

NJIT

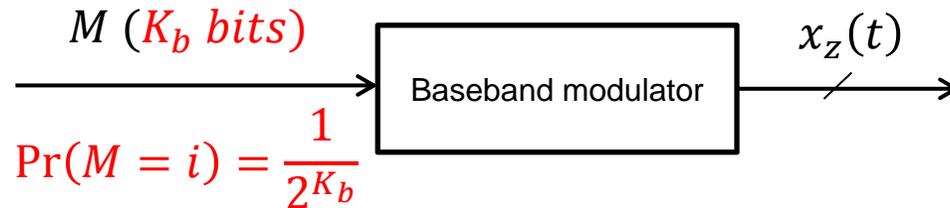


New Jersey's Science &  
Technology University

*THE EDGE IN KNOWLEDGE*

# Transmitting more than one bit (Ch. 14-15)

- Baseband modulator:



$$x_{z,i}(t) \quad \text{if } M = i$$

$$\text{with } E_i = \int |x_{z,i}(t)|^2 dt$$

$$\text{for } i = 0, 1, \dots, 2^{K_b} - 1$$

- Average energy:

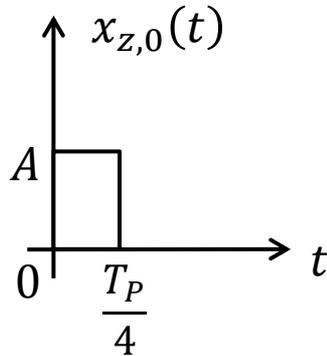
$$E_s = \frac{1}{2^{K_b}} \sum_{i=0}^{2^{K_b}-1} E_i = K_b E_b \quad \leftarrow \text{energy per bit}$$

- If  $x_z(t)$  is of duration  $T_P$ ,

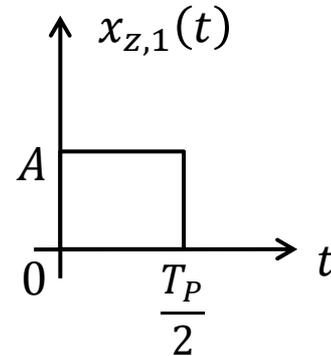
- Average power =  $P_s = \frac{E_s}{T_P}$

- Bit rate =  $\frac{K_b}{T_P}$

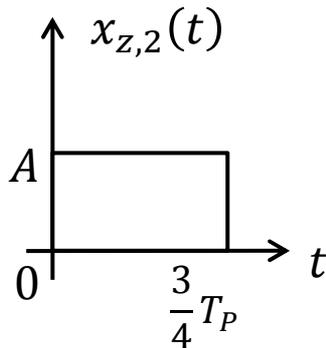
Ex.: a) Pulse Width Modulation (PWM) with  $K_b = 2$  (4-PWM)



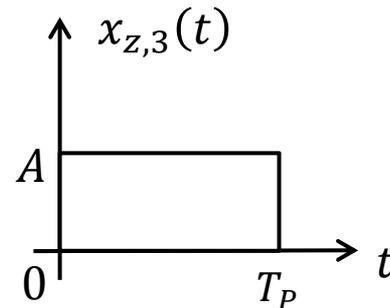
$$E_0 = A^2 \frac{T_P}{4}$$



$$E_1 = A^2 \frac{T_P}{2}$$



$$E_2 = A^2 \frac{3}{4}T_P$$



$$E_3 = A^2 T_P$$

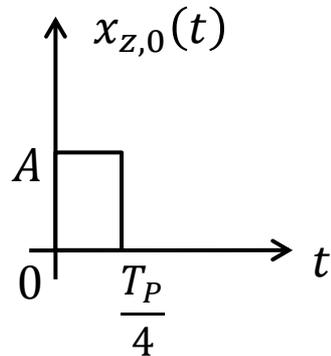
- Average energy:

$$\begin{aligned} E_s &= \frac{1}{4} A^2 T_P \left( \frac{1}{4} + \frac{1}{2} + \frac{3}{4} + 1 \right) \\ &= A^2 T_P \frac{10}{16} \end{aligned}$$

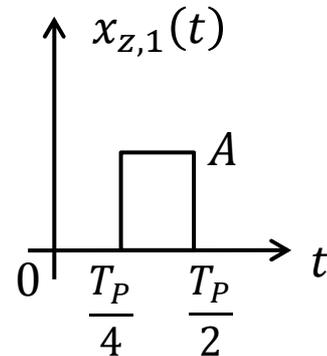
- We now impose that  $E_s = 2E_b$  since  $K_b = 2$  in order to find  $A$ :

$$A^2 T_P \frac{10}{16} = 2E_b \Rightarrow A = \sqrt{\frac{16 E_b}{5 T_P}}$$

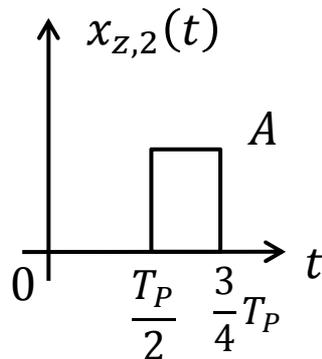
b) Pulse Position Modulation (PPM) with  $K_b = 2$  (4-PPM)



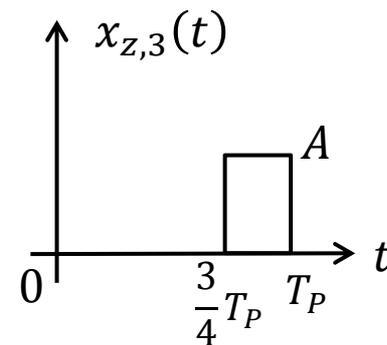
$$E_0 = A^2 \frac{T_P}{4}$$



$$E_1 = A^2 \frac{T_P}{2}$$



$$E_2 = A^2 \frac{T_P}{4}$$



$$E_3 = A^2 \frac{T_P}{4}$$

- Average energy:

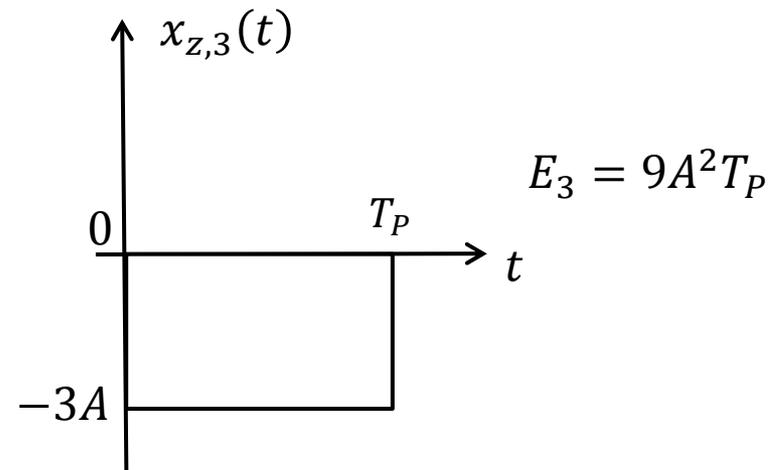
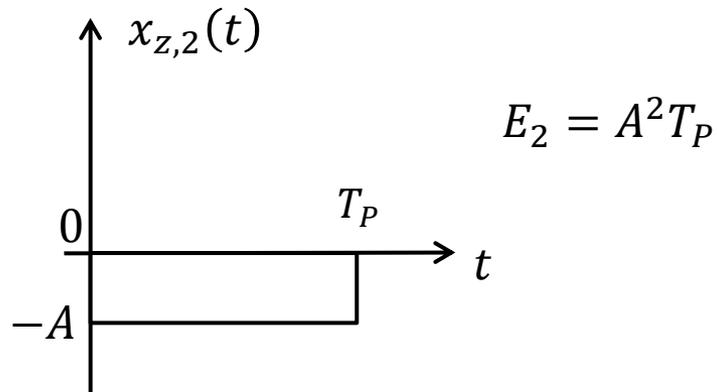
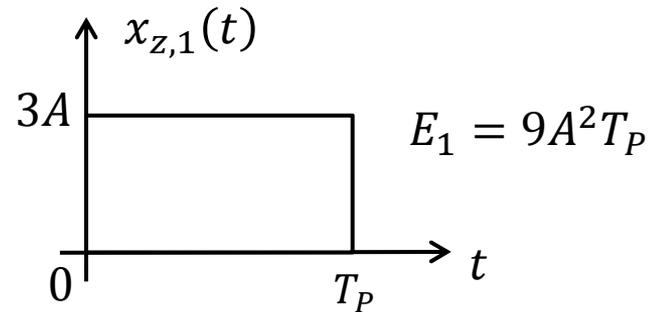
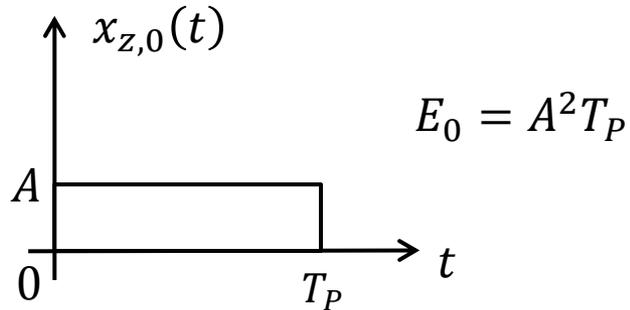
$$E_s = A^2 \frac{T_P}{4}$$

- We now impose that  $E_s = 2E_b$  since  $K_b = 2$  in order to find  $A$ :

$$A^2 \frac{T_P}{4} = 2E_b \Rightarrow A = \sqrt{\frac{8 E_b}{T_P}}$$

### c) Pulse Amplitude Modulation ( $2^{K_b}$ -PAM)

- As an example, 4-PAM with  $K_b = \log_2 4 = 2$ , we have



- Average energy:

$$E_s = \frac{A^2 T_P}{4} (2 \times 1 + 2 \times 9) = \frac{20}{4} A^2 T_P = 5A^2 T_P$$

- We now impose that  $E_s = 2E_b$  since  $K_b = 2$  in order to find  $A$ :

$$5A^2 T_P = 2E_b \Rightarrow A = \sqrt{\frac{2 E_b}{5 T_P}}$$

## d) Phase Shift Keying ( $2^{K_b}$ -PAM)

- $$x_{z,i}(t) = \begin{cases} Ae^{j\theta(i)} & 0 \leq t \leq T_P \\ 0 & \text{elsewhere} \end{cases} \quad i = 0, 1, \dots, 2^{K_b} - 1$$

with  $\theta(i) = \pi \frac{2i+1}{2^{K_b}}$

- As an example, if  $K_b = \log_2 4 = 2$ , we have

$$\theta(0) = \frac{\pi}{4} \Rightarrow x_{z,0}(t) = \begin{cases} \frac{1}{\sqrt{2}}(A + jA) & 0 \leq t \leq T_P \\ 0 & \text{elsewhere} \end{cases}$$

$$\theta(1) = \frac{3\pi}{4} \Rightarrow x_{z,1}(t) = \begin{cases} \frac{1}{\sqrt{2}}(-A + jA) & 0 \leq t \leq T_P \\ 0 & \text{elsewhere} \end{cases}$$

$$\theta(2) = \frac{5\pi}{4} \Rightarrow x_{z,2}(t) = \begin{cases} \frac{1}{\sqrt{2}}(-A - jA) & 0 \leq t \leq T_P \\ 0 & \text{elsewhere} \end{cases}$$

$$\theta(3) = \frac{7\pi}{4} \Rightarrow x_{z,3}(t) = \begin{cases} \frac{1}{\sqrt{2}}(A - jA) & 0 \leq t \leq T_P \\ 0 & \text{elsewhere} \end{cases}$$

- Imposing  $E_s = K_b E_b$ , we easily get

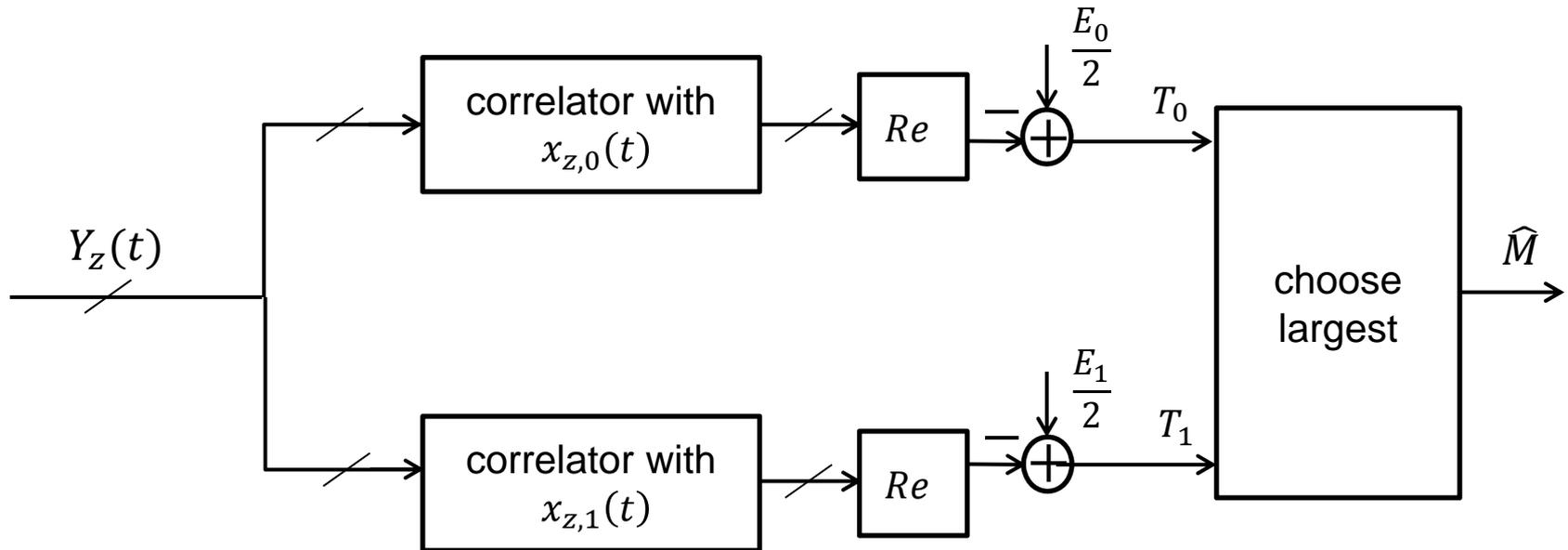
$$A = \sqrt{\frac{K_b E_b}{T_P}}$$



# Quick Quiz

Find  $x_{I,i}(t)$  and  $x_{Q,i}(t)$  for 4-PSK ( $i = 0,1,2,3$ ).

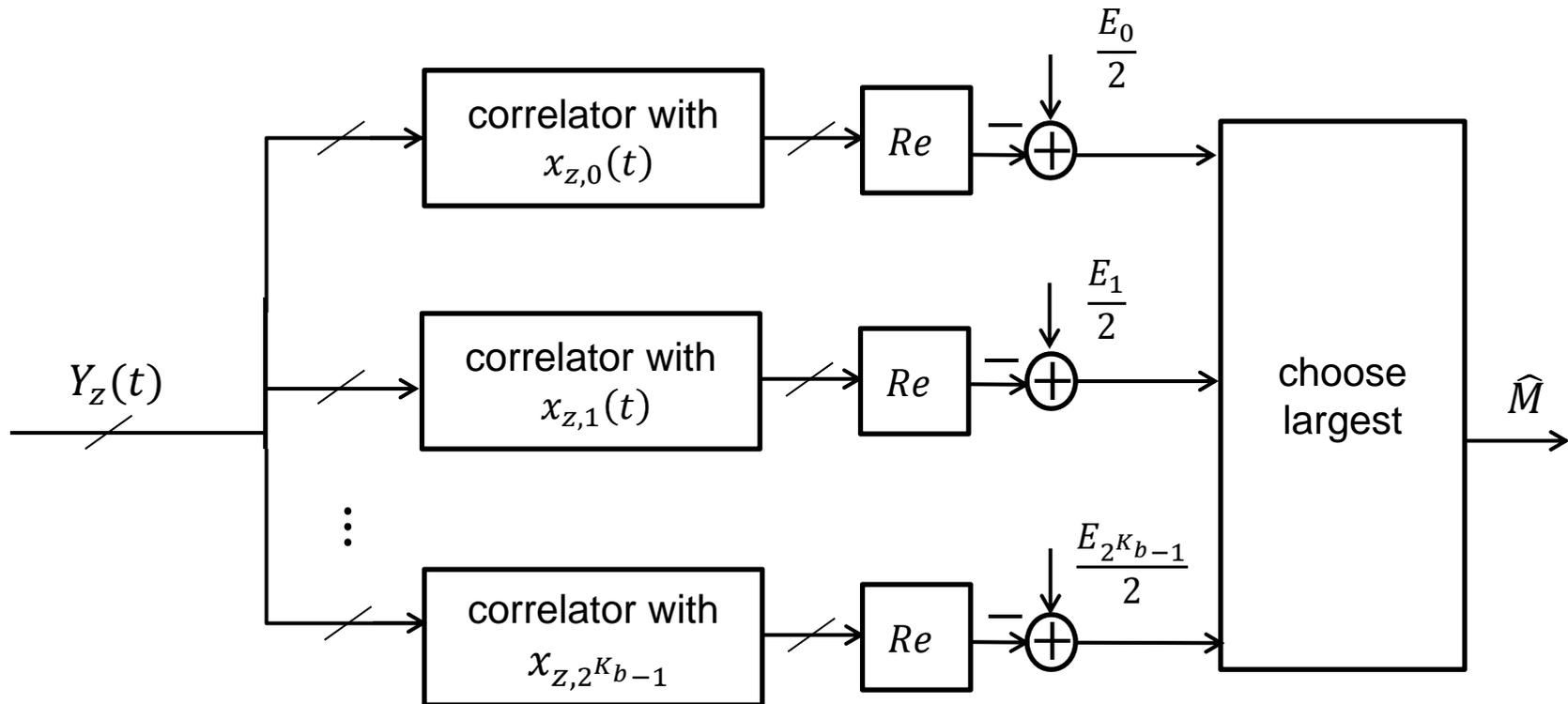
- Recall the optimal demodulator for  $K_b = 1$  (and  $\pi_i = \frac{1}{2}$ ):  
**MAXIMUM LIKELIHOOD BIT DEMODULATOR (MLBD)**



where  $T_j$  is the likelihood metric for message  $M = j$

- Generalizing to  $K_b \geq 1$  (and  $\pi_i = \frac{1}{2^{K_b}}$ ):

## MAXIMUM LIKELIHOOD WORD DEMODULATOR (MLWD)

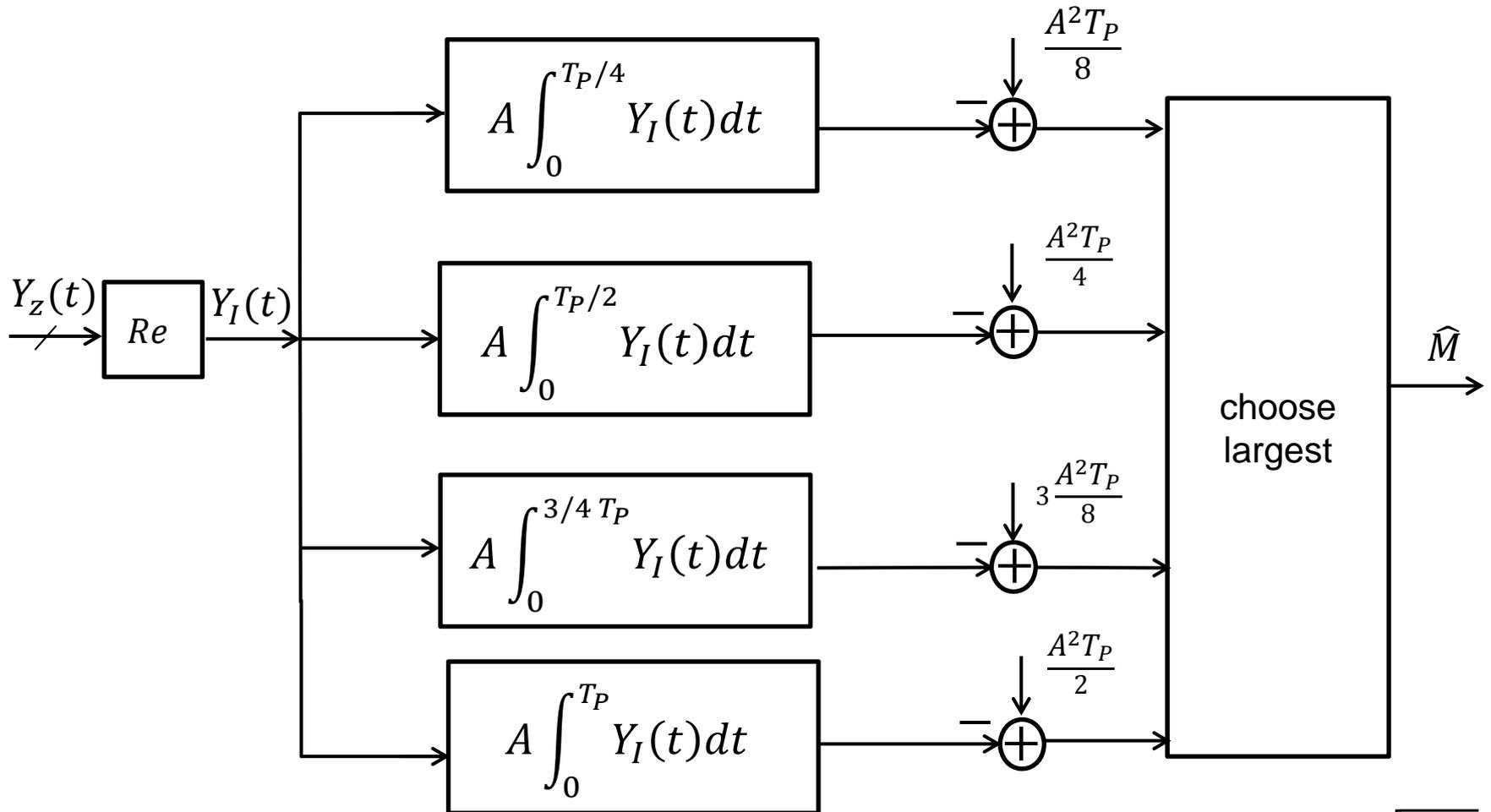


where  $T_j$  is the likelihood metric for message  $M = j$

## Remarks:

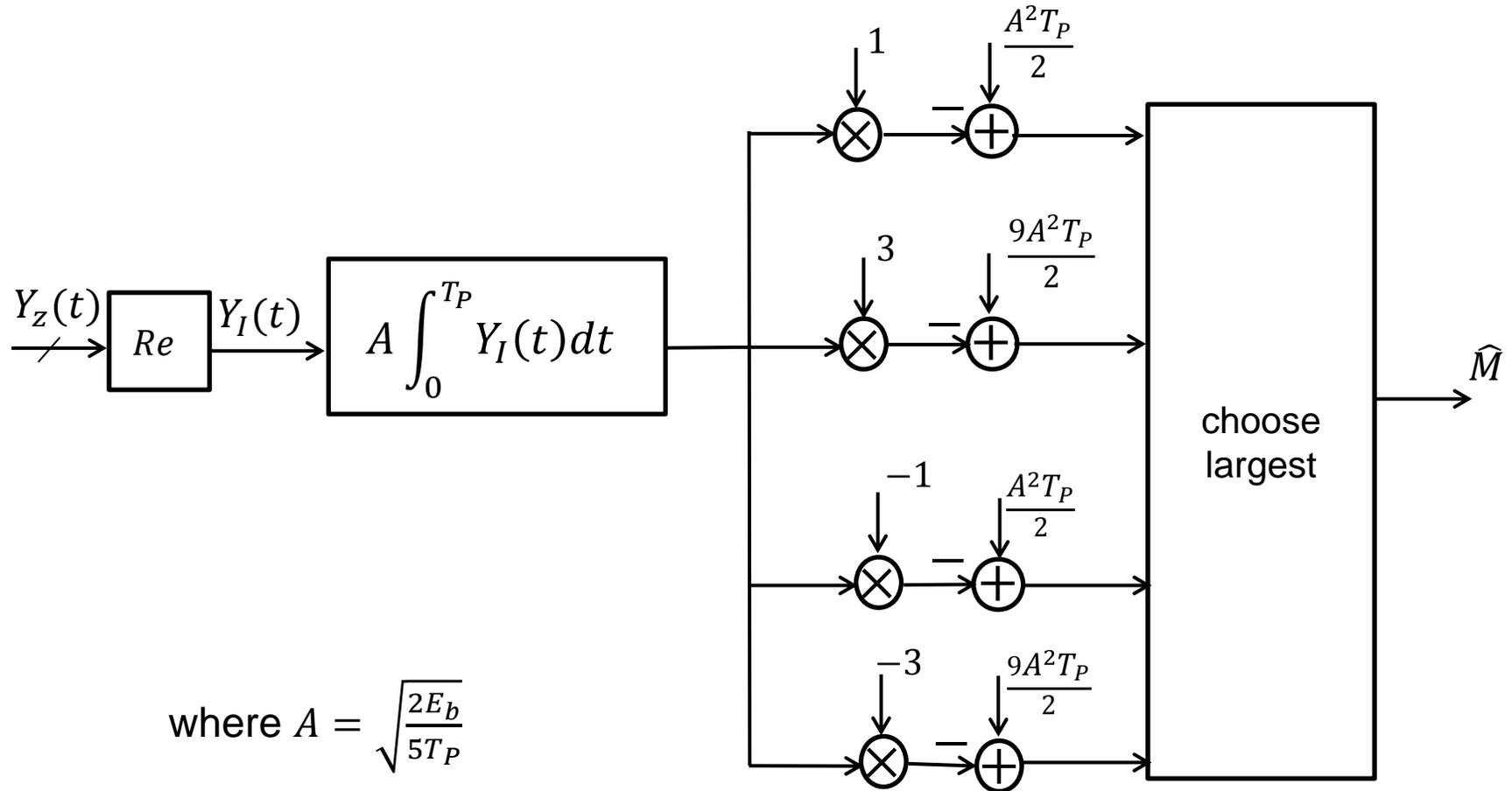
- $2^{K_b}$  branches  $\Rightarrow$  the complexity grows exponentially with  $K_b$
- This complexity can be mitigated if the waveforms have special properties, e.g., real or proportional to one another.

Ex.: a) 4-PWM ( $K_b = 2$ ) – The MLWD can be simplified as:



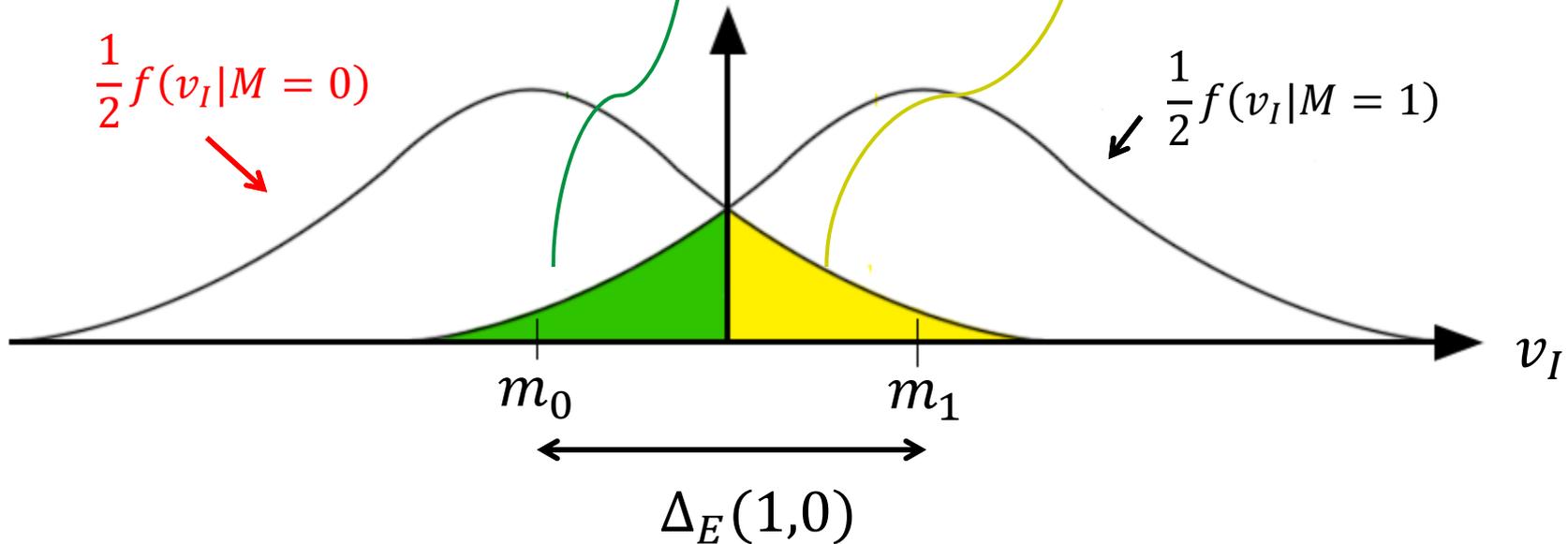
$$A = \sqrt{\frac{16 E_b}{5 T_P}}$$

b) 4-PAM – The MLWD can be simplified as:



- We are now interested in calculating the word error probability (WEP):  $P_W(E) = \Pr(\hat{M} \neq M)$
- Recall that with  $K_b = 1$ :

$$P_W(E) = P_B(E) = \underbrace{\frac{1}{2} \Pr(\hat{M} \neq 1 | M = 1)}_{\text{green}} + \underbrace{\frac{1}{2} \Pr(\hat{M} \neq 0 | M = 0)}_{\text{yellow}} = Q \left( \sqrt{\frac{\Delta_E(1,0)}{2N_0}} \right)$$



- Generalizing to  $K_b \geq 1$ , we get

$$P_W(E) = \frac{1}{2^{K_b}} \sum_{j=0}^{2^{K_b}-1} \Pr(\hat{M} \neq j | M = j)$$

- Calculating  $\Pr(\hat{M} \neq j | M = j)$  is more difficult than for  $K_b = 1$ :

$$\begin{aligned} \Pr(\hat{M} \neq j | M = j) &= \Pr(\underbrace{T_0 > T_j \text{ or } T_2 > T_j \text{ or } \dots \text{ or } T_{2^{K_b}-1} > T_j}_{\text{excluding } M = j} | M = j) \\ &= \Pr(\bigcup_{i \neq j} \{T_i > T_j\} | M = j) \end{aligned}$$

- We have:

$$\Pr(\hat{M} \neq j | M = j) = \Pr\left(\bigcup_{i \neq j} \{T_i > T_j\} | M = j\right)$$

$$\leq \sum_{\substack{i=0 \\ i \neq j}}^{2^{K_b}-1} \Pr(T_i > T_j | M = j)$$

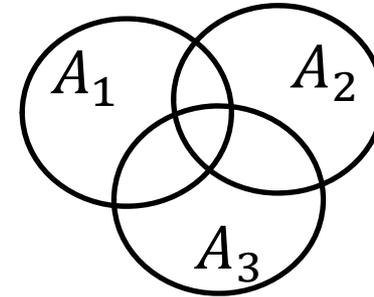
$$= \sum_{\substack{i=0 \\ i \neq j}}^{2^{K_b}-1} Q\left(\sqrt{\frac{\Delta_E(i, j)}{2N_0}}\right)$$

pairwise error probability

The union bound generally states that for any events  $A_i$ , the following holds:

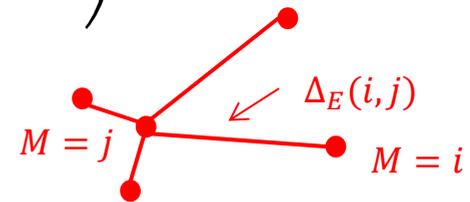
$$\Pr\left(\bigcup_i A_i\right) \leq \sum_i \Pr(A_i)$$

with equality if and only if the events  $A_i$  are disjoint.



- This leads to the so called union bound on  $P_W(E)$ :

$$P_W(E) \leq P_{WUB}(E) = \frac{1}{2^{K_b}} \sum_{j=0}^{2^{K_b}-1} \sum_{i \neq j} Q \left( \sqrt{\frac{\Delta_E(i,j)}{2N_0}} \right)$$



different signals  $i$  may have the same distance from a signal  $j$

$N_j$  = number of different distances from node  $j$

$$= \frac{1}{2^{K_b}} \sum_{j=0}^{2^{K_b}-1} \sum_{k=1}^{N_j} A_{dj}(k) Q \left( \sqrt{\frac{\Delta_{Ej}(k)}{2N_0}} \right)$$

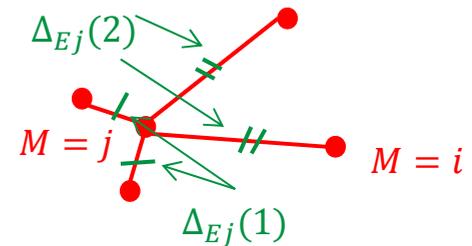
number of signals at distance  $\Delta_{Ej}(k)$  from  $j$ : conditional distance spectrum

grouping pairs of signals  $(i,j)$  with the same distance

$$= \frac{1}{2^{K_b}} \sum_{k=1}^N A_d(k) \operatorname{erfc} \left( \sqrt{\frac{\Delta_E(k)}{2N_0}} \right)$$

$N$  = number of different distances

number of pairs of signals at distance  $\Delta_E(k)$ : distance spectrum



- Remark: We have

$$\sum_{k=1}^{N_j} A_{dj}(k) = 2^{K_b} - 1$$

$$\sum_{k=1}^N A_d(k) = 2^{K_b} (2^{K_b} - 1)$$

(Why?)

Ex.: 4-PWM ( $K_b = 2$ )

- Conditional distance spectrum:

– for  $M = 0$

$$\Delta_E(0,1) = \int_{T_P/4}^{T_P/2} \frac{16 E_b}{5 T_P} dt = \frac{4E_b}{5}$$

$$\Delta_E(0,2) = \frac{8E_b}{5}$$

$$\Delta_E(0,3) = \frac{12E_b}{5}$$

$$\Rightarrow \left[ \left\{ \frac{4E_b}{5}, 1 \right\}, \left\{ \frac{8E_b}{5}, 1 \right\}, \left\{ \frac{12E_b}{5}, 1 \right\} \right] \quad (N_0 = 3)$$

$\Delta_{E0}(1)$        $A_{d0}(1)$

– for  $M = 1$  and  $M = 2$

$$\Rightarrow \left[ \left\{ \frac{4E_b}{5}, 2 \right\}, \left\{ \frac{8E_b}{5}, 1 \right\} \right] \quad (N_1 = N_2 = 2)$$

– for  $M = 3$

$$\Rightarrow \left[ \left\{ \frac{4E_b}{5}, 1 \right\}, \left\{ \frac{8E_b}{5}, 1 \right\}, \left\{ \frac{12E_b}{5}, 1 \right\} \right] \quad (N_3 = 3)$$

- Distance spectrum

$$\left[ \left\{ \frac{4E_b}{5}, 6 \right\}, \left\{ \frac{8E_b}{5}, 4 \right\}, \left\{ \frac{12E_b}{5}, 2 \right\} \right] \quad (N = 3)$$

$\Delta_E(1)$        $A_d(1)$

- Union bound:

$$\begin{aligned} P_W(E) \leq P_{WUB}(E) &= \frac{1}{4} \left( 6Q \left( \sqrt{\frac{4E_b}{10N_0}} \right) + 4Q \left( \sqrt{\frac{8E_b}{10N_0}} \right) + 2Q \left( \sqrt{\frac{12E_b}{10N_0}} \right) \right) \\ &= \frac{3}{2} Q \left( \sqrt{\frac{2E_b}{5N_0}} \right) + Q \left( \sqrt{\frac{4E_b}{5N_0}} \right) + \frac{1}{2} Q \left( \sqrt{\frac{6E_b}{5N_0}} \right) \end{aligned}$$



NJIT

*THE EDGE IN KNOWLEDGE*