy
- Key idea: solve this problem with math

- Consider generating test inputs to cover a path
 If we could do that, branch/statement/etc coverage is easy
- Key idea: solve this problem with math

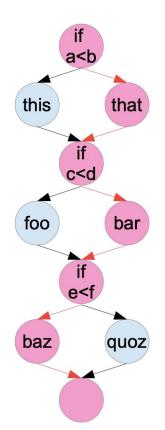
Definition: a *path predicate* (or *path condition*, or *path constraint*) is a boolean formula over program variables that is true when the program executes the given path

- Consider the highlighted (in pink) path
 i.e., "false, false, true"
- What is its path predicate?



- Consider the highlighted (in pink) path
 i.e., "false, false, true"
- What is its path predicate?

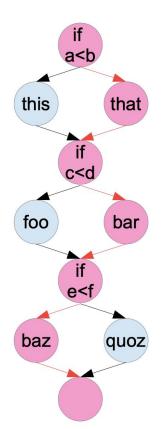
o a >= b && c >= d && e < f



- Consider the highlighted (in pink) path
 i.e., "false, false, true"
- What is its path predicate?

o a >= b && c >= d && e < f

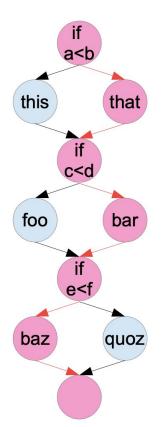
• When the path predicate is true, control flow will follow the given path



- Consider the highlighted (in pink) path
 i.e., "false, false, true"
- What is its path predicate?

o a >= b && c >= d && e < f

- When the path predicate is true, control flow will follow the given path
- So, given a path predicate, how do we choose a test input that covers the path?



Definition: A *satisfying assignment* is a mapping from variables to values that makes a predicate true.

Definition: A *satisfying assignment* is a mapping from variables to values that makes a predicate true.

• What is a satisfying assignment for

 \circ a >= b && c >= d && e < f?

Definition: A *satisfying assignment* is a mapping from variables to values that makes a predicate true.

• What is a satisfying assignment for

$$\circ$$
 a >= b && c >= d && e < f?

■ a=0, b=0, c=0, d=0, e=0, f=1

... many more

• How do we find satisfying assignments in general?

- How do we find satisfying assignments in general?
 - Option 1: ask humans
 - labor-intensive, slow, expensive, etc.

- How do we find satisfying assignments in general?
 - Option 1: ask humans
 - labor-intensive, slow, expensive, etc.
 - Option 2: repeatedly guess randomly
 - works surprisingly well (when answers are not sparse)

- How do we find satisfying assignments in general?
 - Option 1: ask humans
 - labor-intensive, slow, expensive, etc.
 - Option 2: repeatedly guess randomly
 - works surprisingly well (when answers are not sparse)
 - Option 3: use an *automated theorem prover*
 - cf. Wolfram Alpha, MatLab, Mathematica, Z3, etc.
 - works very well for a restricted class of equations (e.g., linear but not arbitrary polynomials, etc.; more detail in week 14)

- How do we find satisfying assignments in general?
 - Option 1: ask humans
 - labor-intensive, slow, expensive, etc.
 - Option 2: repeatedly guess randomly
 - works surprisingly well (when answers are not sparse)
 - Option 3: use an *automated theorem prover*
 - cf. Wolfram Alpha, MatLab, Mathematica, Z3, etc.
 - works very well for a restricted class of equations (e.g., linear but not arbitrary polynomials, etc.; more detail in week 14)

• Consider generating high-branch-coverage tests for a method:

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate
- For each path predicate, solve it

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate
- For each path predicate, solve it
 - $\circ~$ A solution is a satisfying assignment of values to input variables \rightarrow those are your test input

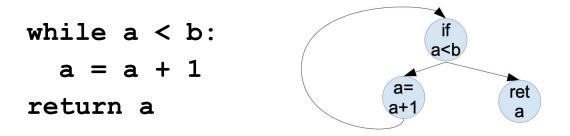
- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate
- For each path predicate, solve it
 - $\circ~$ A solution is a satisfying assignment of values to input variables \rightarrow those are your test input
 - None found? Dead code, tough predicate, etc.

Lens of Logic: enumerating paths

• What could go wrong with enumerating paths in a method?

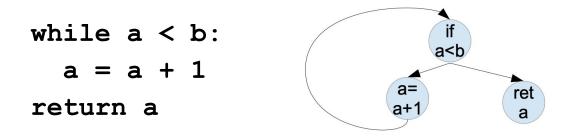
Lens of Logic: enumerating paths

- What could go wrong with enumerating paths in a method?
- There could be **infinitely many**!



Lens of Logic: enumerating paths

- What could go wrong with enumerating paths in a method?
- There could be **infinitely many**!



• One path corresponds to executing the loop once, another to twice, another to three times, etc.

• Key idea: don't enumerate all paths, approximate instead

- Key idea: don't enumerate all paths, approximate instead
- Typical Approximations:

- Key idea: don't enumerate all paths, approximate instead
- Typical Approximations:
 - Consider only acyclic paths (corresponds to taking each loop zero times or one time)

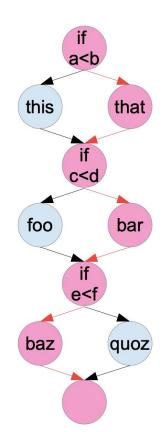
- Key idea: don't enumerate all paths, approximate instead
- Typical Approximations:
 - Consider only acyclic paths (corresponds to taking each loop zero times or one time)
 - Consider only taking each loop **at most** *k* times

- Key idea: don't enumerate all paths, approximate instead
- Typical Approximations:
 - Consider only acyclic paths (corresponds to taking each loop zero times or one time)
 - Consider only taking each loop **at most** *k* times
 - Enumerate paths breadth-first or depth-first and stop after k paths have been enumerated

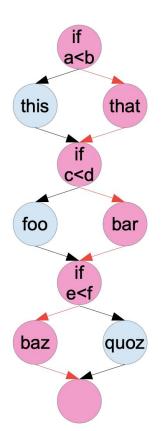
- Key idea: don't enumerate all paths, approximate instead
- Typical Approximations:
 - Consider only acyclic paths (corresponds to taking each loop zero times or one time)
 - Consider only taking each loop **at most** *k* times
 - Enumerate paths breadth-first or depth-first and **stop after** *k* paths have been enumerated
 - Concretely execute the program and see what it does (we'll come back to this later when we discuss concolic testing)

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate
- For each path predicate, solve it
 - $\circ~$ A solution is a satisfying assignment of values to input variables \rightarrow those are your test input
 - None found? Dead code, tough predicate, etc.

- Now we have a path through the program
- What could go wrong with **collecting** the path predicate?

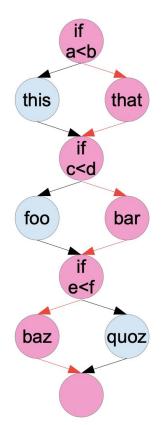


- Now we have a path through the program
- What could go wrong with collecting the path predicate?
 - The path predicate may not be expressible in terms of the inputs we control



- Now we have a path through the program
- What could go wrong with collecting the path predicate?
 - The path predicate may not be **expressible** in terms of the inputs we control

```
foo(a,b):
    str1 = read_from_url("abc.com")
    str2 = read_from_url("xyz.com")
    if (str1 == str2): bar()
```

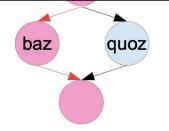


- Now we have a path through the program
- What could go wrong with collecting th predicate?
 - The path predicate may not be expr terms of the inputs we control

```
foo(a,b):
```

if (str1 == str2): bar()

Suppose we want to exercise the path that calls bar. One predicate is str1==str2. What do you assign to a and b?



a<b

• When we can't solve for a path predicate, what can we do?

When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence



- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence \rightarrow **no guarantee** either way

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence \rightarrow no guarantee either way
- So, we make a **best effort**:

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence \rightarrow no guarantee either way
- So, we make a **best effort**:
 - Collect the path predicates as best we can

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence \rightarrow no guarantee either way
- So, we make a **best effort**:
 - Collect the path predicates as best we can
 - Ask the solver to find a solution in terms of the input variables

- When we can't solve for a path predicate, what can we do?
 Ignore the problem (i.e., don't generate a test)
- Remember, testing can show the presence of bugs, but not their absence \rightarrow no guarantee either way
- So, we make a **best effort**:
 - Collect the path predicates as best we can
 - Ask the solver to find a solution in terms of the input variables
 - If it can't (because the math is too hard, we don't control the input, etc.), we give up

- Consider generating high-branch-coverage tests for a method:
- Enumerate "all" paths in the method
- For each path, **collect** the path predicate
- For each path predicate, solve it
 - $\circ~$ A solution is a satisfying assignment of values to input variables \rightarrow those are your test input
 - None found? Dead code, tough predicate, etc.

• Recall: we want to automatically generate test cases

- Recall: we want to automatically generate test cases
- We have an approach that works well in practice:
 - Enumerate some paths
 - Extract their path constraints
 - **Solve** those path constraints

Symbolic execution in practice

- symbolic execution was invented in the 1970s
 - but theorem provers of the time could rarely solve predicates, and the available hardware could enumerate few paths
- modern SMT solvers can handle the first problem, while second problem is less relevant due to Moore's Law

Aside: SMT solvers

• the path predicates I've used as examples today have mostly been **boolean formulas**

- the path predicates I've used as examples today have mostly been **boolean formulas**
 - how hard is it to check if a boolean formula is satisfiable?

- the path predicates I've used as examples today have mostly been **boolean formulas**
 - how hard is it to check if a boolean formula is satisfiable?
 - boolean satisfiability is the classic NP-complete problem

- the path predicates I've used as examples today have mostly been **boolean formulas**
 - how hard is it to check if a boolean formula is satisfiable?
 - boolean satisfiability is the classic NP-complete problem
- in practice, path predicates also include other kinds of expressions besides booleans
 - e.g., linear arithmetic, checking whether a pointer is null, etc.

- the path predicates I've used as examples today have mostly been **boolean formulas**
 - how hard is it to check if a boolean formula is satisfiable?
 - boolean satisfiability is the classic NP-complete problem
- in practice, path predicates also include other kinds of expressions besides booleans
 - e.g., linear arithmetic, checking whether a pointer is null, etc.
- an SMT solver is a **generalization** of a SAT solver that uses **theories** of other kinds of expressions to handle real programs
 - we'll come back to SMT solvers in week 14

- the path predicates I've used as examples today have mostly been **boolean formulas**
 - how hard is it to check if a boolean
 - boolean satisfiability is the clas
- in practice, path predicates also include besides booleans

Modern SMT solvers (e.g., Z3, cvc5) are **extraordinarily effective** at solving most instances (millions or billions of clauses in < 30s)

- \circ e.g., linear arithmetic, checking whe clauses in < 30s)
- an SMT solver is a generalization of a SAT solver that uses theories of other kinds of expressions to handle real programs
 we'll come back to SMT solvers in week 14

Symbolic execution in practice

- symbolic execution was invented in the 1970s
 - but theorem provers of the time could rarely solve predicates, and the available hardware could enumerate few paths
- modern SMT solvers can handle the first problem, while second problem is less relevant due to Moore's Law

• Moore's Law says that the available computing power for a given price increases exponentially

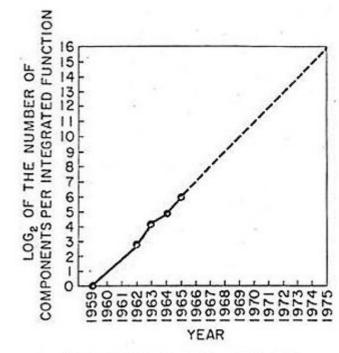
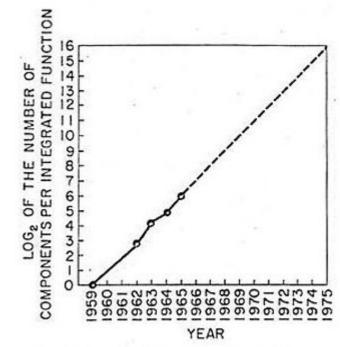


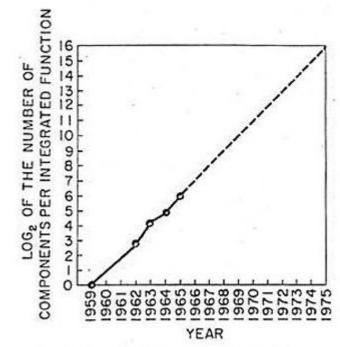
Fig. 2 Number of components per Integrated function for minimum cost per component extrapolated vs time.

- Moore's Law says that the available computing power for a given price increases exponentially
- not actually a law, but an observation that has been true basically since the invention of the transistor



'ig. 2 Number of components per Integrated function for minimum cost per component extrapolated vs time.

- Moore's Law says that the available computing power for a given price increases exponentially
- not actually a law, but an observation that has been true basically since the invention of the transistor
 - but might not last forever



^{&#}x27;ig. 2 Number of components per Integrated function for minimum cost per component extrapolated vs time.

- Moore's Law says that the available computing power for a given price increases exponentially
- not actually a law, but an observation that has been true basically since the invention of the transistor
 - but might not last forever
- implication: amount of computing power available today is huge compared to the 1970s

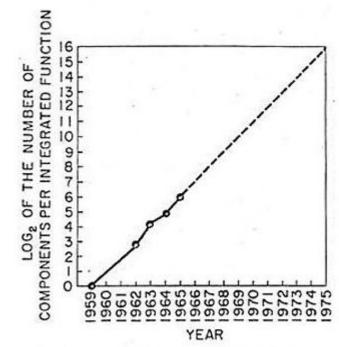


Fig. 2 Number of components per Integrated function for minimum cost per component extrapolated vs time.

Symbolic execution in practice

- symbolic execution was invented in the 1970s
 - but theorem provers of the time could rarely solve predicates, and the available hardware could enumerate few paths
- modern SMT solvers can handle the first problem, while second problem is less relevant due to Moore's Law
- implication: symbolic execution has been widely deployed in industry since the early 2000s
 - e.g., PREfix (Microsoft), Coverity, KLEE

 the biggest strength of symbolic execution is that it produces no false positives

- the biggest strength of symbolic execution is that it produces no false positives
 - that is, every test it generates really does lead to a violation of whatever policy it is enforcing (e.g., really leads to a crash)

- the biggest strength of symbolic execution is that it produces no false positives
 - that is, every test it generates really does lead to a violation of whatever policy it is enforcing (e.g., really leads to a crash)
- there are two serious downsides:

- the biggest strength of symbolic execution is that it produces no false positives
 - that is, every test it generates really does lead to a violation of whatever policy it is enforcing (e.g., really leads to a crash)
- there are two serious downsides:
 - it is expensive (recall: it relies on solving an NP-complete problem repeatedly!)

- the biggest strength of symbolic execution is that it produces no false positives
 - that is, every test it generates really does lead to a violation of whatever policy it is enforcing (e.g., really leads to a crash)
- there are two serious downsides:
 - it is expensive (recall: it relies on solving an NP-complete problem repeatedly!)
 - it cannot cover many parts of programs (recall: solving path predicates is NP-complete, so solvers sometimes fail!)

More test input generation: agenda

- Other approaches that use random testing
 - "feedback-directed" random testing
 - brief introduction to mutation testing
 - EvoSuite: mutation testing + a genetic algorithm
- Lens of Logic: symbolic execution for test input generation
 concolic testing

Limits of Symbolic Execution

- however, symbolic execution has serious limitations
 - for example, consider the function to the right:

```
testme(int x, int y) {
    if (bbox(x) == y) {
        ERROR;
    } else {
        // OK
    }
```

Limits of Symbolic Execution

- however, symbolic execution has serious limitations
 - for example, consider the function to the right:
- if bbox(x) is uninterpretable, then symbolic execution cannot determine if the ERROR statement is reachable

```
testme(int x, int y) {
   if (bbox(x) == y) {
      ERROR;
   } else {
      // OK
```

}

Limits of Symbolic Execution

- however, symbolic execution has serious limitations
 - for example, consider the function to the right:
- if bbox(x) is uninterpretable, then symbolic execution cannot determine if the ERROR statement is reachable

```
testme(int x, int y) {
    if (bbox(x) == y) {
        ERROR;
    } else {
        // OK
    }
}
```

Key question: how could we get around this limitation?

Concolic Testing

Definition: *concolic testing* combines concrete execution of the program (via other test generation techniques) with symbolic execution

Concolic Testing

Definition: *concolic testing* combines concrete execution of the program (via other test generation techniques) with symbolic execution

• "concolic" is a portmanteau of "*conc*rete" and "symb*olic*"

Concolic Testing

Definition: *concolic testing* combines concrete execution of the program (via other test generation techniques) with symbolic execution

- "concolic" is a portmanteau of "*conc*rete" and "symb*olic*"
- **key idea**: when symbolic execution gets stuck, actually execute the program and record what values the uninterpretable code **actually produces**

 symbolic execution determines that bbox(x) is uninterpretable

```
testme(int x, int y) {
    if (bbox(x) == y) {
        ERROR;
    } else {
        // OK
    }
}
```

- symbolic execution determines that bbox(x) is uninterpretable
- choose a random value of x, then execute the program

```
testme(int x, int y) {
   if (bbox(x) == y) {
      ERROR;
   } else {
      // OK
   }
```

- symbolic execution determines that bbox(x) is uninterpretable
- choose a random value of x, then execute the program
- replace the call to bbox(x) with whatever it returned on the concrete execution

```
testme(int x, int y) {
    if (bbox(x) == y) {
        ERROR;
    } else {
        // OK
     }
```

- symbolic execution determines that bbox(x) is uninterpretable
- choose a random value of x, then execute the program
- replace the call to bbox(x) with whatever it returned on the concrete execution
- let symbolic execution solve for y

```
testme(int x, int y) {
    if (bbox(x) == y) {
        ERROR;
    } else {
        // OK
        // OK
```

HW4 in-class

- in today's in-class/homework, you'll run EvoSuite
 - in general, students usually report that this assignment is easier and less time-consuming than HW3
 - however, there are two full length papers to read for next week
 - so you'll have plenty to do this week :)